

man pages section 5: Standards, Environments, and Macros

Sun Microsystems, Inc. 4150 Network Circle Santa Clara, CA 95054 U.S.A.

Part No: 817–0699–10 December 2003 Copyright 2003 Sun Microsystems, Inc. 4150 Network Circle, Santa Clara, CA 95054 U.S.A. All rights reserved.

This product or document is protected by copyright and distributed under licenses restricting its use, copying, distribution, and decompilation. No part of this product or document may be reproduced in any form by any means without prior written authorization of Sun and its licensors, if any. Third-party software, including font technology, is copyrighted and licensed from Sun suppliers.

Parts of the product may be derived from Berkeley BSD systems, licensed from the University of California. UNIX is a registered trademark in the U.S. and other countries, exclusively licensed through X/Open Company, Ltd.

Sun, Sun Microsystems, the Sun logo, docs.sun.com, AnswerBook, AnswerBook2, and Solaris are trademarks, registered trademarks, or service marks of Sun Microsystems, Inc. in the U.S. and other countries. All SPARC trademarks are used under license and are trademarks or registered trademarks of SPARC International, Inc. in the U.S. and other countries. Products bearing SPARC trademarks are based upon an architecture developed by Sun Microsystems, Inc.

The OPEN LOOK and Sun™ Graphical User Interface was developed by Sun Microsystems, Inc. for its users and licensees. Sun acknowledges the pioneering efforts of Xerox in researching and developing the concept of visual or graphical user interfaces for the computer industry. Sun holds a non-exclusive license from Xerox to the Xerox Graphical User Interface, which license also covers Sun's licensees who implement OPEN LOOK GUIs and otherwise comply with Sun's written license agreements.

Federal Acquisitions: Commercial Software-Government Users Subject to Standard License Terms and Conditions.

DOCUMENTATION IS PROVIDED "AS IS" AND ALL EXPRESS OR IMPLIED CONDITIONS, REPRESENTATIONS AND WARRANTIES, INCLUDING ANY IMPLIED WARRANTY OF MERCHANTABILITY, FITNESS FOR A PARTICULAR PURPOSE OR NON-INFRINGEMENT, ARE DISCLAIMED, EXCEPT TO THE EXTENT THAT SUCH DISCLAIMERS ARE HELD TO BE LEGALLY INVALID.

Copyright 2003 Sun Microsystems, Inc. 4150 Network Circle, Santa Clara, CA 95054 U.S.A. Tous droits réservés.

Ce produit ou document est protégé par un copyright et distribué avec des licences qui en restreignent l'utilisation, la copie, la distribution, et la décompilation. Aucune partie de ce produit ou document ne peut être reproduite sous aucune forme, par quelque moyen que ce soit, sans l'autorisation préalable et écrite de Sun et de ses bailleurs de licence, s'il y en a. Le logiciel détenu par des tiers, et qui comprend la technologie relative aux polices de caractères, est protégé par un copyright et licencié par des fournisseurs de Sun.

Des parties de ce produit pourront être dérivées du système Berkeley BSD licenciés par l'Université de Californie. UNIX est une marque déposée aux Etats-Unis et dans d'autres pays et licenciée exclusivement par X/Open Company, Ltd.

Sun, Sun Microsystems, le logo Sun, docs.sun.com, AnswerBook, AnswerBook2, et Solaris sont des marques de fabrique ou des marques déposées, ou marques de service, de Sun Microsystems, Inc. aux Etats-Unis et dans d'autres pays. Toutes les marques SPARC sont utilisées sous licence et sont des marques de fabrique ou des marques déposées de SPARC International, Inc. aux Etats-Unis et dans d'autres pays. Les produits portant les marques SPARC sont basés sur une architecture développée par Sun Microsystems, Inc.

L'interface d'utilisation graphique OPEN LOOK et Sun™ a été développée par Sun Microsystems, Inc. pour ses utilisateurs et licenciés. Sun reconnaît les efforts de pionniers de Xerox pour la recherche et le développement du concept des interfaces d'utilisation visuelle ou graphique pour l'industrie de l'informatique. Sun détient une licence non exclusive de Xerox sur l'interface d'utilisation graphique Xerox, cette licence couvrant également les licenciés de Sun qui mettent en place l'interface d'utilisation graphique OPEN LOOK et qui en outre se conforment aux licences écrites de Sun.

CETTE PUBLICATION EST FOURNIE "EN L'ETAT" ET AUCUNE GARANTIE, EXPRESSE OU IMPLICITE, N'EST ACCORDEE, Y COMPRIS DES GARANTIES CONCERNANT LA VALEUR MARCHANDE, L'APITITUDE DE LA PUBLICATION A REPONDRE A UNE UTILISATION PARTICULIERE, OU LE FAIT QU'ELLE NE SOIT PAS CONTREFAISANTE DE PRODUIT DE TIERS. CE DENI DE GARANTIE NE S'APPLIQUERAIT PAS, DANS LA MESURE OU IL SERAIT TENU JURIDIQUEMENT NUL ET NON AVENU.





030827@6671

Contents

Preface 7

Introduction13Intro(5)14

Standards, Environments, and Macros 15

ascii(5) 16 attributes(5) 18 26 charmap(5) crypt_bsdbf(5) 29 crypt_bsdmd5(5) 30 crypt_sunmd5(5) 31 crypt_unix(5) 32 dhcp(5) 33 dhcp_modules(5) 35 environ(5) 37 eqnchar(5) 42 extensions(5) 43 filesystem(5) 44 fnmatch(5) 63 fns(5) 67 fns_dns(5) 69 fns_files(5) 71 fns_initial_context(5) 73 $fns_nis+(5)$ 77

fns_nis(5) 79 fns_policies(5) 81 fns_references(5) 85 fns_x500(5) 88 91 formats(5) fsattr(5) 96 iconv_1250(5) 107 iconv_1251(5) 113 iconv(5) 122 iconv_646(5) 126 iconv_852(5) 129 iconv_8859-1(5) 135 iconv_8859-2(5) 141 iconv_8859-5(5) 147 iconv_dhn(5) 155 iconv_koi8-r(5) 159 iconv_mac_cyr(5) 167 iconv_maz(5) 175 iconv_pc_cyr(5) 179 iconv_unicode(5) 185 190 isalist(5) largefile(5) 192 196 lf64(5) lfcompile(5) 203 lfcompile64(5) 206 live_upgrade(5) 208 locale(5) 212 238 man(5)242 mansun(5) me(5) 246 mm(5) 251 ms(5)258 nfssec(5)263 pam_authtok_check(5) 265 pam_authtok_get(5) 267 pam_authtok_store(5) 269 pam_dhkeys(5) 270 pam_dial_auth(5) 272

⁴ man pages section 5: Standards, Environments, and Macros • December 2003

273 pam_krb5(5) pam_ldap(5) 278 pam_passwd_auth(5) 283 pam_projects(5) 285 pam_rhosts_auth(5) 286 pam_roles(5) 287 pam_sample(5) 289 pam_smartcard(5) 291 293 pam_unix(5) 296 pam_unix_account(5) pam_unix_auth(5) 297 pam_unix_session(5) 299 prof(5) 300 301 rbac(5) 304 regex(5) 313 regexp(5) 320 SEAM(5) 322 sgml(5) smartcard(5) 326 standards(5) 328 sticky(5) 333 term(5) 334 vgrindefs(5) 338 341 wbem(5)

Index 345

6 man pages section 5: Standards, Environments, and Macros • December 2003

Preface

Both novice users and those familar with the SunOS operating system can use online man pages to obtain information about the system and its features. A man page is intended to answer concisely the question "What does it do?" The man pages in general comprise a reference manual. They are not intended to be a tutorial.

Overview

The following contains a brief description of each man page section and the information it references:

- Section 1 describes, in alphabetical order, commands available with the operating system.
- Section 1M describes, in alphabetical order, commands that are used chiefly for system maintenance and administration purposes.
- Section 2 describes all of the system calls. Most of these calls have one or more error returns. An error condition is indicated by an otherwise impossible returned value.
- Section 3 describes functions found in various libraries, other than those functions that directly invoke UNIX system primitives, which are described in Section 2.
- Section 4 outlines the formats of various files. The C structure declarations for the file formats are given where applicable.
- Section 5 contains miscellaneous documentation such as character-set tables.
- Section 6 contains available games and demos.
- Section 7 describes various special files that refer to specific hardware peripherals and device drivers. STREAMS software drivers, modules and the STREAMS-generic set of system calls are also described.

- Section 9 provides reference information needed to write device drivers in the kernel environment. It describes two device driver interface specifications: the Device Driver Interface (DDI) and the Driver/Kernel Interface (DKI).
- Section 9E describes the DDI/DKI, DDI-only, and DKI-only entry-point routines a developer can include in a device driver.
- Section 9F describes the kernel functions available for use by device drivers.
- Section 9S describes the data structures used by drivers to share information between the driver and the kernel.

Below is a generic format for man pages. The man pages of each manual section generally follow this order, but include only needed headings. For example, if there are no bugs to report, there is no BUGS section. See the intro pages for more information and detail about each section, and man(1) for more information about man pages in general.

NAME	functior	ection gives the names of the commands or ons documented, followed by a brief otion of what they do.						
SYNOPSIS	functior in the st Options single le argume	This section shows the syntax of commands or functions. When a command or file does not exist in the standard path, its full path name is shown. Options and arguments are alphabetized, with single letter arguments first, and options with arguments next, unless a different argument order is required.						
	The foll section:	owing special characters are used in this						
	[]	Brackets. The option or argument enclosed in these brackets is optional. If the brackets are omitted, the argument must be specified.						
		Ellipses. Several values can be provided for the previous argument, or the previous argument can be specified multiple times, for example, "filename ".						
	I	Separator. Only one of the arguments separated by this character can be specified at a time.						
	{ }	Braces. The options and/or arguments enclosed within braces are interdependent, such that everything enclosed must be treated as a unit.						

PROTOCOL	This section occurs only in subsection 3R to indicate the protocol description file.
DESCRIPTION	This section defines the functionality and behavior of the service. Thus it describes concisely what the command does. It does not discuss OPTIONS or cite EXAMPLES. Interactive commands, subcommands, requests, macros, and functions are described under USAGE.
IOCTL	This section appears on pages in Section 7 only. Only the device class that supplies appropriate parameters to the ioctl(2) system call is called ioctl and generates its own heading. ioctl calls for a specific device are listed alphabetically (on the man page for that specific device). ioctl calls are used for a particular class of devices all of which have an io ending, such as mtio(7I).
OPTIONS	This secton lists the command options with a concise summary of what each option does. The options are listed literally and in the order they appear in the SYNOPSIS section. Possible arguments to options are discussed under the option, and where appropriate, default values are supplied.
OPERANDS	This section lists the command operands and describes how they affect the actions of the command.
OUTPUT	This section describes the output – standard output, standard error, or output files – generated by the command.
RETURN VALUES	If the man page documents functions that return values, this section lists these values and describes the conditions under which they are returned. If a function can return only constant values, such as 0 or -1 , these values are listed in tagged paragraphs. Otherwise, a single paragraph describes the return values of each function. Functions declared void do not return values, so they are not discussed in RETURN VALUES.
ERRORS	On failure, most functions place an error code in the global variable errno indicating why they failed. This section lists alphabetically all error codes a function can generate and describes the conditions that cause each error. When more than

	one condition can cause the same error, each condition is described in a separate paragraph under the error code.
USAGE	This section lists special rules, features, and commands that require in-depth explanations. The subsections listed here are used to explain built-in functionality:
	Commands Modifiers Variables Expressions Input Grammar
EXAMPLES	This section provides examples of usage or of how to use a command or function. Wherever possible a complete example including command-line entry and machine response is shown. Whenever an example is given, the prompt is shown as example%, or if the user must be superuser, example%, or if the user must be superuser, example#. Examples are followed by explanations, variable substitution rules, or returned values. Most examples illustrate concepts from the SYNOPSIS, DESCRIPTION, OPTIONS, and USAGE sections.
ENVIRONMENT VARIABLES	This section lists any environment variables that the command or function affects, followed by a brief description of the effect.
EXIT STATUS	This section lists the values the command returns to the calling program or shell and the conditions that cause these values to be returned. Usually, zero is returned for successful completion, and values other than zero for various error conditions.
FILES	This section lists all file names referred to by the man page, files of interest, and files created or required by commands. Each is followed by a descriptive summary or explanation.
ATTRIBUTES	This section lists characteristics of commands, utilities, and device drivers by defining the attribute type and its corresponding value. See attributes(5) for more information.
SEE ALSO	This section lists references to other man pages, in-house documentation, and outside publications.

10 man pages section 5: Standards, Environments, and Macros • December 2003

DIAGNOSTICS	This section lists diagnostic messages with a brief explanation of the condition causing the error.
WARNINGS	This section lists warnings about special conditions which could seriously affect your working conditions. This is not a list of diagnostics.
NOTES	This section lists additional information that does not belong anywhere else on the page. It takes the form of an aside to the user, covering points of special interest. Critical information is never covered here.
BUGS	This section describes known bugs and, wherever possible, suggests workarounds.

12 man pages section 5: Standards, Environments, and Macros • December 2003

Introduction

Intro(5)

NAME	Intro – introduction to miscellany				
DESCRIPTION	Among the topics presented in this section are:				
	Standards	The POSIX (IEEE) Standards and the X/Open Specifications are described on the standards page.			
	Environments	The user environment (environ), the subset of the user environment that depends on language and cultural conventions (locale), the large file compilation environment (lfcompile), and the transitional compilation environment (lfcompile64) are described.			
	Macros	The macros to format Reference Manual pages (man and mansun) as well as other text format macros (me, mm, and ms) are described.			
	Characters	Tables of character sets (ascii, charmap, eqnchar, and iconv), file format notation (formats), file name pattern matching (fnmatch), and regular expressions (regex and regexp) are presented.			
	FNS	Topics concerning the Federated Naming Service (fns, fns_initial_context, fns_policies, and fns_references) are discussed.			

Standards, Environments, and Macros

ascii(5)

NAME | ascii – map of ASCII character set

SYNOPSIS

cat /usr/pub/ascii

DESCRIPTION

/usr/pub/ascii is a map of the ASCII character set, to be printed as needed. It contains octal and hexadecimal values for each character. While not included in that file, a chart of decimal values is also shown here.

Octal - Character

000 NUL	001 SOH	002 STX	003 ETX	004 EOT	005 ENQ	006 ACK	007 BEL
010 BS	011 HT	012 NL	013 VT	014 NP	015 CR	016 SO	017 SI
020 DLE	021 DC1	022 DC2	023 DC3	024 DC4	025 NAK	026 SYN	027 ETB
030 CAN	031 EM	032 SUB	033 ESC	034 FS	035 GS	036 RS	037 US
040 SP	041 !	042 "	043 #	044 \$	045 %	046 &	047 ′
050 (051)	052 *	053 +	054 ,	055 -	056 .	057 /
060 0	061 1	062 2	063 3	064 4	065 5	066 6	067 7
070 8	071 9	072 :	073 ;	074 <	075 =	076 >	077 ?
100 @	101 A	102 B	103 C	104 D	105 E	106 F	107 G
110 H	111 I	112 J	113 K	114 L	115 M	116 N	117 O
120 P	121 Q	122 R	123 S	124 T	125 U	126 V	127 W
130 X	131 Y	132 Z	133 [134 \	135]	136 ^	137 _
140 `	141 a	142 b	143 c	144 d	145 e	146 f	147 g
150 h	151 i	152 ј	153 k	154 l	155 m	156 n	157 o
160 p	161 q	162 r	163 s	164 t	165 u	166 v	167 w
170 x	171 y	172 z	173 {	174	175 }	176 ~	177 DEL

Hexadecimal - Character

	00	NUL	01	SOH	02	STX	03	ETX	04	EOT	05	ENQ	06	ACK	07	BEL
	08	BS	09	HT	0A	NL	0B	VT	0 C	NP	0D	CR	0E	SO	0F	SI
	10	DLE	11	DC1	12	DC2	13	DC3	14	DC4	15	NAK	16	SYN	17	ETB
	18	CAN	19	EM	1A	SUB	1B	ESC	1C	FS	1D	GS	1E	RS	1F	US
	20	SP	21	!	22	"	23	#	24	\$	25	%	26	&	27	'
	28	(29)	2A	*	2B	+	2C	,	2D	-	2E		2F	/
	30	0	31	1	32	2	33	3	34	4	35	5	36	6	37	7
	38	8	39	9	3A	:	3B	;	3 C	<	3D	=	ЗE	>	ЗF	?
	40	@	41	A	42	В	43	С	44	D	45	Е	46	F	47	G
	48	Н	49	I	4A	J	4B	K	4C	L	4D	М	$4\mathrm{E}$	N	$4 \mathrm{F}$	0
	50	Ρ	51	Q	52	R	53	S	54	Т	55	U	56	V	57	W
	58	Х	59	Y	5A	Z	5B	[5C	\	5D]	5E	^	5F	_
	60	`	61	a	62	b	63	С	64	d	65	е	66	f	67	g
	68	h	69	i	6A	j	6B	k	6C	1	6D	m	6E	n	6F	0
	70	р	71	q	72	r	73	s	74	t	75	u	76	v	77	w
	78	x	79	У	7A	Z	7B	{	7C		7D	}	7E	~	7F	DEL
							-	- 1								
						Decin	na⊥	- Chai	cact	ler						
	,) NUL	1	SOH	2	STX	2	PTV	4	FOT	F	ENO	c	ACK	-	BEL
		B BS		SON HT		NL		VT		NP		CR		SO		SI
		5 DLE		DC1		DC2		DC3		DC4		NAK		SYN		ETB
		4 CAN		EM		SUB		ESC		FS		GS		RS		US
		2 SP	33		34		35		36			%	38			, ,
	4		41		42		43			ې ،		_	46			/
		30	49	,	50		4J 51		52		53		40 54			7
		58	49 57		58		59		52 60		53 61		54 62			?
		58 4@	57 65		58 66		59 67	,	60 68		69		62 70			G
		± @ 2 Н	65 73		66 74		67 75		68 76		69 77		78			0
I	/ .	2 H	13	T	/4	U	15	17	10	Ц	//	1*1	18	TN	19	0

16 man pages section 5: Standards, Environments, and Macros • Last Revised 19 Apr 2002

ascii(5)

80	Ρ	81	Q	82	R	83	S	84	Т	85	U	86	V	87	W
88	Х	89	Y	90	Ζ	91	[92	\setminus	93]	94	^	95	_
96	`	97	а	98	b	99	С	100	d	101	е	102	f	103	g
104	h	105	i	106	j	107	k	108	1	109	m	110	n	111	0
112	р	113	q	114	r	115	s	116	t	117	u	118	v	119	W
120	x	121	У	122	z	123	{	124		125	}	126	~	127	DEL

FILES

/usr/pub/ascii

On-line chart of octal and hexadecimal values for the ASCII character set.

Standards, Environments, and Macros 17

attributes, architecture, availability, CSI, stability, MT-Level - attributes of interfaces NAME

DESCRIPTION

The ATTRIBUTES section of a manual page contains a table (see below) defining attribute types and their corresponding values.

	ATTRIBUTE TYPE	ATTRIBUTE VALUE								
	Architecture	SPARC								
	Availability	SUNWcsu								
	CSI	Enabled								
	Interface Stability	Unstable								
	MT-Level	Safe								
Architecture	Architecture defines processor or specific hardware. (See -p option of uname(1)). In some cases, it may indicate required adapters or peripherals.									
Availability	This refers to the software package which contains the command or component being described on the man page. To be able to use the command, the indicated package must have been installed. For information on how to add a package see pkgadd(1M).									
Code Set Independence (CSI)	OS utilities and libraries which are free of dependencies on the properties of any code sets are said to have Code Set Independence (CSI). They have the attribute of being CSI enabled. This is in contrast to many commands and utilities in Solaris, for example, that work only with Extended Unix Codesets (EUC), an encoding method that allows concurrent support for up to four code sets and is commonly used to represent Asian character sets.									
	However, for practical reasons, this independence is not absolute. Certain assumptions are still applied to the current CSI implementation:									
	• File code is a superset of ASCII.									
	 To support multi-byte characters and n and / (slash) characters cannot be part 	ull-terminated UNIX file names, the NULL of any multi-byte characters.								
	 Only "stateless" file code encodings are supported. Stateless encoding avoids shift, locking shift, designation, invocation, and so forth, although single shift is not excluded. 									
	 Process code (wchar_t values) is implementation dependent and can change over time or between implementations or between locales. 									
		 Not every object in Solaris 2 and Solaris 7can have names composed of arbitrary characters. The names of the following objects must be composed of ASCII 								
	 User names, group name, and pass System name Names of printers and special device 									

	 Names of terminals (/dev/tty*) Process ID numbers Message queues, semaphores, and shared memory labels. The following may be composed of ISO Latin-1 or EUC characters: File names Directory names Command names Shell variables and environmental variable names Mount points for file systems NIS key names and domain names The names of NFS shared files should be composed of ASCII characters. Although
	files and directories may have names and contents composed of characters from non-ASCII code sets, using only the ASCII codeset allows NFS mounting across any machine, regardless of localization. For the commands and utilities that are CSI enabled, all can handle single-byte and multi-byte locales released in 2.6. For applications to get full support of internationalization services, dynamic binding has to be applied. Statically bound programs will only get support for C and POSIX locales.
Interface Stability	Sun often provides developers with early access to new technologies, which allows developers to evaluate with them as soon as possible. Unfortunately, new technologies are prone to changes and standardization often results in interface incompatibility from previous versions.
	To make reasonable risk assessments, developers need to know how likely an interface is to change in future releases. To aid developers in making these assessments, interface stability information is included on some manual pages for commands, entry-points, and file formats.
	The more stable interfaces can safely be used by nearly all applications, because Sun will endeavor to ensure that these continue to work in future minor releases. Applications that depend only on Standard and Stable interfaces should reliably continue to function correctly on future minor releases (but not necessarily on earlier major releases).
	The less stable interfaces allow experimentation and prototyping, but should be used only with the understanding that they might change incompatibly or even be dropped or replaced with alternatives in future minor releases.
	"Interfaces" that Sun does not document (for example, most kernel data structures and some symbols in system header files) may be implementation artifacts. Such internal interfaces are not only subject to incompatible change or removal, but we are unlikely to mention such a change in release notes.
Release Levels	Products are given release levels, as well as names, to aid compatibility discussions. Each release level may also include changes suitable for lower levels.

	Release	Version	Significance
	Major	x.0	Likely to contain major feature additions; adhere to different, possibly incompatible Standard revisions; and though unlikely, could change, drop, or replace Standard or Stable interfaces. Initial product releases are usually 1.0.
	Minor	x.y	Compared to an x.0 or earlier release (y!=0), it's likely to contain: minor feature additions, compatible Standard and Stable interfaces, possibly incompatible Evolving interfaces, or likely incompatible Unstable interfaces.
	Micro	x.y.z	Intended to be interface compatible with the previous release (z!=0), but likely to add bug fixes, performance enhancements, and support for additional hardware.
Classifications	level. The first colu Level for Incompa	umn lists the Stability Level. The table Changes, and the third co	el classifications relate to release e second column lists the Release lumn lists other comments. For a see the appropriate subsection below.
	Stability	Release	Comments
	Standard Stable	Major (x.0) Major (x.0)	Actual or de facto. Incompatibilities are exceptional.
	Evolving	Minor (x.y)	Migration advice might accompany an incompatibility.
	Unstable	Minor (x.y)	Experimental or transitional: incompatibilities are common.
	External	Micro (x.y.z)	Not controlled by Sun: intrarelease incompatibilities are common.

20 man pages section 5: Standards, Environments, and Macros • Last Revised 15 Aug 2001

Stability	Release	Comments
Obsolete	Minor (x.y)	Deprecated interface: likely to be removed in a future minor release.

The interface stability levels described in this manual page apply to both source and binary interfaces unless otherwise stated. The stability level of each interface is unknown unless explicitly stated.

Standard[: [organization_name,] standard_name, version]

The documented interface complies with the standard listed. If a standard is not specified the interface is defined by several standards. This is usually the hierarchy built up from the C Language (defined by ISO/IEC or K&R), SVID 3 and associated ABIs (defined by AT&T), the POSIX standards (defined by IEEE and ISO/IEC), and the Single UNIX Specifications (defined by The Open Group). See standards(5) for a complete list of these standards.

Most of these interfaces are defined by a formal standard, and controlled by a standards development organization. Changes will usually be made in accordance with approved changes to that standard. This stability level can also apply to interfaces that have been adopted (without a formal standard) by an "industry convention."

Support is provided for only the specified version(s) of a standard; support for later versions is not guaranteed. If the standards development organization approves a non-upward-compatible change to a Standard interface that Sun decides to support, Sun will announce a compatibility and migration strategy.

Programmers producing portable applications should rely on the interface descriptions present in the standard or specification to which the application is intended to conform, rather than the manual page descriptions of Standard interfaces. When the standard or specification allows alternative implementation choices, the manual page usually only describes the alternative implemented by Sun. The manual page also describes any compatible extensions to the base definition of Standard interfaces provided by Sun.

Stable

A Stable interface is a mature interface under Sun's control. Sun will try to avoid non-upwards-compatible changes to these interfaces, especially in minor or micro releases.

If support of a Stable interface must be discontinued, Sun will attempt to provide notification and the stability level changes to Obsolete.

Evolving

An Evolving interface may eventually become Standard or Stable but is still in transition.

	Sun will make reasonable efforts to ensure compatibility with previous releases as it evolves. When non-upwards compatible changes become necessary, they will occur in minor and major releases; such changes will be avoided in micro releases whenever possible. If such a change is necessary, it will be documented in the release notes for the affected release, and when feasible, Sun will provide migration aids for binary compatibility and continued source development.
	External An External interface is controlled by an entity other than Sun. At Sun's discretion, Sun can deliver as part of any release updated and possibly incompatible versions of such interfaces, subject to their availability from the controlling entity. This classification is typically applied to publicly available "freeware" and similar objects.
	For External interfaces, Sun makes no claims regarding either source or binary compatibility between any two releases. Applications based on these interfaces might not work in future releases, including patches that contain External interfaces.
	Unstable An Unstable interface is provided to give developers early access to new or rapidly changing technology or as an interim solution to a problem for which a more stable solution is anticipated in the future.
	For Unstable interfaces, Sun no claims about either source or binary compatibility from one minor release to another. Applications developed based on these interfaces may not work in future minor releases.
	Obsolete: Scheduled for removal after <i>event</i> An Obsolete interface is supported in the current release, but is scheduled to be removed in a future (minor) release. When support of an interface is to be discontinued, Sun will attempt to provide notification before discontinuing support. Use of an Obsolete interface may produce warning messages.
MT-Level	Libraries are classified into four categories which define their ability to support multiple threads. Manual pages containing routines that are of multiple or differing levels show this within their NOTES or USAGEsection.
	Safe Safe is an attribute of code that can be called from a multithreaded application. The effect of calling into a Safe interface or a safe code segment is that the results are valid even when called by multiple threads. Often overlooked is the fact that the result of this Safe interface or safe code segment can have global consequences that affect all threads. For example, the action of opening or closing a file from one thread is visible by all the threads within a process. A multi-threaded application has the responsibility for using these interfaces in a safe manner, which is different from whether or not the interface is Safe. For example, a multi-threaded application that closes a file that is still in use by other threads within the application is not using the close(2) interface safely.

Unsafe

An Unsafe library contains global and static data that is not protected. It is not safe to use unless the application arranges for only one thread at time to execute within the library. Unsafe libraries may contain routines that are Safe; however, most of the library's routines are unsafe to call.

The following table contains reentrant counterparts for Unsafe functions. This table is subject to change by Sun.

Reentrant functions for libc:

Unsafe Function	Reentrant counterpart
ctime	ctime_r
localtime	localtime_r
asctime	asctime_r
gmtime	gmtime_r
ctermid	ctermid_r
getlogin	getlogin_r
rand	rand_r
readdir	readdir_r
strtok	strtok_r
tmpnam	tmpnam_r

MT-Safe

An MT-Safe library is fully prepared for multithreaded access. It protects its global and static data with locks, and can provide a reasonable amount of concurrency. Note that a library can be safe to use, but not MT-Safe. For example, surrounding an entire library with a monitor makes the library Safe, but it supports no concurrency so it is not considered MT-Safe. An MT-Safe library must permit a reasonable amount of concurrency. (This definition's purpose is to give precision to what is meant when a library is described as Safe. The definition of a Safe library does not specify if the library supports concurrency. The MT-Safe definition makes it clear that the library is Safe, and supports some concurrency. This clarifies the Safe definition, which can mean anything from being single threaded to being any degree of multithreaded.)

Async-Signal-Safe

Async-Signal-Safe refers to particular library routines that can be safely called from a signal handler. A thread that is executing an Async-Signal-Safe routine will not deadlock with itself if interrupted by a signal. Signals are only a problem for MT-Safe routines that acquire locks.

Signals are disabled when locks are acquired in Async-Signal-Safe routines. This prevents a signal handler that might acquire the same lock from being called. The list of Async-Signal-Safe functions includes:

_exit	access	aio_error
aio_return	aio_suspend	alarm
cfgetispeed	cfgetospeed	cfsetispeed
cfsetospeed	chdir	chmod
chown	clock_gettime	close
creat	dup	dup2
execle	execve	fcntl
fdatasync	fork	fstat
fsync	getegid	geteuid
getgid	getgroups	getpgrp
getpid	getppid	getuid
kill	link	lseek
mkdir	mkfifo	open
pathconf	pause	pipe
read	rename	rmdir
sem_post	sema_post	setgid
setpgid	setsid	setuid
sigaction	sigaddset	sigdelset
sigemptyset	sigfillset	sigismember
sigpending	sigprocmask	sigqueue
sigsuspend	sleep	stat
sysconf	tcdrain	tcflow
tcflush	tcgetattr	tcgetpgrp
tcsendbreak	tcsetattr	tcsetpgrp
thr_kill	thr_sigsetmask	time
timer_getoverrun	timer_gettime	timer_settime
times	umask	uname

24 man pages section 5: Standards, Environments, and Macros • Last Revised 15 Aug 2001

unlink	utime	wait	
waitpid	write		

MT-Safe with Exceptions

See the NOTES or USAGE sections of these pages for a description of the exceptions.

Safe with Exceptions

See the NOTES or USAGE sections of these pages for a description of the exceptions.

Fork1-Safe

A Fork1-Safe library releases the locks it had held whenever fork1(2) is called in a Solaris thread program, or fork(2) in a POSIX (see standards(5)) thread program. Calling fork(2) in a POSIX thread program has the same semantic as calling fork1(2) in a Solaris thread program. All system calls, libpthread, and libthread are Fork1-Safe. Otherwise, you should handle the locking clean-up yourself (see pthread_atfork(3C)).

Cancel-Safety

If a multi-threaded application uses pthread cancel(3THR) to cancel (that is, kill) a thread, it is possible that the target thread is killed while holding a resource, such as a lock or allocated memory. If the thread has not installed the appropriate cancellation cleanup handlers to release the resources appropriately (see pthread cancel(3THR)), the application is "cancel-unsafe", that is, it is not safe with respect to cancellation. This unsafety could result in deadlocks due to locks not released by a thread that gets cancelled, or resource leaks; for example, memory not being freed on thread cancellation. All applications that use pthread cancel(3THR) should ensure that they operate in a Cancel-Safe environment. Libraries that have cancellation points and which acquire resources such as locks or allocate memory dynamically, also contribute to the cancel-unsafety of applications that are linked with these libraries. This introduces another level of safety for libraries in a multi-threaded program: Cancel-Safety. There are two sub-categories of Cancel-Safety: Deferred-Cancel-Safety, and Asynchronous-Cancel-Safety. An application is considered to be Deferred-Cancel-Safe when it is Cancel-Safe for threads whose cancellation type is PTHREAD CANCEL DEFERRED. An application is considered to be Asynchronous-Cancel-Safe when it is Cancel-Safe for threads whose cancellation type is PTHREAD CANCEL ASYNCHRONOUS. Deferred-Cancel-Safety is easier to achieve than Asynchronous-Cancel-Safety, since a thread with the deferred cancellation type can be cancelled only at well-defined cancellation points, whereas a thread with the asynchronous cancellation type can be cancelled anywhere. Since all threads are created by default to have the deferred cancellation type, it may never be necessary to worry about asynchronous cancel safety. Indeed, most applications and libraries are expected to always be Asynchronous-Cancel-Unsafe. An application which is Asynchronous-Cancel-Safe is also, by definition, Deferred-Cancel-Safe.

SEE ALSO uname(1), pkgadd(1M), Intro(3), standards(5)

charmap(5)

charmap(5)						
NAME	charmap – cl	naracter set de	escription file			
DESCRIPTION	set. Other int character set	formation abo	out the coded c ues are defined	haracter set n	nay also be in th	coded character ne file. Coded ames followed by
	The characte	r set descripti	on file provide	es:		
	character character	classes) indep s in the portal	pendent of cha	racter set enco et. This makes	such as collation oding, and usin it possible to c re the portable	g only the reate generic
					the portable cha gardless of enco	
Symbolic Names	(except for the characters consupported by encoding val	nose symbolic ommonly asso y the impleme ues are inclue	names that ar ciated with the entation, the sy ded in the file.	e shown with e symbolic name mbolic names Some of the e	identical glyph mes in the follo s and their corre ncodings assoc	esponding
	<ack></ack>	<dc2></dc2>	<enq></enq>	<fs></fs>	<is4></is4>	<soh></soh>
	<bel></bel>	<dc3></dc3>	<eot></eot>	<gs></gs>	<lf></lf>	<stx></stx>
	<bs></bs>	<dc4></dc4>	<esc></esc>	<ht></ht>	<nak></nak>	
	<can></can>		<etb></etb>	<is1></is1>	<rs></rs>	<syn></syn>
	<cr></cr>	<dle></dle>	<etx></etx>	<is2></is2>	<si></si>	<us></us>
	<dc1></dc1>		<ff></ff>	<is3></is3>	<so></so>	<vt></vt>
Declarations	the symbol s surrounding	hown in the f	ollowing list, s owed by one c	tarting in colu	lefinitions. Eacl 1mn 1, includin characters, follo	
	<code_set_na< th=""><th>me></th><th></th><th></th><th>character set fo n file is defined</th><th></th></code_set_na<>	me>			character set fo n file is defined	
	<mb_cur_ma:< th=""><th>r></th><th></th><th>mum number . This defaults</th><th>of bytes in a m to 1.</th><th>ulti-byte</th></mb_cur_ma:<>	r>		mum number . This defaults	of bytes in a m to 1.	ulti-byte
	<mb_cur_min< th=""><th>2></th><th>minimun</th><th></th><th>nteger value tha ytes in a charac</th><th></th></mb_cur_min<>	2>	minimun		nteger value tha ytes in a charac	

<escape_char></escape_char>	The escape character used to indicate that the characters following will be interpreted in a special way, as defined later in this section. This defaults to backslash (\thinsp;), which is the character glyph used in all the following text and examples, unless otherwise noted.
<comment_char></comment_char>	The character that when placed in column 1 of a charmap line, is used to indicate that the line is to be ignored. The default character is the number sign (#).
The shear of an estimation of	- Confidence - Still the still (the line of the set of the table della set in a set

Format The character set mapping definitions will be all the lines immediately following an identifier line containing the string CHARMAP starting in column 1, and preceding a trailer line containing the string END CHARMAP starting in column 1. Empty lines and lines containing a *<comment_char>* in the first column will be ignored. Each non-comment line of the character set mapping definition (that is, between the CHARMAP and END CHARMAP lines of the file) must be in either of two forms:

"%s %s %s\n", <symbolic-name>,<encoding>,<comments>

or

"%s. . .%s %s %s \n", <symbolic-name>,<symbolic-name>, <encoding>,<comments>

In the first format, the line in the character set mapping definition defines a single symbolic name and a corresponding encoding. A character following an escape character is interpreted as itself; for example, the sequence <\i\> represents the symbolic name \ enclosed between angle brackets.

In the second format, the line in the character set mapping definition defines a range of one or more symbolic names. In this form, the symbolic names must consist of zero or more non-numeric characters, followed by an integer formed by one or more decimal digits. The characters preceding the integer must be identical in the two symbolic names, and the integer formed by the digits in the second symbolic name must be equal to or greater than the integer formed by the digits in the first name. This is interpreted as a series of symbolic names formed from the common part and each of the integers between the first and the second integer, inclusive. As an example, <j0101>. . .<j0104> is interpreted as the symbolic names <j0101>, <j0102>, <j0103>, and <j0104>, in that order.

A character set mapping definition line must exist for all symbolic names and must define the coded character value that corresponds to the character glyph indicated in the table, or the coded character value that corresponds with the control character symbolic name. If the control characters commonly associated with the symbolic

Standards, Environments, and Macros 27

charmap(5)	
	names are supported by the implementation, the symbolic name and the corresponding encoding value must be included in the file. Additional unique symbolic names may be included. A coded character value can be represented by more than one symbolic name.
	The encoding part is expressed as one (for single-byte character values) or more concatenated decimal, octal or hexadecimal constants in the following formats:
	"%cd%d" , <escape_char>,<decimal byte="" value=""> "%cx%x" , <escape_char>,<hexadecimal byte="" value=""> "%c%0" , <escape_char>,<octal byte="" value=""></octal></escape_char></hexadecimal></escape_char></decimal></escape_char>
Decimal Constants	Decimal constants must be represented by two or three decimal digits, preceded by the escape character and the lower-case letter d; for example, $\05$, $\07$, or $\143$. Hexadecimal constants must be represented by two hexadecimal digits, preceded by the escape character and the lower-case letter x; for example, $\x05$, $\x61$, or $\x8f$. Octal constants must be represented by two or three octal digits, preceded by the escape character; for example, $\05$, $\141$, or $\217$. In a portable charmap file, each constant must represent an 8-bit byte. Implementations supporting other byte sizes may allow constants to represent values larger than those that can be represented in 8-bit bytes, and to allow additional digits in constants. When constants are concatenated for multi-byte character values, they must be of the same type, and interpreted in byte order from first to last with the least significant byte of the multi-byte character specified by the last constant.
Ranges of Symbolic Names	In lines defining ranges of symbolic names, the encoded value is the value for the first symbolic name in the range (the symbolic name preceding the ellipsis). Subsequent symbolic names defined by the range will have encoding values in increasing order. For example, the line
	<j0101><j0104> \d129\d254</j0104></j0101>
	will be interpreted as:
	<j0101> \d129\d254 <j0102> \d129\d255 <j0103> \d130\d0 <j0104> \d130\d1</j0104></j0103></j0102></j0101>
	Note that this line will be interpreted as the example even on systems with bytes larger than 8 bits. The comment is optional.
SEE ALSO	<pre>locale(1) localedef(1) nl_langinfo(3C) extensions(5), locale(5)</pre>

crypt_bsdbf(5)

NAME	crypt_bsdbf – password hashing module us	sing Blowfish cryptographic algorithm	(-)
SYNOPSIS	/usr/lib/security/\$ISA/crypt_bsdbf.so		
DESCRIPTION	The crypt_bsdmd5 module is a one-way p crypt(3C) that uses the Blowfish cryptogra for crypt.conf(4) and policy.conf(4)	aphic algorithm. The algorithm identifier	1
ATTRIBUTES	See attributes(5) for descriptions of the	following attributes:	
	ATTRIBUTE TYPE	ATTRIBUTE VALUE	
	MT-Level	Safe	
SEE ALSO	<pre>passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5)</pre>		
		Standards, Environments, and Macros	29

crypt_bsdmd5(5)

NAME	crypt_bsdmd5 – password hashing module	e using MD5 message hash algorithm
SYNOPSIS	/usr/lib/security/\$ISA/crypt_bsdmd5.sc	
DESCRIPTION	The crypt_bsdmd5 module is a one-way crypt(3C) that uses the MD5 message has crypt.conf(4) and policy.conf(4) is 1 on BSD and Linux systems. See attributes(5) for descriptions of the	h algorithm. The algorithm identifier for . The output is compatible with md5crypt
ATTRIDCT LO		following attributes.
	ATTRIBUTE TYPE	ATTRIBUTE VALUE
	MT-Level	Safe
SEE ALSO	<pre>passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5)</pre>	<pre>impl(3C), crypt_gensalt(3C), case(3C), crypt.conf(4), passwd(4),</pre>

crypt_sunmd5(5)

	crypt_sunmd5 - password hashing module	using MD5 message hash algorithm
SYNOPSIS	/usr/lib/security/\$ISA/crypt_sunmd5.so	
DESCRIPTION	The crypt_sunmd5 module is a one-way p crypt(3C) that uses the MD5 message hash crypt.conf(4) and policy.conf(4) is me	n algorithm. The algorithm identifier for
	This module is designed to make it difficult attacks based on high speed MD5 impleme loops, and table lookup.	
ATTRIBUTES	See attributes(5) for descriptions of the	following attributes:
	ATTRIBUTE TYPE	ATTRIBUTE VALUE
	MT-Level	Safe
SEE ALSO	<pre>passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5)</pre>	

crypt_unix(5)		
NAME	crypt_unix – traditional UNIX crypt algorit	hm
DESCRIPTION	The crypt_unix algorithm is the tradition considered sufficiently secure for current sy compatibility. The crypt_sunmd5(5), cryp algorithm should be used instead.	stems and is provided for backwards
	The algorithm identifier for policy.conf crypt.conf(4) for this algorithm.	(4) isunix There is no entry in
	The crypt_unix algorithm is internal to 1 function used by crypt(3C) when the first	1 0 0
	This algorithm is based on a one-way encry (among other things) to frustrate use of har Only the first eight characters of the key pa algorithm; the rest are silently ignored. The the set [a-zA-Z0-9./]. This string is used to 4096 different ways.	dware implementations of a key search. ssed to crypt() are used with this salt is a two-character string chosen from
USAGE	The return value of the crypt_unix algoristandard-conforming systems. See standa	
ATTRIBUTES	See attributes(5) for descriptions of the	following attributes:
	ATTRIBUTE TYPE	ATTRIBUTE VALUE
	MT-Level	Safe
SEE ALSO	MT-Level passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5), crypt_ crypt_sunmd5(5), standards(5)	Safe impl(3C), crypt_gensalt(3C), ase(3C), crypt.conf(4), passwd(4),
SEE ALSO	passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5), crypt_	Safe impl(3C), crypt_gensalt(3C), ase(3C), crypt.conf(4), passwd(4),
SEE ALSO	passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5), crypt_	Safe impl(3C), crypt_gensalt(3C), ase(3C), crypt.conf(4), passwd(4),
SEE ALSO	passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5), crypt_	Safe impl(3C), crypt_gensalt(3C), ase(3C), crypt.conf(4), passwd(4),
SEE ALSO	passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5), crypt_	Safe impl(3C), crypt_gensalt(3C), ase(3C), crypt.conf(4), passwd(4),
SEE ALSO	passwd(1), crypt(3C), crypt_genhash_ crypt_gensalt_impl(3C), getpassphr policy.conf(4), attributes(5), crypt_	Safe impl(3C), crypt_gensalt(3C), ase(3C), crypt.conf(4), passwd(4),

32 man pages section 5: Standards, Environments, and Macros • Last Revised 10 Jun 2002

NAME	dhcp – Dynamic Host Configuration Protocol		
DESCRIPTION	Dynamic Host Configuration Protocol (DHCP) enables host systems in a TCP/IP network to be configured automatically for the network as they boot. DHCP uses a client/server mechanism: servers store configuration information for clients, and provide that information upon a client's request. The information can include the client's IP address and information about network services available to the client.		
	This manual page provides a brief summary of the Solaris DHCP implementation.		
Solaris DHCP Client	The Solaris DHCP client is implemented as background daemon, dhcpagent(1M). This daemon is started automatically during bootup if there exists at least one dhcp. <i>interface</i> file in /etc. Only interfaces with a corresponding /etc/dhcp. <i>interface</i> file are automatically configured during boot. Network parameters needed for system configuration during bootup are extracted from the information recieved by the daemon through the use of the dhcpinfo(1) command. The daemon's default behavior can be altered by changing the tunables in the /etc/default/dhcpagent file. The daemon is controlled by the ifconfig(1M) utility. Check the status of the daemon using the netstat(1M) and ifconfig(1M) commands.		
Solaris DHCP Server	The Solaris DHCP server is implemented as a background daemon, in.dhcpd(1M). This daemon can deliver network configuration information to either BOOTP or DHCP clients. The Solaris DHCP service can be managed using the dhcpmgr(1M) GUI or the command line utilities dhcpconfig(1M), dhtadm(1M), and pntadm(1M).		
DHCP Configuration Tables	The Solaris DHCP server stores client configuration information in the following two types of tables:		
	dhcptab tables	Contain macros and options (also known as symbols), used to construct a package of configuration information to send to each DHCP client. There exists only one dhcptab for the DHCP service. The dhcptab(4) can be viewed and modified using the dhtadm(1M) command or dhcpmgr(1M) graphical utility. See dhcptab(4) for more information about the syntax of dhcptab records. See dhcp_inittab(4) for more information about the DHCP options and symbols.	
	DHCP network tables	DHCP network tables, which contain mappings of client IDs to IP addresses and parameters associated with those addresses. Network tables are named with the IP address of the network, and can be created, viewed, and modified using the pntadm command or dhcpmgr graphical utility. See dhcp_network(4) for more information about network tables.	

dhcp(5)

Solaris DHCP Service Developer's Guide

Alexander, S., and R. Droms. *RFC 2132, DHCP Options and BOOTP Vendor Extensions*. Silicon Graphics, Inc. Bucknell University. March 1997.

Droms, R. *RFC 1534, Interoperation Between DHCP and BOOTP*. Bucknell University. October 1993.

Droms, R. *RFC 2131, Dynamic Host Configuration Protocol*. Bucknell University. March 1997.

Wimer, W.RFC 1542, Clarifications and Extensions for the Bootstrap Protocol. Carnegie Mellon University. October 1993.

NAME	dhcp_modules – data storage modules for the DHCP service	
DESCRIPTION	This man page describes the characteristics of data storage modules (public modules for use by the Solaris Dynamic Host Configuration Protocol (DHCP) service.	
	Public modules are the part of the DHCP service architecture that encapsulate the details of storing DHCP service data in a data storage service. Examples of data storage services are NIS+, Oracle, and ufs file systems.	
	Public modules are dynamic objects which can be shipped separately from the Solaris DHCP service. Once installed, a public module is visible to the DHCP service, and can be selected for use by the service through the DHCP service management interfaces (dhcpmgr(1M), dhcpconfig(1M), dhtadm(1M), and pntadm(1M)).	
	Public modules may be provided by Sun Microsystems, Inc or by third parties.	
	The Solaris DHCP service management architecture provides a mechanism for plugging in public module-specific administration functionality into the dhcpmgr(1M) and dhcpconfig(1M) utilities. This functionality is in the form of a Java Bean, which is provided by the public module vendor. This Java Bean collects public module-specific configuration from the user (you) and provides it to the Solaris DHCP service.	
	The Solaris DHCP service bundles three modules with the service, which are described below. There are three dhcpsvc.conf(4) DHCP service configuration parameters pertaining to public modules: RESOURCE, PATH, and RESOURCE_CONFIG. See dhcpsvc.conf(4) for more information about these parameters.	
SUNWfiles	This module stores its data in ASCII files. Although the format is ASCII, hand-editing is discouraged. It is useful for DHCP service environments that support several hundred to a couple thousand of clients and lease times are a few hours or more.	
	This module's data may be shared between DHCP servers through the use of NFS.	
SUNWbinfiles	This module stores its data in binary files. It is useful for DHCP service environments with many networks and many thousands of clients. This module provides an order of magnitude increase in performance and capacity over SUNWfiles.	
	This module's data cannot be shared between DHCP servers.	
SUNWnisplus	This module stores its data within a NIS+ domain. It is useful in environments where NIS+ is already deployed and facilitates sharing among multiple DHCP servers. This module suports several hundred to a few thousand clients with lease times of several hours or more.	
	The NIS+ service should be hosted on a machine with ample CPU power, memory, and disk space, as the load on NIS+ is significant when it is used to store DHCP data. Periodic checkpointing of the NIS+ service is necessary in order to roll the transaction logs and keep the NIS+ service operating at its highest efficiency. See nisping(1M) and crontab(1) for more information.	

dhcp_modules(5)

SEE ALSO crontab(1), dhcpconfig(1M), dhcpmgr(1M), dhtadm(1M), nisping(1M),pntadm(1M), dhcpsvc.conf(4), dhcp(5)Solaris DHCP Service Developer's Guide

NAME | environ – user environment

DESCRIPTION

When a process begins execution, one of the exec family of functions makes available an array of strings called the environment; see exec(2). By convention, these strings have the form *variable=value*, for example, PATH=/sbin:/usr/sbin. These environmental variables provide a way to make information about a program's environment available to programs.

A name may be placed in the environment by the export command and *name=value* arguments in sh(1), or by one of the exec functions. It is unwise to conflict with certain shell variables such as MAIL, PS1, PS2, and IFS that are frequently exported by .profile files; see profile(4).

The following environmental variables can be used by applications and are expected to be set in the target run-time environment.

HOME

The name of the user's login directory, set by login(1) from the password file; see passwd(4).

LANG

The string used to specify internationalization information that allows users to work with different national conventions. The setlocale(3C) function checks the LANG environment variable when it is called with "" as the locale argument. LANG is used as the default locale if the corresponding environment variable for a particular category is unset or null. If, however, LC_ALL is set to a valid, non-empty value, its contents are used to override both the LANG and the other LC_* variables. For example, when invoked as setlocale(LC_CTYPE, ""), setlocale() will query the LC_CTYPE environment variable first to see if it is set and non-null. If LC_CTYPE is not set or null, then setlocale() will check the LANG environment variable to see if it is set and non-null. If both LANG and LC_CTYPE are unset or NULL, the default "C" locale will be used to set the LC_CTYPE category.

Most commands will invoke setlocale (LC_ALL, "") prior to any other processing. This allows the command to be used with different national conventions by setting the appropriate environment variables.

The following environment variables correspond to each category of setlocale(3C):

LC ALL

If set to a valid, non-empty string value, override the values of LANG and all the other LC_*variables.

LC_COLLATE

This category specifies the character collation sequence being used. The information corresponding to this category is stored in a database created by the localedef(1) command. This environment variable affects strcoll(3C) and strxfrm(3C).

environ(5)

LC CTYPE

This category specifies character classification, character conversion, and widths of multibyte characters. When LC_CTYPE is set to a valid value, the calling utility can display and handle text and file names containing valid characters for that locale; Extended Unix Code (EUC) characters where any individual character can be 1, 2, or 3 bytes wide; and EUC characters of 1, 2, or 3 column widths. The default "C" locale corresponds to the 7-bit ASCII character set; only characters from ISO 8859-1 are valid. The information corresponding to this category is stored in a database created by the localedef() command. This environment variable is used by ctype(3C), mblen(3C), and many commands, such as cat(1), ed(1), ls(1), and vi(1).

LC_MESSAGES

This category specifies the language of the message database being used. For example, an application may have one message database with French messages, and another database with German messages. Message databases are created by the mkmsgs(1) command. This environment variable is used by exstr(1), gettxt(1), srchtxt(1), gettxt(3C), and gettext(3C).

LC MONETARY

This category specifies the monetary symbols and delimiters used for a particular locale. The information corresponding to this category is stored in a database created by the localedef(1) command. This environment variable is used by localeconv(3C).

LC NUMERIC

This category specifies the decimal and thousands delimiters. The information corresponding to this category is stored in a database created by the localedef() command. The default C locale corresponds to "." as the decimal delimiter and no thousands delimiter. This environment variable is used by localeconv(3C), printf(3C), and strtod(3C).

LC TIME

This category specifies date and time formats. The information corresponding to this category is stored in a database specified in localedef(). The default C locale corresponds to U.S. date and time formats. This environment variable is used by many commands and functions; for example: at(1), calendar(1), date(1), strftime(3C), and getdate(3C).

MSGVERB

Controls which standard format message components fmtmsg selects when messages are displayed to stderr; see fmtmsg(1) and fmtmsg(3C).

NETPATH

A colon-separated list of network identifiers. A network identifier is a character string used by the Network Selection component of the system to provide application-specific default network search paths. A network identifier must consist of non-null characters and must have a length of at least 1. No maximum length is specified. Network identifiers are normally chosen by the system administrator. A network identifier is also the first field in any /etc/netconfig file entry. NETPATH thus provides a link into the /etc/netconfig file and the information about a network contained in that network's entry. /etc/netconfig is maintained by the system administrator. The library routines described in getnetpath(3NSL) access the NETPATH environment variable.

NLSPATH

Contains a sequence of templates which catopen(3C) and gettext(3C) use when attempting to locate message catalogs. Each template consists of an optional prefix, one or more substitution fields, a filename and an optional suffix. For example:

NLSPATH="/system/nlslib/%N.cat"

defines that catopen() should look for all message catalogs in the directory /system/nlslib, where the catalog name should be constructed from the *name* parameter passed to catopen(), %N, with the suffix .cat.

Substitution fields consist of a % symbol, followed by a single-letter keyword. The following keywords are currently defined:

- %N The value of the *name* parameter passed to catopen().
- %L The value of LANG or LC MESSAGES.
- %l The language element from LANG or LC_MESSAGES.
- %t The territory element from LANG or LC MESSAGES.
- %c The codeset element from LANG or LC MESSAGES.
- %% A single % character.

An empty string is substituted if the specified value is not currently defined. The separators "" and "." are not included in %t and %c substitutions.

Templates defined in NLSPATH are separated by colons (:). A leading colon or two adjacent colons (::) is equivalent to specifying %N. For example:

NLSPATH=":%N.cat:/nlslib/%L/%N.cat"

indicates to catopen() that it should look for the requested message catalog in name, name.cat and /nlslib/\$LANG/name.cat. For gettext(), %N automatically maps to "messages".

If NLSPATH is unset or NULL, catopen() and gettext() call setlocale(3C), which checks LANG and the LC * variables to locate the message catalogs.

NLSPATH will normally be set up on a system wide basis (in /etc/profile) and thus makes the location and naming conventions associated with message catalogs transparent to both programs and users.

PATH

The sequence of directory prefixes that sh(1), time(1), nice(1), nohup(1), and other utilities apply in searching for a file known by an incomplete path name. The prefixes are separated by colons (:). login(1) sets PATH=/usr/bin. For more detail, see sh(1).

Standards, Environments, and Macros 39

environ(5)

SEV LEVEL

Define severity levels and associate and print strings with them in standard format error messages; see addseverity(3C), fmtmsg(1), and fmtmsg(3C).

TERM

The kind of terminal for which output is to be prepared. This information is used by commands, such as vi(1), which may exploit special capabilities of that terminal.

ΤZ

Timezone information. The contents of this environment variable are used by the functions ctime(3C), localtime(3C), strftime(3C), and mktime(3C) to override the default timezone. If TZ is not in the following form, it designates a path to a timezone database file relative to /usr/share/lib/zoneinfo/, ignoring the first character if it is a colon (:). Otherwise, TZ has the form:

stdoffset [dst [offset] [, start [/time], end [/time]]]

std and dst	Three or more bytes that are the designation for the standard (<i>std</i>) and daylight savings time (<i>dst</i>) timezones. Only <i>std</i> is required. If <i>dst</i> is missing, then daylight savings time does not apply in this locale. Upper- and lower-case letters from the portable character set are explicitly allowed. Any graphic characters from the portable character set except a leading colon (:) or digits, the comma (,), the minus (-), the plus (+), and the null character are permitted to appear in these fields, but their meaning is unspecified.
offset	Indicates the value one must add to the local time to arrive at Coordinated Universal Time. The offset has the form:
	<i>hh</i> [: <i>mm</i> [: <i>ss</i>]]
	The minutes (<i>mm</i>) and seconds (<i>ss</i>) are optional. The hour (<i>hh</i>) is required and may be a single digit. The <i>offset</i> following <i>std</i> is required. If no <i>offset</i> follows <i>dst</i> , daylight savings time is assumed to be one hour ahead of standard time. One or more digits may be used. The value is always interpreted as a decimal number. The hour must be between 0 and 24, and the minutes (and seconds), if present, must be between 0 and 59. Out of range values may cause unpredictable behavior. If preceded by a "–", the timezone is east of the Prime Meridian. Otherwise, it is west of the Prime Meridian (which may be indicated by an optional preceding "+" sign).

environ(5)

	start/time, end/time	savings tin change fro occurs, an happens. I	when to change to and back from daylight ne, where <i>start/time</i> describes when the om standard time to daylight savings time d <i>end/time</i> describes when the change back Each time field describes when, in current the change is made.
		The forma	ts of <i>start</i> and <i>end</i> are one of the following:
		Jn	The Julian day n ($1 \le n \le 365$). Leap days are not counted. That is, in all years, February 28 is day 59 and March 1 is day 60. It is impossible to refer to the occasional February 29.
		п	The zero-based Julian day ($0 \le n \le 365$). Leap days are counted, and it is possible to refer to February 29.
		Mm.n.d	The d^{th} day, $(0 \le d \le 6)$ of week <i>n</i> of month <i>m</i> of the year $(1 \le n \le 5, 1 \le m \le 12)$, where week 5 means "the last <i>d</i> -day in month <i>m</i> " which may occur in either the fourth or the fifth week). Week 1 is the first week in which the d^{th} day occurs. Day zero is Sunday.
			tation specific defaults are used for <i>start</i> these optional fields are not given.
		leading sig	has the same format as <i>offset</i> except that no gn ("—" or "+" is allowed. The default, if ot given is 02:00:00.
SEE ALSO	<pre>nice(1), nohup(1), sh(1), sort(1) catopen(3C), ctime(3C), ctype getnetpath(3NSL), gettext(3 mktime(3C), printf(3C), setle</pre>	l), time(1), e(3C), fmtr (C), gettxt ocale(3C),	t(3C), localeconv(3C), mblen(3C),

Standards, Environments, and Macros 41

eqnchar	(5)
---------	-----

[incluar(0)							
NAME	eqnchar –	special characte	er definitions	s for e	eqn		
SYNOPSIS	eqn /usr/share/lib/pub/eqnchar filename troff options						
		r/share/lib/pu	-				
DESCRIPTION	construction definitions		at are not av ntended for	ailab use w	le on the Gr	aphic Syste	er definitions for oms typesetter. These). It contains
	ciplus	\oplus		Ш		square	
	citimes	\otimes	langle	ì		circle	0
	wig	~	rangle	ħ		blot	
	-wig	~	hbar mrd			bullet	•
	>wig <wig< th=""><th>> ~ ~ ≅</th><th>ppd <-></th><th>$\stackrel{\perp}{\leftrightarrow}$</th><th></th><th>prop empty</th><th>∝ Ø</th></wig<>	> ~ ~ ≅	ppd <->	$\stackrel{\perp}{\leftrightarrow}$		prop empty	∝ Ø
	=wig	2 ~II	<=>	\Leftrightarrow		member	e
	star	*	<	≮		потет	¢
	bigstar	*	>	≯		сир	U
	=dot	÷	ang			сар	\cap
	orsign	X	rang	Ŀ		incl	
	andsign	X	3dot	Ξ.		subset	
	=del oppA	$ \begin{array}{c} \Delta \\ = \\ - \chi \\ \exists \\ \mathring{A} \end{array} $	thf quarter	· · · 1/4		supset !subset	\supset
	oppE		3quarter	3/4		!supset	
	angstrom	Å	degree	o			=
FILES	/usr/sha	are/lib/pub/	eqnchar				
ATTRIBUTES	See attri	butes(5) for d	escriptions o	of the	following a	ttributes:	
		ATTRIBUTE T	YPE			ATTRIBUT	E VALUE
	Availability	У			SUNWdoc		
SEE ALSO	eqn(1), nr	off(1),troff((1),attribu	ites(5)		

42 man pages section 5: Standards, Environments, and Macros • Last Revised 12 Jul 2002

extensions(5)

NAME	extensions – localedef extensions description file
DESCRIPTION	A localedef extensions description file or <i>extensions</i> file defines various extensions for the localedef(1) command.
	The localedef extensions description file provides:
	EUC code set width information via the cswidth keyword: cswidth bc1 : sw1, bc2 : sw2, bc3 : sw3where bc1, bc2, and bc3 indicate the number of bytes (byte count) per character for EUC codesets 1, 2, and 3, respectively. sw1, sw2, and sw3 indicate screen width for EUC codesets 1, 2, and 3, respectively.
	 Other extensions which will be documented in a future release.
SEE ALSO	<pre>locale(1), localedef(1), environ(5), locale(5)</pre>

NAME | filesystem – file system organization

SYNOPSIS

/usr

DESCRIPTION

The file system tree is organized for administrative convenience. Distinct areas within the file system tree are provided for files that are private to one machine, files that can be shared by multiple machines of a common platform, files that can be shared by all machines, and home directories. This organization allows sharable files to be stored on one machine but accessed by many machines using a remote file access mechanism such as NFS. Grouping together similar files makes the file system tree easier to upgrade and manage.

The file system tree consists of a root file system and a collection of mountable file systems. The mount(2) program attaches mountable file systems to the file system tree at mount points (directory entries) in the root file system or other previously mounted file systems. Two file systems, / (the root) and /usr, must be mounted in order to have a completely functional system. The root file system is mounted automatically by the kernel at boot time; the /usr file system is mounted by the system start-up script, which is run as part of the booting process.

Certain locations, noted below, are approved installation locations for bundled Foundation Solaris software. In some cases, the approved locations for bundled software are also approved locations for add-on system software or for applications. The following descriptions make clear where the two locations differ. For example, /etc is the installation location for platform‐dependent configuration files that are bundled with Solaris software. The analogous location for applications is /etc/opt/*packagename*.

In the following descriptions, *subsystem* is a category of application or system software, such as a window system (dt) or a language (java1.2)

The following descriptions make use of the terms *platform, platform‐dependent, platform‐independent,* and *platform‐specific.* Platform*‐independent,* and *platform‐specific.* Platform refers to a machines Instruction Set Architecture or processor type, such as is returned by uname -i. *Platform‐dependent* refers to a file that is installed on all platforms and whose contents vary depending on the platform. Like a platform‐dependent file, a *platform‐independent* file is installed on all platforms. However, the contents of the latter type remains the same on all platforms. An example of a platform‐independent file is a standard configuration file, such as /etc/hosts. Unlike a platform‐dependent or a platform‐independent file, the *platform‐specific* file is installed only on a subset of supported platforms. Most platform-specific files are gathered under /platform and /usr/platform.

Root File System The root file system contains files that are unique to each machine. It contains the following directories:

44 man pages section 5: Standards, Environments, and Macros • Last Revised 15 Nov 2001

Root of the overall file system name space.

/dev

Primary location for special files. Typically, device files are built to match the kernel and hardware configuration of the machine.

/dev/cfg

Symbolic links to physical ap_ids.

/dev/cua

Device files for uucp.

/dev/dsk

Block disk devices.

/dev/fbs

Frame buffer device files.

/dev/fd

File descriptors.

/dev/md

Logical volume management meta-disk devices.

/dev/printers USB printer device files.

/dev/pts Pseudo-terminal devices.

/dev/rdsk Raw disk devices.

/dev/rmt

Raw tape devices.

/dev/sad

Entry points for the STREAMS Administrative driver.

/dev/sound

Audio device and audio device control files.

/dev/swap

Default swap device.

/dev/term

Terminal devices.

/devices

Physical device files.

/etc

Platform‐dependent administrative and configuration files and databases that are not shared among systems. /etc may be viewed as the directory that

defines the machine's identity. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /etc/opt/*packagename*.

/etc/acct

Accounting system configuration information.

/etc/apache

Apache configuration files.

/etc/cron.d

Configuration information for cron(1M).

/etc/default

Defaults information for various programs.

/etc/dfs

Configuration information for shared file systems.

/etc/dhcp

Dynamic Host Configuration Protocol (DHCP) configuration files.

/etc/dmi

Solstice Enterprise Agents configuration files.

/etc/fn

Federated Naming Service and X.500 support files.

/etc/fs

Binaries organized by file system types for operations required before /usr is mounted.

/etc/gss

Generic Security Service (GSS) Application Program Interface configuration files.

/etc/gtk

GNOME (GNU Network Object Model Environment) configuration files.

/etc/inet

Configuration files for Internet services.

/etc/init.d

Shell scripts for transitioning between run levels.

/etc/iplanet

iPlanet configuration files.

/etc/krb5

Kerberos configuration files.

/etc/lib

Shared libraries needed during booting.

/etc/lp

Configuration information for the printer subsystem.

```
/etc/llc2
  Logical link control (11c2) driver configuration files.
/etc/lp
  Configuration information for the printer subsystem.
/etc/lu
  Solaris Live Upgrade configuration files.
/etc/lvm
  Solaris Logical Volume Manager configuration files.
/etc/mail
  Mail subsystem configuration.
/etc/nca
  Solaris Network Cache and Accelerator (NCA) configuration files.
/etc/net
  Configuration information for transport independent network services.
/etc/nfs
  NFS server logging configuration file.
/etc/openwin
  OpenWindows configuration files.
/etc/opt
  Configuration information for optional packages.
/etc/ppp
  Solaris PPP configuration files.
/etc/rc0.d
  Scripts for entering or leaving run level 0. See init(1M).
/etc/rc1.d
  Scripts for entering or leaving run level 1. See init(1M).
/etc/rc2.d
  Scripts for entering or leaving run level 2. See init(1M).
/etc/rc3.d
  Scripts for entering or leaving run level 3. See init(1M).
/etc/rcS.d
  Scripts for bringing the system up in single user mode.
/etc/rcm
  Directory for reconfiguration manager (RCM) custom scripts.
/etc/rpcsec
  This directory might contain an NIS+ authentication configuration file.
/etc/saf
  Service Access Facility files.
```

| /etc/security

Basic Security Module (BSM) configuration files.

/etc/sfw

Samba configuration files.

/etc/skel

Default profile scripts for new user accounts. See useradd(1M).

/etc/smartcard

Solaris SmartCard configuration files.

/etc/snmp

Solstice Enterprise Agents configuration files.

/etc/ssh

Secure Shell configuration files. See ssh(1)

/etc/sysevent

syseventd configuration files.

/etc/subsystem

Platform‐dependent *subsystem* configuration files that are not shared among systems. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /etc/opt/*packagename*.

/etc/tm

Trademark files; contents displayed at boot time.

/etc/usb

USB configuration information.

/etc/uucp

UUCP configuration information. See uucp(1C).

/etc/wrsm

WCI Remote Shared Memory (WRSM) configuration information. See wrsmconf(1M)

/export

Default root of the shared file system tree.

/home

Default root of a subtree for user directories.

/kernel

Subtree of platform‐dependent loadable kernel modules required as part of the boot process. It includes the generic part of the core kernel that is platform‐independent, /kernel/genunix. See kernel(1M) An approved installation location for bundled Solaris software and for add-on system software.

/kernel/drv 32-bit device drivers.

48 man pages section 5: Standards, Environments, and Macros • Last Revised 15 Nov 2001

/kernel/drv/sparcv9

64-bit SPARC device drivers.

/kernel/genunix

Platform‐independent kernel.

/kernel/subsystem/ia64

64-bit x86 platform‐dependent modules required for boot. An approved installation location for bundled Solaris software and for add-on system software. Note that ia64 is an example name; the actual name might be different.

/kernel/subsystem/sparcv9

64-bit SPARC platform‐dependent modules required for boot. An approved installation location for bundled Solaris software and for add-on system software.

/mnt

Default temporary mount point for file systems. This is an empty directory on which file systems can be temporarily mounted.

/opt

Root of a subtree for add-on application packages.

/platform

Subtree of platform‐specific objects which need to reside on the root filesystem. It contains a series of directories, one per supported platform. The semantics of the series of directories is equivalent to / (root).

/platform/`uname -i`/kernel

Platform‐specific modules required for boot. These modules have semantics equivalent to /kernel. It includes the file unix, the core kernel. See kernel(1M). An approved installation location for bundled Solaris software and for add-on system software.

/platform/`uname -m`/kernel

Hardware class-specific modules required for boot. An approved installation location for bundled Solaris software and for add-on system software.

/platform/`uname -i`/kernel/subsystem/ia64

x86 64-bit, platform‐dependent modules required for boot. Note that ia64 is an example name; the actual name might be different. An approved installation location for bundled Solaris software.

/platform/`uname -i`/kernel/subsystem/sparcv9 SPARC 64-bit platform‐specific modules required for boot. An approved installation location for bundled Solaris software.

/platform/`uname -i`/kernel/sparcv9/unix
 64-bit platform‐dependent kernel.

/platform/`uname -i`/kernel/unix
32-bit platform‐dependent kernel.

/platform/`uname -i`/lib

Platform‐specific shared objects required for boot. Semantics are equivalent to /lib. An approved installation location for bundled Solaris software and for add-on system software.

/platform/`uname -i`/sbin

Platform‐specific administrative utilities required for boot. Semantics are equivalent to /sbin. An approved installation location for bundled Solaris software and for add-on system software.

/proc

Root of a subtree for the process file system.

/sbin

Essential executables used in the booting process and in manual system recovery. The full complement of utilities is available only after /usr is mounted. /sbin is an approved installation location for bundled Solaris software.

/tmp

Temporary files; cleared during the boot operation.

/usr

Mount point for the /usr file system. See description of /usr file system, below.

/var

Root of a subtree for varying files. Varying files are files that are unique to a machine but that can grow to an arbitrary (that is, variable) size. An example is a log file. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /var/opt/*packagename*.

/var/adm

System logging and accounting files.

/var/apache

Scripts, icons, logs, and cache pages for Apache web server.

/var/audit

Basic Security Module (BSM) audit files.

/var/crash

Default depository for kernel crash dumps.

/var/cron

Log files for cron(1M).

/var/dmi

Solstice Enterprise Agents (SEA) Desktop Management Interface (DMI) run-time components.

/var/dt

dtlogin configuration files.

/var/ftp FTP server directory.

/var/inet

IPv6 router state files.

/var/krb5 Database and log files for Kerberos.

/var/ld

Configuration files for runtime linker.

/var/ldap

LDAP client configuration files.

/var/log

System log files.

/var/lp

Line printer subsystem logging information.

/var/mail

Directory where users' mail is kept.

/var/news

Community service messages. This is not the same as USENET-style news.

/var/nfs

NFS server log files.

/var/nis NIS+ databases.

/var/ntp

Network Time Protocol (NTP) server state directory.

/var/opt

Root of a subtree for varying files associated with optional software packages. An approved installation location for add-on system software and applications.

/var/preserve

Backup files for vi(1) and ex(1).

/var/run

Temporary files which are not needed across reboots. Only root may modify the contents of this directory.

/var/sadm

Databases maintained by the software package management utilities.

/var/sadm/system/logs

Status log files produced by software management functions and/or applications. For example, log files produced for product installation. An approved installation location for bundled Solaris software and for add-on system software and applications.

Standards, Environments, and Macros 51

/var/saf Service access facility logging and accounting files. /var/samba Log and lock files for Samba. /var/snmp SNMP status and configuration information. /var/spool Contains directories for files used in printer spooling, mail delivery, cron(1M), at(1), and so forth. /var/spool/clientmqueue sendmail(1M) client files. /var/spool/cron cron(1M) and at(1) spooling files. /var/spool/locks Spooling lock files. /var/spool/lp Line printer spool files. See lp(1). /var/spool/mqueue Mail queued for delivery. /var/spool/pkg Spooled packages. /var/spool/print LP print service client-side request staging area. /var/spool/samba Samba print queue. /var/spool/uucp Queued uucp(1C) jobs. /var/spool/uucppublic Files deposited by uucp(1C). /var/statmon Network status monitor files. /var/tmp Files that vary in size or presence during normal system operations. This directory is *not* cleared during the boot operation. An approved installation location for bundled Solaris software and for add-on system software and applications. /var/uucp uucp(1C) log and status files.

52 man pages section 5: Standards, Environments, and Macros • Last Revised 15 Nov 2001

	mesystem(3)
	/var/yp Databases needed for backwards compatibility with NIS and ypbind(1M); unnecessary after full transition to NIS+.
/usr File System	Because it is desirable to keep the root file system small and not volatile, on disk-based systems larger file systems are often mounted on /home, /opt, /usr, and /var.
	The file system mounted on /usr contains platform-dependent and platform-independent sharable files. The subtree rooted at /usr/share contains platform-independent sharable files; the rest of the /usr tree contains platform-dependent files. By mounting a common remote file system, a group of machines with a common platform may share a single /usr file system. A single /usr/share file system can be shared by machines of any platform. A machine acting as a file server can share many different /usr file systems to support several different architectures and operating system releases. Clients usually mount /usr read-only so that they do not accidentally change any shared files.
	The /usr file system contains the following subdirectories:
	/usr/4lib a.out libraries for the Binary Compatibility Package.
	/usr/5bin Symbolic link to the /usr/bin directory.
	/usr/X Symbolic link to the /usr/openwin directory.
	/usr/adm Symbolic link to the /var/adm directory.
	/usr/apache Apache executables, loadable modules, and documentation.
	/usr/aset Directory for Automated Security Enhancement Tools (ASET) programs and files.
	/usr/bin Platform‐dependent, user-invoked executables. These are commands users expect to be run as part of their normal \$PATH. For executables that are different on a 64-bit system than on a 32-bit system, a wrapper that selects the appropriate executable is placed here. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/ <i>packagename</i> /bin.
	/usr/bin/ia64 x86 64-bit, platform‐dependent, user-invoked executables. Note that ia64 is an example name; the actual name might be different. This directory should not be part of a user's \$PATH. A wrapper in /usr/bin should invoke the executable in this directory. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/ <i>packagename</i> /bin/ia64.

/usr/bin/sparcv9

SPARC 64-bit, platform‐dependent, user-invoked executables. This directory should not be part of a user's \$PATH. A wrapper in /usr/bin should invoke the executable in this directory. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/packagename/bin/sparcv9.

/usr/bin/subsystem

Platform‐dependent user-invoked executables that are associated with *subsystem*. These are commands users expect to be run as part of their normal \$PATH. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/bin.

/usr/bin/subsystem/ia64

x86 64-bit, platform‐dependent, user-invoked executables. Note that ia64 is an example name; the actual name might be different. This directory should not be part of a user's \$PATH. A wrapper in /usr/bin should invoke the executable in this directory. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/bin/ia64.

/usr/bin/subsystem/sparcv9

SPARC 64-bit, platform‐dependent, user-invoked executables. This directory should not be part of a user's \$PATH. A wrapper in /usr/bin should invoke the executable in this directory. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/packagename/bin/sparcv9.

/usr/subsystem/bin

Platform‐dependent user-invoked executables that are associated with *subsystem*. These are commands users expect to be run as part of their normal \$PATH. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/bin.

/usr/subsystem/bin/ia64

x86 64-bit, platform‐dependent, user-invoked executables. Note that ia64 is an example name; the actual name might be different. This directory should not be part of a user's \$PATH. A wrapper in /usr/bin should invoke the executable in this directory. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/bin/ia64.

/usr/subsystem/bin/sparcv9

SPARC 64-bit, platform‐dependent, user-invoked executables. This directory should not be part of a user's \$PATH. A wrapper in /usr/bin should invoke the executable in this directory. See isaexec(3C). An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/packagename/bin/sparcv9.

/usr/ccs C compilation system.

/usr/ccs/bin

C compilation commands and system utilities.

/usr/ccs/lib Symbolic link to /usr/lib.

/usr/demo

Demo programs and data.

/usr/dict

Symbolic link to the /usr/share/lib/dict directory, which contains the dictionary file used by the UNIX spell program.

/usr/dt

root of a subtree for CDE software.

/usr/dt/bin

Primary location for CDE system utilities.

/usr/dt/include Header files for CDE software.

/usr/dt/lib Libraries for CDE software.

/usr/dt/share/man

On-line reference manual pages for CDE software.

/usr/games

An empty directory, a remnant of the SunOS 4.0/4.1 software.

/usr/include

Include headers (for C programs).

/usr/iplanet

Directory server executables, loadable modules, and documentation.

/usr/j2se

Java 2 SDK executables, loadable modules, and documentation.

/usr/java*

Directories containing Java programs and libraries.

/usr/kernel

Subtree of platform‐dependent loadable kernel modules, not needed in the root filesystem. An approved installation location for bundled Solaris software.

/usr/kvm

A mount point, retained for backward compatibility, that formerly contained platform-specific binaries and libraries.

/usr/lib

Platform‐dependent libraries, various databases, commands and daemons not invoked directly by a human user. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib.

/usr/lib/64

Symbolic link to the most portable 64-bit Solaris interfaces.

/usr/lib/acct

Accounting scripts and binaries. See acct(1M).

/usr/lib/class

Scheduling‐ class-specific directories containing executables for priocntl(1) and dispadmin(1M).

/usr/lib/dict

Database files for spell(1).

/usr/lib/font

troff(1) font description files.

/usr/lib/fs

File system type dependent modules; generally not intended to be invoked directly by the user.

/usr/lib/ia64

x86 64–bit, platform‐dependent libraries, various databases, commands and daemons not invoked directly by a human user. Note that ia64 is an example name; the actual name might be different. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib/ia64.

/usr/lib/iconv

Conversion tables for iconv(1).

/usr/lib/libp Profiled libraries.

/usr/lib/locale

Localization databases.

/usr/lib/lp

Line printer subsystem databases and back-end executables.

/usr/lib/mail

Auxiliary programs for the mail(1) subsystem.

/usr/lib/netsvc

Internet network services.

/usr/lib/nfs

Auxiliary NFS-related programs and daemons.

/usr/lib/pics

Position Independent Code (PIC) archives needed to rebuild the run-time linker.

/usr/lib/refer

Auxiliary programs for refer(1).

/usr/lib/sa

Scripts and commands for the system activity report package. See sar(1).

/usr/lib/saf

Auxiliary programs and daemons related to the service access facility.

/usr/lib/sparcv9

SPARC 64-bit, platform‐dependent libraries, various databases, commands and daemons not invoked directly by a human user. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib/sparcv9.

/usr/lib/spell

Auxiliary programs and databases for spell(1). This directory is only present when the Binary Compatibility Package is installed.

/usr/lib/uucp

Auxiliary programs and daemons for uucp(1C).

/usr/lib/*subsystem*

Platform‐dependent libraries, various databases, commands and daemons that are associated with *subsystem* and that are not invoked directly by a human user. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib.

/usr/lib/subsystem/ia64

x86 64-bit, platform‐dependent libraries, various databases, commands and daemons that are associated with *subsystem* and that are not invoked directly by a human user. Note that ia64 is an example name; the actual name might be different. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib/ia64.

/usr/lib/subsystem/sparcv9

SPARC 64-bit, platform‐dependent libraries, various databases, commands and daemons that are associated with *subsystem* and that are not invoked directly by a human user. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib/sparcv9.

/usr/subsystem/lib

Platform‐dependent libraries, various databases, commands and daemons not invoked directly by a human user. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib.

/usr/subsystem/lib/ia64

x86 64–bit, platform‐dependent libraries, various databases, commands and daemons that are associated with *subsystem* and that are not invoked directly by a human user. Note that ia64 is an example name; the actual name might be different. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib/ia64.

/usr/subsystem/lib/sparcv9

SPARC 64-bit, platform‐dependent libraries, various databases, commands and daemons that are associated with *subsystem* and that are not invoked directly by a human user. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/lib/sparcv9.

/usr/local

Not part of the SVR4-based Solaris distribution. The /usr directory is exclusively for software bundled with the Solaris operating system. If needed for storing machine-local add-on software, create the directory /opt/local and make /usr/local a symbolic link to /opt/local. The /opt directory or filesystem is for storing add-on software to the system.

/usr/mail

Symbolic link to the /var/mail directory.

/usr/man

Symbolic link to the /usr/share/man directory.

/usr/net/servers

Entry points for foreign name service requests relayed using the network listener. See listen(1M).

/usr/news

Symbolic link to the /var/news directory.

/usr/oasys

Commands and files related to the Form and Menu Language Interpreter (FMLI) execution environment. See face(1).

/usr/old

Programs that are being phased out.

/usr/openwin

Installation or mount point for the OpenWindows software.

/usr/perl5

Perl 5 programs and documentation

/usr/platform

Subtree of platform‐specific objects which does not need to reside on the root filesystem. It contains a series of directories, one per supported platform. The

semantics of the series of directories is equivalent to /platform, except for subdirectories which do not provide utility under one or the other (for example, /platform/include is not needed).

/usr/platform/`uname -i`/include

Symbolic link to /../`uname -i`/include.Platform‐specific system (sys, vm) header files with semantics equivalent to /usr/include. An approved installation location for bundled Solaris software and for add-on system software.

/usr/platform/`uname -i`/kernel

Platform‐specific modules with semantics equivalent to /usr/kernel. An approved installation location for bundled Solaris software and for add-on system software.

```
/usr/platform/`uname -i`/lib
```

Platform‐specific daemon and shared objects with semantics equivalent to /usr/lib. An approved installation location for bundled Solaris software and for add-on system software.

/usr/platform/`uname -i`/lib/ia64

x86 64–bit, platform‐specific daemon and shared objects. Note that ia64 is an example name; the actual name might be different. An approved installation location for bundled Solaris software and for add-on system software.

/usr/platform/`uname -i`/lib/sparcv9

SPARC 64-bit, platform‐ specific daemon and shared objects. An approved installation location for bundled Solaris software and for add-on system software.

/usr/platform/`uname -i`/[s]mannum

Where *num* can be one of 3x, 1m, 4, 7d, or 9e. Platform‐specific system manual pages for documenting platform‐specific, shared objects, administration utilities, configuration files, special files/modules, and header files. An approved installation location for bundled Solaris software and for add-on system software.

/usr/platform/`uname -i`/sbin

Platform-specific system administration utilities with semantics equivalent to /usr/sbin. An approved installation location for bundled Solaris software and for add-on system software.

/usr/preserve

Symbolic link to the /var/preserve directory.

/usr/proc

Directory for the proc tools.

/usr/proc/bin

Contains links to SPARC Version 8 binaries in /usr/bin.

/usr/pub

Files for online man page and character processing.

/usr/sadm

System administration files and directories.

Standards, Environments, and Macros 59

/usr/sadm/bin

Binaries for the Form and Menu Language Interpreter (FMLI) scripts. See fmli(1).

/usr/sadm/install

Executables and scripts for package management.

/usr/sbin

Platform‐dependent executables for system administration, expected to be run only by system administrators. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/sbin.

/usr/sbin/install.d

Custom Jumpstart scripts and executables.

/usr/sbin/static

Statically linked version of selected programs from /usr/bin and /usr/sbin. These are used to recover from broken dynamic linking and before all pieces necessary for dynamic linking are present.

/usr/sbin/sparc7 and sparc9

32-bit and 64-bit versions of commands.

/usr/sfw

GNU and open source executables, libraries, and documentation.

/usr/sbin/subsystem

Platform‐dependent executables for system administration, expected to be run only by system administrators, and associated with *subsystem*. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/sbin.

/usr/subsystem/sbin

Platform‐dependent executables for system administration, expected to be run only by system administrators, and associated with *subsystem*. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/sbin.

/usr/share

Platform‐independent sharable files. An approved installation location for bundled Solaris software.

/usr/share/admserv5.1

iPlanet Console and Administration Server documentation.

/usr/share/audio

Sample audio files.

/usr/share/ds5

iPlanet Server documentation.

/usr/share/lib

Platform‐independent sharable databases. An approved installation location for bundled Solaris software.

60 man pages section 5: Standards, Environments, and Macros • Last Revised 15 Nov 2001

/usr/share/lib/dict Contains word list for spell(1).

/usr/share/lib/keytables Keyboard layout description tables.

/usr/share/lib/mailx Help files for mailx(1).

/usr/share/lib/nterm
nroff(1) terminal tables.

/usr/share/lib/pub Character set data files.

/usr/share/lib/tabset Tab setting escape sequences.

/usr/share/lib/terminfo
Terminal description files for terminfo(4).

/usr/share/lib/tmac

Macro packages and related files for text processing tools, for example, nroff(1) and troff(1).

/usr/share/lib/zoneinfo Time zone information.

/usr/share/[s]man

Platform‐ independent sharable manual pages. An approved installation location for bundled Solaris software. The analogous location for add-on system software or for applications is /opt/*packagename*/[s]man.

/usr/share/src

Source code for kernel, utilities, and libraries.

/usr/snadm

Files related to system and network administration.

/usr/spool

Symbolic link to the /var/spool directory.

/usr/src

Symbolic link to the /usr/share/src directory.

/usr/tmp

Symbolic link to the var/tmp directory.

/usr/ucb

Berkeley compatibility package binaries.

/usr/ucbinclude

Berkeley compatibility package headers.

/usr/ucblib

Berkeley compatibility package libraries.

Standards, Environments, and Macros 61

/usr/vmsys

Commands and files related to the Framed Access Command Environment (FACE) programs. See face(1).

/usr/xpg4

Directory for POSIX-compliant utilities.

NAME	fnmatch – file name pattern matching
DESCRIPTION	The pattern matching notation described below is used to specify patterns for matching strings in the shell. Historically, pattern matching notation is related to, but slightly different from, the regular expression notation. For this reason, the description of the rules for this pattern matching notation is based on the description of regular expression notation described on the regex(5) manual page.
Patterns Matching a Single Character	The following <i>patterns matching a single character</i> match a single character: <i>ordinary characters, special pattern characters</i> and <i>pattern bracket expressions</i> . The pattern bracket expression will also match a single collating element.
	An ordinary character is a pattern that matches itself. It can be any character in the supported character set except for NUL, those special shell characters that require quoting, and the following three special pattern characters. Matching is based on the bit pattern used for encoding the character, not on the graphic representation of the character. If any character (ordinary, shell special, or pattern special) is quoted, that pattern will match the character itself. The shell special characters always require quoting.
	When unquoted and outside a bracket expression, the following three characters will have special meaning in the specification of patterns:
	? A question-mark is a pattern that will match any character.
	* An asterisk is a pattern that will match multiple characters, as described in Patterns Matching Multiple Characters, below.
	[The open bracket will introduce a pattern bracket expression.
	The description of basic regular expression bracket expressions on the regex(5) manual page also applies to the pattern bracket expression, except that the exclamation-mark character (!) replaces the circumflex character (^) in its role in a <i>non-matching list</i> in the regular expression notation. A bracket expression starting with an unquoted circumflex character produces unspecified results.
	The restriction on a circumflex in a bracket expression is to allow implementations that support pattern matching using the circumflex as the negation character in addition to the exclamation-mark. A portable application must use something like $[\^2]$ to match either character.
	When pattern matching is used where shell quote removal is not performed (such as in the argument to the find -name primary when find is being called using one of the exec functions, or in the <i>pattern</i> argument to the fnmatch(3C) function, special characters can be escaped to remove their special meaning by preceding them with a backslash character. This escaping backslash will be discarded. The sequence \\ represents one literal backslash. All of the requirements and effects of quoting on ordinary, shell special and special pattern characters will apply to escaping in this context.

fnmatch(5)

Both quoting and escaping are described here because pattern matching must work in three separate circumstances:

 Calling directly upon the shell, such as in pathname expansion or in a case statement. All of the following will match the string or file abc:

abc	"abc"	a"b"c	a\bc	a[b]c
a["b"]c	a[\b]c	a["\b"]c	a?c	a*c
The following	ng will not:			
"a?c"		a*c	a\[b]	C

- Calling a utility or function without going through a shell, as described for find(1) and the function fnmatch(3C)
- Calling utilities such as find, cpio, tar or pax through the shell command line. In this case, shell quote removal is performed before the utility sees the argument. For example, in:

find /bin -name e\c[\h]o -print after quote removal, the backslashes are presented to find and it treats them as escape characters. Both precede ordinary characters, so the c and h represent themselves and echo would be found on many historical systems (that have it in /bin). To find a file name that contained shell special characters or pattern characters, both quoting and escaping are required, such as:

pax -r . . . "*a\ (\?"to extract a filename ending with a (?.

Conforming applications are required to quote or escape the shell special characters (sometimes called metacharacters). If used without this protection, syntax errors can result or implementation extensions can be triggered. For example, the KornShell supports a series of extensions based on parentheses in patterns; see ksh(1)

Patterns Matching Multiple Characters The following rules are used to construct *patterns matching multiple characters* from *patterns matching a single character*:

- The asterisk (*) is a pattern that will match any string, including the null string.
- The concatenation of *patterns matching a single character* is a valid pattern that will match the concatenation of the single characters or collating elements matched by each of the concatenated patterns.
- The concatenation of one or more *patterns matching a single character* with one or more asterisks is a valid pattern. In such patterns, each asterisk will match a string of zero or more characters, matching the greatest possible number of characters that still allows the remainder of the pattern to match the string.
- 64 man pages section 5: Standards, Environments, and Macros Last Revised 28 Mar 1995

Since each asterisk matches zero or more occurrences, the patterns a*b and a**b have identical functionality.

Examples:

a[bc]	matches the strings ab and ac.
a*d	matches the strings ad, abd and abcd, but not the string abc.
a*d*	matches the strings ad, abcd, abcdef, aaaad and adddd.
*a*d	matches the strings ad, abcd, efabcd, aaaad and adddd.

Patterns Used for Filename Expansion The rules described so far in Patterns Matching Multiple Characters and Patterns Matching a Single Character are qualified by the following rules that apply when pattern matching notation is used for filename expansion.

- The slash character in a pathname must be explicitly matched by using one or more slashes in the pattern; it cannot be matched by the asterisk or question-mark special characters or by a bracket expression. Slashes in the pattern are identified before bracket expressions; thus, a slash cannot be included in a pattern bracket expression used for filename expansion. For example, the pattern a [b/c] d will not match such pathnames as abd or a/d. It will only match a pathname of literally a [b/c] d.
- 2. If a filename begins with a period (.), the period must be explicitly matched by using a period as the first character of the pattern or immediately following a slash character. The leading period will not be matched by:
 - the asterisk or question-mark special characters
 - a bracket expression containing a non-matching list, such as:
 - [!a] a range expression, such as:
 - [%–0] or a character class expression, such as:

[[:punct:]] It is unspecified whether an explicit period in a bracket expression matching list, such as:

- [.abc] can match a leading period in a filename.
- 3. Specified patterns are matched against existing filenames and pathnames, as appropriate. Each component that contains a pattern character requires read permission in the directory containing that component. Any component, except the last, that does not contain a pattern character requires search permission. For example, given the pattern:

/foo/bar/x*/bamsearch permission is needed for directories / and foo, search and read permissions are needed for directory bar, and search permission is

fnmatch(5)

needed for each x^* directory. If the pattern matches any existing filenames or pathnames, the pattern will be replaced with those filenames and pathnames, sorted according to the collating sequence in effect in the current locale. If the pattern contains an invalid bracket expression or does not match any existing filenames or pathnames, the pattern string is left unchanged.

SEE ALSO find(1), ksh(1), fnmatch(3C), regex(5)

fns – overview of FNS
Federated Naming Service (FNS) provides a method for federating multiple naming services under a single, simple interface for the basic naming operations. The service supports resolution of <i>composite</i> names, names that span multiple naming systems, through the naming interface. In addition to the naming interface, FNS also specifies <i>policies</i> for composing names in the enterprise namespace. See fns_policies(5) and fns_initial_context(5).
Fundamental to the FNS model are the notions of composite names and <i>contexts</i> . A context provides operations for:
 associating (binding) names to objects resolving names to objects removing bindings, listing names, renaming and so on.
A context contains a set of names to reference bindings. A reference contains a list of communication end-points. Every naming operation in the FNS interface is performed on a context object.
The federated naming system is formed by contexts from one naming system being bound in the contexts of another naming system. Resolution of a composite name proceeds from contexts within one naming system to those in the next, until the name is resolved.
XFN is X/Open Federated Naming. The programming interface and policies that FNS supports are specified by XFN. See xfn(3XFN) and fns_policies(5).
A composite name is a name that spans multiple naming systems. It consists of an ordered list of components. Each component is a name from the namespace of a single naming system. FNS defines the syntax for constructing a composite name using names from component naming systems. Individual naming systems are responsible for the syntax of each component.
The syntax for composite names is that components are composed left to right using the slash character ('/') as the component separator. For example, the composite name /Wiz.Com/site/Oceanview.East consists of four components: , Wiz.COM, site, and Oceanview.East.See fns_policies(5) and fns_initial_context(5) for more examples of composite names.
FNS is useful for the following reasons:
 A single uniform naming interface is provided to clients for accessing naming services. Consequently, the addition of new naming services does not require changes to applications or existing naming services. Furthermore, applications that use FNS will be portable across platforms because the interface exported by FNS is XFN, a public, open interface endorsed by other vendors and by the X/Open Company.

fns(5)

	 Names can be composed in a uniform way (that is, FNS supports a model in which composite names are constructed in a uniform syntactic way and can have any number of components).
	 Coherent naming is encouraged through the use of shared contexts and shared names.
FNS and Naming Systems	FNS has support for NIS+, NIS, and files as enterprise-level naming services. This means that FNS implements the enterprise-level policies using NIS+, NIS, and files. FNS also supports DNS and X.500 (via DAP or LDAP) as global naming services, as well as support for federating NIS+ and NIS with DNS and X.500. See the corresponding individual man page for information about the implementation for a specific naming service.
SEE ALSO	<pre>nis+(1), xfn(3XFN), fns_dns(5), fns_files(5), fns_initial_context(5), fns_nis(5), fns_nis+(5), fns_policies(5), fns_references(5), fns_x500(5)</pre>

NAME | fns_dns - overview of FNS over DNS implementation

DESCRIPTION

Federated Naming Service (FNS) provides a method for federating multiple naming services under a single, simple interface for the basic naming operations. One of the naming services supported by FNS is the Internet Domain Name System, or DNS (see in.named(1M)). DNS is a hierarchical collection of name servers that provide the Internet community with host and domain name resolution. FNS uses DNS to name entities globally. Names can be constructed for any enterprise that is accessible on the Internet; consequently, names can also be constructed for objects exported by these enterprises.

FNS provides the XFN interface for performing naming resolution on DNS domains and hosts. In addition, enterprise namespaces such as those served by NIS+ and NIS can be federated with DNS by adding TXT records to DNS. To federate an NIS+ or NIS namespace under DNS, you first obtain the root reference for the NIS+ hierarchy or NIS domain. This reference is referred to as the *next naming system reference* because it refers to the *next* naming system beneath the DNS domain. This reference contains information about how to communicate with the NIS+ or NIS servers and has the following format:

<domainname> <server name> [<server address>]

where <domainname> is the fully qualified domain name. Notice that NIS+ and NIS have slightly different syntaxes for domain names. For NIS+, the fully qualified domain name is case-insensitive and terminated by a dot character ('.'). For NIS, the fully qualified domain name is case-sensitive and is *not* terminated by a dot character. For both NIS+ and NIS, *<server address>* is optional. If it is not supplied, a host name lookup will be performed to get the machine's address.

For example, if the machine wiz-nisplus-server with address 133.33.33 serves the NIS+ domain wiz.com., the reference would look like this:

wiz.com.wiz-nisplus-server133.33.33.33

For NIS, the reference information is of the form:

<domainname> <server name>

For example, if the machine woz-nis-server serves the NIS domain Woz.COM, the reference would look like this:

Woz.COM woz-nis-server

Standards, Environments, and Macros 69

fns_dns(5)

After obtaining this information, you then edit the DNS table (see in.named(1M)) and add a TXT record with this reference information. The TXT record must be associated with a DNS domain that includes an NIS record. For example, the reference information shown in the examples above would be entered as follows. For NIS+: TXT "XFNNISPLUS wiz.com. wiz-nisplus-server 133.33.33.33" For NIS: TXT "XFNNIS woz.com woz-nis-server" Note the mandatory double quotes (' " ') delimiting the contents of the TXT record. After making any changes to the DNS table, you must notify the server by either restarting it or sending it a signal to reread the table: #kill -HUP `cat /etc/named.pid` This update effectively adds the next naming system reference to DNS. You can look up this reference using fnlookup(1) to see if the information has been added properly. For example, the following command looks up the next naming system reference of the DNS domain Wiz.COM: #fnlookup -v .../Wiz.COM/ Note the mandatory trailing slash ('/'). After this administrative step has been taken, clients outside of the NIS+ hierarchy or NIS domain can access and perform operations on the contexts in the NIS+ hierarchy or NIS domain. Foreign NIS+ clients access the hierarchy as unauthenticated NIS+ clients. Continuing the example above, and assuming that NIS+ is federated underneath the DNS domain Wiz. COM, you can now list the root of the NIS+ enterprise using the command: #fnlist .../Wiz.COM/ SEE ALSO fnlist(1), fnlookup(1), nis+(1), in.named(1M), ypserv(1M), xfn(3XFN), fns(5), fns nis(5), fns nis+(5), fns references(5), fns x500(5)

70 man pages section 5: Standards, Environments, and Macros • Last Revised 22 Nov 1996

NAME	fns_files - overview of FNS over files implementation
DESCRIPTION	The Federated Naming Service (FNS) provides a method for federating multiple naming services under a single, simple interface for the basic naming operations. One of the naming services supported by FNS is /etc files. FNS provides the XFN interface for performing naming and attribute operations on FNS enterprise objects (organization, site, user, host, and service objects), using files as the naming service. FNS stores bindings for these objects in files and uses them in conjunction with existing /etc files objects.
FNS Policies and /etc Files	FNS defines policies for naming objects in the federated namespace (see fns_policies(5)). At the enterprise level, FNS policies specify naming for organizations, hosts, users, sites, and services. The enterprise-level naming service provides contexts to allow other objects to be named relative to these objects.
	The organizational unit namespace provides a hierarchical namespace for naming subunits of an enterprise. In /etc files, there is no concept of an organization. Hence, with respect to /etc files as the naming service, there is a single organizational unit context that represents the entire system. Users in an FNS organizational unit correspond to the users in the /etc/passwd file. FNS provides a context for each user in the /etc/passwd file.
	Hosts in an FNS organizational unit correspond to the hosts in the /etc/hosts file. FNS provides a context for each host in the /etc/hosts file.
Security Considerations	Changes to the FNS information (using the commands fncreate(1M), fncreate_fs(1M), fnbind(1), fndestroy(1M) and fnunbind(1)) can be performed only by the privileged users on the system that exports the /var/fn directory. Also, based on the UNIX user IDs, users are allowed to modify their own contexts, bindings, and attributes, from any machine that mounts the /var/fn directory.
	For example, the command fncreate(1M) creates FNS related files and directories in the system on which the command is executed. Hence, the invoker of the fncreate(1M) command must have super-user privileges in order to create the user, host, site, and service contexts. However, a user could use the fnunbind(1) command to create calendar bindings in the user's own context, as in this example:
	fnbind -r thisuser/service/calendar onc_calendar onc_cal_str jsmith@beatrix
	The files object name that corresponds to an FNS composite name can be obtained using fnlookup(1) and fnlist(1).
USAGE	The files used for storing FNS information are placed in the directory /var/fn. The machine on which /var/fn is located has access to the FNS file. The FNS information can be made accessible to other machines by exporting /var/fn. Client machines that NFS mount the /var/fn directory would then be able to access the FNS information.

fns_files(5)

SEE ALSO	<pre>fnbind(1), fnlist(1), fnlookup(1), fnunbind(1), fncreate(1M), fncreate_fs(1M), fndestroy(1M), xfn(3XFN), fns(5),</pre>
	<pre>fncreate_rs(nw), fndestroy(nw), xfn(3xnA), fns(3), fns_initial_context(5), fns_nis(5), fns_nis+(5), fns_policies(5),</pre>
	<pre>fns_references(5)</pre>

72 man pages section 5: Standards, Environments, and Macros • Last Revised 13 Dec 1996

NAME | fns_initial_context - overview of the FNS Initial Context

DESCRIPTION

Every FNS name is interpreted relative to some context, and every FNS naming operation is performed on a context object. The FNS programming interface (XFN) provides a function that allows the client to obtain an Initial Context object. The Initial Context provides the initial pathway to other FNS contexts. FNS defines a set of bindings that the client can expect to find in this context,

FNS assumes that for every process:

- 1. There is a user associated with the process when
 - fn ctx handle from initial() is invoked. This association is based on the effective uid of the process. In the following discussion this user is denoted by *U*. The association of user to process may change during the life of a process but does not affect the context handle originally returned by fn ctx handle from initial().
- 2. The process is running on a host when fn ctx handle from initial() is invoked. In the following discussion this host is denoted by *H*.

The following atomic names can appear in the Initial Context:

	thishost	thisorgunit
thisens	myself	myorgunit
myens	orgunit	site
user	host	

Except for . . . , these names with an added underscore ('_') prefix are also in the Initial Context and have the same binding as their counterpart (for example, thishost and thishost have the same binding). In addition, org has the same binding as orgunit, and thisuser has the same binding as myself. The bindings for these names are summarized in the following table.

Some of these names may not necessarily appear in all Initial Contexts. For example, a process owned by the super-user of a machine does not have any of the user-related bindings. Or, for another example, an installation that has not set up a site namespace will not have the site-related bindings.

	global context for resolving DNS or X.500 names. Synonym: /
thishost	<i>H</i> 's host context. Synonym: _thishost
thisens	the enterprise root of <i>H</i> . Synonym: _thisens
thisorgunit	<i>H</i> 's distinguished organizational unit context. In Solaris, this is <i>H</i> 's NIS+ home domain. Synonym: _thisorgunit

Standards, Environments, and Macros 73

fns_initial_context(5)

	myself	U's user context. Synonyms: _myself, thisuser
	myens	the enterprise root of U. Synonym: _myens
	myorgunit	<i>U</i> 's distinguished organizational unit context. In Solaris, this is <i>U</i> 's NIS+ home domain. Synonym: _myorgunit
	user	the context in which users in the same organizational unit as <i>H</i> are named. Synonym: _user
	host	the context in which hosts in the same organizational unit as <i>H</i> are named. Synonym: _host
	org	the root context of the organizational unit namespace in <i>H</i> 's enterprise. In Solaris, this corresponds to the NIS+ root domain. Synonyms: orgunit, _orgunit
	site	the root context of the site namespace in <i>H</i> 's enterprise, if the site namespace has been configured. Synonym: _site
EXAMPLES	EXAMPLE 1 Names beginning with the enterprise root	
	The types of objects that may be named relative to the enterprise root are user, host, service, organizational unit, file, and site. Here are some examples of names that begin with the enterprise root:	
	thisens/orgunit/multimedia.servers.engineering names an organizational unit multimedia.servers.engineering in H's enterprise.	
	thisens/site/northwing names the north wing site, enterprise.	.floor3.admin on the third floor of the administrations building in H 's
	myens/user/hdiffie names the user hdiffie in	n <i>U</i> 's enterprise.
	myens/service/teletax names the teletax service	e of U 's enterprise.
	EXAMPLE 2 Names beginning wi	th organizational unit names
	user, host, service, file, and sit organizational unit names (eit	be named relative to an organizational unit name are: e. Here are some examples of names that begin with her explicitly via org, or implicitly via thisorgunit or ts relative to organizational unit names when resolved in
	names a conference room v	inance/site/videoconference.northwing videoconference located in the north wing of the site vational unit accounts_payable.finance.

74 man pages section 5: Standards, Environments, and Macros • Last Revised 1 Nov 1994

fns_initial_context(5)

EXAMPLE 2 Names beginning with organizational unit names (<i>Continued</i>)
org/finance/user/mjones names a user mjones in the organizational unit finance.
<pre>org/finance/host/inmail names a machine inmail belonging to the organizational unit finance.</pre>
org/accounts_payable.finance/fs/pub/blue-and-whites/FY92-124 names a file pub/blue-and-whites/FY92-124 belonging to the organizational unit accounts_payable.finance.
<pre>org/accounts_payable.finance/service/calendar names the calendar service of the organizational unit accounts_payable.finance. This might manage the meeting schedules of the organizational unit.</pre>
thisorgunit/user/cmead names the user cmead in <i>H</i> 's organizational unit.
<pre>myorgunit/fs/pub/project_plans/widget.ps names the file pub/project_plans/widget.ps exported by U's organizational unit's file system.</pre>
EXAMPLE 3 Names beginning with site names
The types of objects that may be named relative to a site name are users, hosts, services, and files. Here are some examples of names that begin with site names via site, and name objects relative to sites when resolved in the Initial Context:
<pre>site/b5.mtv/service/printer/speedy names a printer speedy in the b5.mtv site.</pre>
<pre>site/admin/fs/usr/dist names a file directory usr/dist available in the site admin.</pre>
EXAMPLE 4 Names beginning with user names
The types of objects that may be named relative to a user name are services and files. Here are some examples of names that begin with user names (explicitly via user or implicitly via thisuser), and name objects relative to users when resolved in the Initial Context:
user/jsmith/service/calendar names the calendar service of the user jsmith.
user/jsmith/fs/bin/games/riddles names the file bin/games/riddles of the user jsmith.
thisuser/service/printer names the printer service of <i>U</i> .

fns_initial_context(5)

	EXAMPLE 5 Names beginning with host names
	The types of objects that may be named relative to a host name are services and files. Here are some examples of names that begin with host names (explicitly via host or implicitly via thishost), and name objects relative to hosts when resolved in the Initial Context:
	host/mailhop/service/mailbox names the mailbox service associated with the machine mailhop.
	<pre>host/mailhop/fs/pub/saf/archives.91 names the directory pub/saf/archives.91 found under the root directory of the machine mailhop.</pre>
	thishost/service/printer names the printer service of <i>H</i> .
SEE ALSO	<pre>nis+(1), geteuid(2), fn_ctx_handle_from_initial(3XFN), xfn(3XFN), fns(5), fns_policies(5)</pre>

NAME	fns_nis+ – overview of FNS over NIS+ implementation
DESCRIPTION	Federated Naming Service (FNS) provides a method for federating multiple naming services under a single, simple interface for the basic naming operations. One of the naming services supported by FNS is NIS+, the enterprise-wide information service in Solaris (see nis+(1)). FNS provides the XFN interface for performing naming and attribute operations on FNS enterprise objects (organization, site, user, host, and service objects) using NIS+. FNS stores bindings for these objects in NIS+ and uses them in conjunction with existing NIS+ objects.
FNS Policies and NIS+	FNS defines policies for naming objects in the federated namespace (see fns_policies(5)). At the enterprise level, FNS policies specify naming for organizations, hosts, users, sites, and services. The enterprise-level naming service provides contexts to allow other objects to be named relative to these objects.
	The organizational unit namespace provides a hierarchical namespace for naming subunits of an enterprise. An organizational unit maps to an NIS+ domain. Organizational unit names can be either fully qualified NIS+ domain names or relatively NIS+ domain names. If a terminal dot is present in the name, it is treated as a fully qualified name. Otherwise, the name is resolved relative to the root NIS+ domain.
	Users in the NIS+ namespace are found in the passwd.org_dir table of an NIS+ domain. Users in an FNS organizational unit correspond to the users in the passwd.org_dir table of the corresponding NIS+ domain. FNS provides a context for each user in the passwd.org_dir table.
	Hosts in the NIS+ namespace are found in the hosts.org_dir table of an NIS+ domain. Hosts in an FNS organizational unit correspond to the hosts in the hosts.org_dir table of the corresponding NIS+ domain. FNS provides a context for each host in the hosts.org_dir table.
	In NIS+, users and hosts have a notion of a <i>home domain</i> . It is the primary NIS+ domain that maintains information associated with them. A user or host's home domain can be determined directly using its NIS+ principal name, which is composed of the atomic user (login) name or the atomic host name, and the name of the NIS+ home domain. For example, user jsmith with home domain wiz.com has an NIS+ principal name, jsmith.wiz.com.
	A user's NIS+ home domain corresponds to the user's FNS organizational unit and determines the binding for myens and myorgunit.
	A host's NIS+ home domain corresponds to the host's FNS organizational unit and determines the binding for thisens, thisorgunit, user, and host.

fns_nis+(5)	
Federating NIS+ with DNS or X.500	Federating NIS+ with the global naming systems DNS or X.500 makes NIS+ contexts accessible outside of an NIS+ hierarchy. To enable the federation, the administrator must first add address information in either DNS or X.500 (see fns_dns(5) and fns_x500(5)). After this administrative step has been taken, clients outside of the NIS+ hierarchy can access contexts and perform operations from outside the hierarchy as an unauthenticated NIS+ client.
NIS+ Security	The command fncreate(1M) creates NIS+ tables and directories in the NIS+ hierarchy associated with the domain of the host on which it executes. The invoker of fncreate(1M) and other FNS commands is expected to have the necessary NIS+ credentials. (See nis+(1) and nisdefaults(1)). The environment variable NIS_GROUP of the process specifies the group owner for the NIS+ objects thus created. In order to facilitate administration of the NIS+ objects, NIS_GROUP should be set to the name of the NIS+ administration group for the domain prior to executing fncreate(1M) and other FNS commands. Changes to NIS+-related properties, including default access control rights, could be effected using NIS+ administration tools and interfaces after the context has been created. The NIS+ object name that corresponds to an FNS composite name can be obtained using fnlookup(1) and fnlist(1).
SEE ALSO	<pre>fnlist(1), fnlookup(1), nis+(1), nischgrp(1), nischmod(1), nischown(1), nisdefaults(1), nisls(1), fncreate(1M), xfn(3XFN), fns(5), fns_dns(5), fns_files(5), fns_initial_context(5), fns_nis(5), fns_policies(5), fns_references(5), fns_x500(5)</pre>

NAME	fns_nis - overview of FNS over NIS (YP) implementation
DESCRIPTION	Federated Naming Service (FNS) provides a method for federating multiple naming services under a single, simple interface for the basic naming operations. One of the naming services supported by FNS is NIS (YP), the enterprise-wide information services in Solaris (see ypcat(1), ypmatch(1), ypfiles(4)). FNS provides the XFN interface for performing naming and attribute operations on FNS enterprise objects (organization, site, user, host and service objects) using NIS. FNS stores bindings for these objects in NIS and uses them in conjunction with existing NIS objects.
FNS Policies and NIS	FNS defines policies for naming objects in the federated namespace (see fns_policies(5)). At the enterprise level, FNS policies specify naming for organizations, hosts, users, sites, and services. The enterprise-level naming service provides contexts to allow other objects to be named relative to these objects.
	The FNS organizational unit namespace provides a hierarchical namespace for naming subunits of an enterprise. However, NIS does not support a hierarchical organizational structure. Therefore, a NIS domain maps to a single organizational unit in the FNS namespace.
	Users in an FNS organizational unit correspond to the users in the passwd.byname map of the corresponding NIS domain. FNS provides a context for each user in the passwd.byname map.
	Hosts in an FNS organizational unit correspond to the hosts in the hosts.byname map of the corresponding NIS domain. FNS provides a context for each host in the hosts.byname map.
Federating NIS with DNS or X.500	Federating NIS with the global naming systems DNS or X.500 makes NIS contexts accessible outside of an NIS domain. To enable the federation, the administrator must first add address information in either DNS or X.500 (see fns_dns(5) and fns_x500(5)). After this administrative step has been taken, clients outside of the NIS domain can access contexts and perform operations.
Security Considerations	Changes to the FNS information (using the commands fncreate(1M), fncreate_fs(1M), fncreate_printer(1M), fnbind(1), fndestroy(1M), fncheck(1M), and fnunbind(1)) can be performed only by the privileged users on the NIS master server that maintains the FNS information.
	For example, the command fncreate(1M) creates the NIS map for the associated NIS domain in the system on which it is executed. Hence, the command must be run by a privileged user either on the NIS master server or on a system that will serve as a NIS master server for FNS.
	The NIS object name that corresponds to an FNS composite name can be obtained using fnlookup(1) and fnlist(1).

fns_nis(5)

SEE ALSO	<pre>fnbind(1), fnlist(1), fnlookup(1), fnunbind(1), ypcat(1), ypmatch(1),</pre>
	<pre>fncheck(1M), fncreate(1M), fncreate_fs(1M), fncreate_printer(1M),</pre>
	<pre>fndestroy(1M), xfn(3XFN), ypfiles(4), fns(5), fns_dns(5), fns_files(5),</pre>
	<pre>fns_initial_context(5), fns_nis+(5), fns_policies(5), fns_references(5),</pre>
	fns_x500(5)

NAME | fns_policies - overview of the FNS Policies

DESCRIPTION

FNS defines policies for naming objects in the federated namespace. The goal of these policies is to allow easy and uniform composition of names. The policies use the basic rule that objects with narrower scopes are named relative to objects with wider scopes.

FNS policies are described in terms of the following three categories: global, enterprise, and application.

Global naming service	A global naming service is a naming service that has world-wide scope. Internet DNS and X.500 are examples of global naming services. The types of objects named at this global level are typically countries, states, provinces, cities, companies, universities, institutions, and government departments and ministries. These entities are referred to as <i>enterprises</i> .
Enterprise-level naming service	Enterprise-level naming services are used to name objects within an enterprise. Within an enterprise, there are naming services that provide contexts for naming common entities such as organizational units, physical sites, human users, and computers. Enterprise-level naming services are bound below the global naming services. Global naming services provide contexts in which the root contexts of enterprise-level naming services can be bound.
Application-level naming service	Application-level naming services are incorporated in applications offering services such as file service, mail service, print service, and so on. Application-level naming services are bound below enterprise naming services. The enterprise-level naming services provide contexts in which contexts of application-level naming services can be bound.

FNS has policies for global and enterprise naming. Naming within applications is left to individual applications or groups of related applications and not specified by FNS.

FNS policy specifies that DNS and X.500 are global naming services that are used to name enterprises. The global namespace is named using the name A DNS name or an X.500 name can appear after the Support for federating global naming services is planned for a future release of FNS.

fns_policies(5)

Within an enterprise, there are namespaces for organizational units, sites, hosts, users, files and services, referred to by the names orgunit, site, host, user, fs, and service. In addition, these namespaces can be named using these names with an added underscore ('_') prefix (for example, host and _host have the same binding). The following table summarizes the FNS policies.

Context	Subordinate	Parent
Туре	Context	Context
org unit	site	enterprise root
	user	
	host	
	file system	
	service	
site	user	enterprise root
	host	org unit
	file system	
	service	
user	service	enterprise root
	file system	org unit
host	service	enterprise root
	file system	org unit
service	not specified	enterprise root
		org unit
		site
		user
		host
file system	none	enterprise root
		org unit
		site
		user
		host

⁸² man pages section 5: Standards, Environments, and Macros • Last Revised 4 Nov 1994

	I
	In Solaris, an organizational unit name corresponds to an NIS+ domain name and is identified using either the fully-qualified form of its NIS+ domain name, or its NIS+ domain name relative to the NIS+ root. Fully-qualified NIS+ domain names have a terminal dot ('.'). For example, assume that the NIS+ root domain is "Wiz.COM." and "sales" is a subdomain of that. Then, the names org/sales.Wiz.COM. and org/sales both refer to the organizational unit corresponding to the same NIS+ domain sales.Wiz.COM.
	User names correspond to names in the corresponding NIS+ <i>passwd.org_dir</i> table. The file system context associated with a user is obtained from his entry in the NIS+ <i>passwd.org_dir</i> table.
	Host names correspond to names in the corresponding NIS+ <i>hosts.org_dir</i> table. The file system context associated with a host corresponds to the files systems exported by the host.
EXAMPLES	EXAMPLE 1 The types of objects that may be named relative to an organizational unit name are: user, host, service, file, and site. Here are some examples of names name objects relative to organizational unit names:
	<pre>org/accounts_payable.finance/site/videoconference.northwing names a conference room videoconference located in the north wing of the site associated with the organizational unit accounts_payable.finance.</pre>
	org/finance/user/mjones names a user mjones in the organizational unit finance.
	org/finance/host/inmail names a machine inmail belonging to the organizational unit finance.
	org/accounts_payable.finance/fs/pub/blue-and-whites/FY92-124 names a file pub/blue-and-whites/FY92-124 belonging to the organizational unit accounts_payable.finance.
	org/accounts_payable.finance/service/calendar names the calendar service of the organizational unit accounts_payable.finance. This might manage the meeting schedules of the organizational unit.
	EXAMPLE 2 The types of objects that may be named relative to a site name are services and files. Here are some examples of names that name objects relative to sites:
	<pre>site/b5.mtv/service/printer/speedy names a printer speedy in the b5.mtv site.</pre>
	<pre>site/admin/fs/usr/dist names a file directory usr/dist available in the site admin.</pre>
	EXAMPLE 3 The types of objects that may be named relative to a user name are services and files. Here are some examples of names that name objects relative to users:
	user/jsmith/service/calendar names the calendar service of the user jsmith.
	Standards, Environments, and Macros 83

fns_policies(5)	
	EXAMPLE 3 The types of objects that may be named relative to a user name are services and files. Here are some examples of names that name objects relative to users: (<i>Continued</i>)
	user/jsmith/fs/bin/games/riddles names the file bin/games/riddles of the user jsmith.
	EXAMPLE 4 The types of objects that may be named relative to a host name are services and files. Here are some examples of names that name objects relative to hosts:
	host/mailhop/service/mailbox names the mailbox service associated with the machine mailhop.
	<pre>host/mailhop/fs/pub/saf/archives.91 names the directory pub/saf/archives.91 found under the root directory of the machine mailhop.</pre>
SEE ALSO	<pre>fncreate(1M), nis+(1), xfn(3XFN), fns(5), fns_initial_context(5), fns_references(5)</pre>

NAME	fns_references - overview of FNS References		
DESCRIPTION	Every composite name in FNS is bound to a <i>reference</i> . A reference consists of a type and a list of addresses. The reference type is used to identify the type of object.		
	An address is something that can be used with some communication mechanism to invoke operations on an object or service. Multiple addresses are intended to identify multiple communication endpoints for a single conceptual object or service. Each address in a reference consists of an address type and an opaque buffer. The address type determines the format and interpretation of the address data. Together, the address's type and data specify how to reach the object. Many communication mechanisms are possible; FNS does not place any restrictions on them.		
		e reference and address types that are currently defined. ed with the Federated Naming Group at SunSoft.	
Reference Types	All reference types use the FN_ID_STRING identifier format unless otherwise qualified.		
	onc_fn_enterprise	Enterprise root context.	
	onc_fn_organization	A context for naming objects related to an organizational unit.	
	onc_fn_hostname	A context for naming hosts.	
	onc_fn_username	A context for naming users.	
	onc_fn_user	A context for naming objects related to a user.	
	onc_fn_host	A context for naming objects related to a computer.	
	onc_fn_site	A context for naming sites.	
	onc_fn_service	A context for naming services.	
	onc_fn_nsid	A context for naming namespace identifiers.	
	onc_fn_generic	A context for naming application-specific objects.	
	onc_fn_fs	A context for naming files, directories, and file systems.	
	onc_fn_printername	A context for naming printers.	
	onc_printers	A printer object. When implemented on top of NIS+, this could also be a context for naming printers.	
	fn_link_ref	An XFN link.	
	inet_domain	An Internet domain.	
Address Types	All address types use the FN_ID_STRING identifier format unless otherwise qualified. The format of address contents is determined by the corresponding address type.		
	onc_fn_nisplus	For an FNS enterprise-level object implemented on top of NIS+. The address	

Standards, Environments, and Macros 85

fns_references(5)

	contains the context type, context representation type (either normal or merged), version number of the reference, and the NIS+ name of the object. The only intended use of this reference is that it be passed to fn_ctx_handle_from_ref(3XFN)
onc_fn_nis	For an FNS enterprise-level object implemented on top of NIS. The address contains the context type and version number of the reference, and the NIS name of the object. The only intended use of this reference is that it be passed to fn_ctx_handle_from_ref(3XFN).
onc_fn_files	For an FNS enterprise-level object implemented on top of /etc files. The address contains the context type and version number of the reference, and the location of the object in the /etc file system. The only intended use of this reference is that it be passed to fn_ctx_handle_from_ref(3XFN).
onc_fn_fs_user	For a user's home directory. The address contains the user's name and the name of the naming service password table where the user's home directory is stored.
onc_fn_fs_user_nisplus	For a user's home directory. The address contains the user's name and the name of the NIS+ password table where the user's home directory is stored.
onc_fn_fs_host	For all file systems exported by a host. The address contains the host's name.
onc_fn_fs_mount	For a single mount point. The address contains the mount options, the name of the servers and the exported path. See mount(1M).
onc_fn_printer_files	For a printer's address in the files naming service.
onc_fn_printer_nis	For a printer's address in the NIS naming service.
onc_fn_printer_nisplus	For a printer's address in the NIS+ naming service.

		fns_references(5)
	fn_link_addr	For an XFN link address. The contents is the string form of the composite name.
	inet_domain	For an Internet domain. The address contains the fully-qualified domain name (for example, "Wiz.COM.")
	inet_ipaddr_string	For an object with an Internet address. The address contains an internet IP address in dotted string form (for example, "192.144.2.3").
	x500	For an X.500 object. The address contains an X.500 Distinguished Name, in the syntax specified in the X/Open DCE: Directory Services.
	osi_paddr	For an object with an OSI presentation address. The address contains the string encoding of an OSI Presentation Address as defined in <i>A string encoding of Presentation</i> <i>Address</i> (RFC 1278).
	onc_printers_bsaddr	For a printer that understands the BSD print protocol. The address contains the machine name and printer name used by the protocol.
	onc_printers_use	For a printer alias. The address contains a printer name.
	onc_printers_all	For a list of printers that are enumerated using the "all" option. The address contains a list of printer names.
	onc_printers_location	For a printer's location. The address format is unspecified.
	onc_printers_type	For a printer's type. The address format is unspecified.
	onc_printers_speed	For a printer's speed. The address format is unspecified.
ALSO	<pre>mount(1M), fn_ctx_handle_from_re fns_policies(5)</pre>	f(3XFN), xfn(3XFN), fns(5),
	Hardcastle-Kille, S.E., <i>A string encoding of</i> College London, November 1991.	Presentation Address, RFC 1278, University

SEE

Standards, Environments, and Macros 87

fns_x500(5)

NAME | fns_x500 – overview of FNS over X.500 implementation

DESCRIPTION

N Federated Naming Service (FNS) provides a method for federating multiple naming services under a single, simple interface for the basic naming operations. One of the naming services supported by FNS is the X.500 Directory Service (see ITU-T X.500 or ISO/IEC 9594). X.500 is a global directory service. Its components cooperate to manage information about a hierarchy of objects on a worldwide scope. Such objects include countries, organizations, people, services, and machines. FNS uses X.500 to name entities globally.

FNS provides the XFN interface for retrieval and modification of information stored in X.500. In addition, enterprise namespaces such as those served by NIS+ and NIS can be federated with X.500 by adding reference information to X.500 describing how to reach the desired next naming service. To federate a NIS+ or NIS namespace under X.500, perform the following steps:

1. Obtain the root reference for the NIS+ hierarchy or NIS domain.

- 2. Enhance the X.500 schema to support the addition of XFN references.
- 3. Create an X.500 entry to store the XFN reference.
- 4. Add the XFN reference.

The root reference is referred to as the *next naming system reference* because it refers to the *next* naming system beneath X.500. This reference contains information about how to communicate with the NIS+ or NIS servers and has the following format:

<domainname> <server name> [<server address>]

where *<domainname>* is the fully qualified domain name. Notice that NIS+ and NIS have slightly different syntaxes for domain names. For NIS+, the fully qualified domain name is case-insensitive and terminated by a dot character ('.'). For NIS, the fully qualified domain name is case-sensitive and *not* terminated by a dot character. For both NIS+ and NIS, *<server address>* is optional. If it is not supplied, a host name lookup will be performed to get the machine's address.

For example, if the machine wiz-nisplus-server with address 133.33.33 serves the NIS+ domain wiz.com., the reference would look like this:

wiz.com.wiz-nisplus-server133.33.33.33

For another example, if the machine woz-nis-server serves the NIS domain Woz.COM, the reference would look like this:

Woz.COM woz-nis-server

Before the next naming system reference can be added to X.500, the X.500 schema must be altered to include the following object class and associated attributes (defined in ASN.1 notation).

88 man pages section 5: Standards, Environments, and Macros • Last Revised 29 Jan 1998

fns_x500(5)

```
xFNSupplement OBJECT-CLASS ::= {
 SUBCLASS OF { top }
 KIND
              auxiliary
 MAY CONTAIN { objectReferenceString | nNSReference-
String }
 ID
             id-oc-xFNSupplement }
id-oc-xFNSupplement OBJECT IDENTIFIER ::= {
 iso member-body(2) ansi(840) sun(113536) 25 }
objectReferenceString ATTRIBUTE ::= {
 WITH SYNTAX
               OCTET STRING
 EQUALITY MATCHING RULE octetStringMatch
 SINGLE VALUE TRUE
                     id-at-objectReferenceString }
 ΤD
id-at-objectReferenceString OBJECT IDENTIFIER ::= {
 iso member-body(2) ansi(840) sun(113536) 30 }
nNSReferenceString ATTRIBUTE ::= {
 WITH SYNTAX OCTET STRING
 EQUALITY MATCHING RULE octetStringMatch
 SINGLE VALUE TRUE
 ТD
                     id-at-nNSReferenceString }
id-at-nNSReferenceString OBJECT IDENTIFIER ::= {
 so member-body(2) ansi(840) sun(113536) 31 }
```

The procedures for altering the X.500 schema will vary from implementation to implementation. Consult *Solstice X.500* or the schema administration guide for your X.500 product.

Once X.500 supports XFN references, the next naming system reference can be added by first creating an X.500 object and then adding the new reference to it. For example, the following commands create entries for the Wiz and Woz organizations in the U.S.A. and add the reference information shown in the examples above to them.

For NIS+:

```
example% fnattr .../c=us/o=wiz -a objectclass \
    top organization xfnsupplement
example% fnbind -r .../c=us/o=wiz/ onc_fn_enterprise \
    onc_fn_nisplus_root "wiz.com. wiz-nisplus-server"
For NIS:
example% fnattr .../c=us/o=woz -a objectclass \
    top organization xfnsupplement
example% fnbind -r .../c=us/o=woz/ onc_fn_enterprise \
    onc_fn_nis_root "Woz.COM woz-nis-server"
Notice the mandatory trailing slash ('/') in the name argument to fnbind(1).
```

fns_	x500	(5)
------	------	-----

This modification effectively adds the next naming system reference to X.500. The reference may be retrieved using fnlookup(1) to see if the information has been added properly. For example, the following command looks up the next naming system reference of the Wiz organization:

example% fnlookup -v .../c=us/o=wiz/

Note the mandatory trailing slash.

After this administrative step has been taken, clients outside of the NIS+ hierarchy or NIS domain can access and perform operations on the contexts in the NIS+ hierarchy or NIS domain. Foreign NIS+ clients access the hierarchy as unauthenticated NIS+ clients. Continuing the example above, and assuming that NIS+ is federated underneath the Wiz organization, the root of the NIS+ enterprise may be listed using the command:

example% fnlist .../c=us/o=wiz/

Note the mandatory trailing slash.

The next naming system reference may be removed using the command:

example% fnunbind .../c=us/o=wiz/

Note the mandatory trailing slash.

SEE ALSO fnattr(1), fnbind(1), fnlist(1), fnlookup(1), nis+(1), ypserv(1M), xfn(3XFN), fns(5), fns_dns(5), fns_nis(5), fns_nis+(5), fns_references(5)

Solstice X.500

NOTES In a 64-bit XFN application, retrieval and modification of information stored in the X.500 directory service is not supported.

90 man pages section 5: Standards, Environments, and Macros • Last Revised 29 Jan 1998

NAME	formats – file format notation			
DESCRIPTION	Utility descriptions use a syntax to describe the data organization within files—stdin, stdout, stderr, input files, and output files—when that organization is not otherwise obvious. The syntax is similar to that used by the printf(3C) function. When used for stdin or input file descriptions, this syntax describes the format that could have been used to write the text to be read, not a format that could be used by the scanf(3C) function to read the input file.			
Format	The description of	an individual	record is as follows:	
	" <format>", [<arg1< th=""><th>>, <arg2>,</arg2></th><th>., <<i>argn></i>]</th></arg1<></format>	>, <arg2>,</arg2>	., < <i>argn></i>]	
	The format is a cl	haracter string	; that contains three types of objects defined below:	
	characters	S	Characters that are not <i>escape sequences</i> or <i>conversion pecifications</i> , as described below, are copied to the output.	
	escape sequences	F	Represent non-graphic characters.	
	conversion specificat		pecifies the output format of each argument. (See elow.)	
	The following cha	racters have tl	ne following special meaning in the format string:	
	"	(An empty character position.) One or more blank characters. Exactly one space character.		
	\wedge			
	empty character po output (not <i>white s</i> utility that reads th	osition in for <i>pace,</i> which ca nat output as :), and so forth	ome flexibility for application output. Note that an mat represents one or more blank characters on the in include newline characters). Therefore, another its input must be prepared to parse the data using . The character is used when exactly one space	
Escape Sequences	The following table lists escape sequences and associated actions on display devices capable of the action.			
	Sequence	Character	Terminal Action	
	\\	backslash	None.	
	\a	alert	Attempts to alert the user through audible or visible notification.	
	\b 	backspace	Moves the printing position to one column before the current position, unless the current position is the start of a line.	

formats(5)

	Sequence	Character	Terminal Action
	\f	form-feed	Moves the printing position to the initial printing position of the next logical page.
	\n	newline	Moves the printing position to the start of the next line.
	\r	carriage-return	Moves the printing position to the start of the current line.
	\t	tab	Moves the printing position to the next tab position on the current line. If there are no more tab positions left on the line, the behavior is undefined.
	\v	vertical-tab	Moves the printing position to the start of the next vertical tab position. If there are no more vertical tab positions left on the page, the behavior is undefined.
Conversion Specifications		specification is in following appear	troduced by the percent-sign character (%). After the in sequence:
	flags		o or more <i>flags,</i> in any order, that modify the aning of the conversion specification.
	field width	mir con is p flag	optional string of decimal digits to specify a nimum <i>field width</i> . For an output field, if the verted value has fewer bytes than the field width, it added on the left (or right, if the left-adjustment g (–), described below, has been given to the field lth).
	precision	d, c leac rad ma: con wri take	tes the minimum number of digits to appear for the b, i, u, x or X conversions (the field is padded with ding zeros), the number of digits to appear after the ix character for the e and f conversions, the ximum number of significant digits for the g version; or the maximum number of bytes to be tten from a string in s conversion. The precision es the form of a period (.) followed by a decimal it string; a null digit string is treated as zero.
	conversion charact		onversion character (see below) that indicates the e of conversion to be applied.
flags	The <i>flags</i> and the	ir meanings are:	
	-	The result of th	e conversion is left-justified within the field.
	+	The result of a –).	signed conversion always begins with a sign (+ or

	<space></space>	If the first character of a signed conversion is not a sign, a space character is prefixed to the result. This means that if the space character and + flags both appear, the space character flag is ignored.
	#	The value is to be converted to an alternative form. For c, d, i, u, and s conversions, the behaviour is undefined. For o conversion, it increases the precision to force the first digit of the result to be a zero. For x or X conversion, a non-zero result has 0x or 0X prefixed to it, respectively. For e, E, f, g, and G conversions, the result always contains a radix character, even if no digits follow the radix character. For g and G conversions, trailing zeros are not removed from the result as they usually are.
	0	For d, i, o, u, x, X, e, E, f, g, and G conversions, leading zeros (following any indication of sign or base) are used to pad to the field width; no space padding is performed. If the 0 and – flags both appear, the 0 flag is ignored. For d, i, o, u, x and X conversions, if a precision is specified, the 0 flag is ignored. For other conversions, the behaviour is undefined.
Conversion Characters	undefined if there while arguments re	naracter results in fetching zero or more arguments. The results are are insufficient arguments for the format. If the format is exhausted emain, the excess arguments are ignored.
	The conversion char	acters and their meanings are:
	d,i,o,u,x,X	The integer argument is written as signed decimal (d or i), unsigned octal (o), unsigned decimal (u), or unsigned hexadecimal notation (x and X). The d and i specifiers convert to signed decimal in the style [–] <i>ddd.</i> The x conversion uses the numbers and letters 0123456789abcdef and the X conversion uses the numbers and letters 0123456789ABCDEF. The <i>precision</i> component of the argument specifies the minimum number of digits to appear. If the value being converted can be represented in fewer digits than the specified minimum, it is expanded with leading zeros. The default precision is 1. The result of converting a zero value with a precision of 0 is no characters. If both the field width and precision are omitted, the implementation may precede, follow or precede and follow numeric arguments of types d, i and u with blank characters; arguments of type o (octal) may be preceded with leading zeros.
		The treatment of integers and spaces is different from the printf(3C) function in that they can be surrounded with blank characters. This was done so that, given a format such as:
		"%d\n",< <i>foo</i> >
		the implementation could use a printf() call such as:

formats(5)

	<pre>printf("%6d\n", foo);</pre>
	and still conform. This notation is thus somewhat like $\verb+scanf()$ addition to <code>printf()</code> .
f	The floating point number argument is written in decimal notation in the style [-] <i>ddd.ddd</i> , where the number of digits after the radii character (shown here as a decimal point) is equal to the <i>precision</i> specification. The LC_NUMERIC locale category determines the radix character to use in this format. If the <i>precision</i> is omitted from the argument, six digits are written after the radix character; if the <i>precision</i> is explicitly 0, no radix character appears.
e,E	The floating point number argument is written in the style $[-] d.ddde\pm dd$ (the symbol \pm indicates either a plus or minus sign where there is one digit before the radix character (shown here a decimal point) and the number of digits after it is equal to the precision. The LC_NUMERIC locale category determines the radix character to use in this format. When the precision is missing, six digits are written after the radix character; if the precision is 0, no radix character appears. The E conversion character produces a number with E instead of e introducing the exponent. The exponent always contains at least two digits. However, if the val to be written requires an exponent greater than two digits, additional exponent digits are written as necessary.
g,G	The floating point number argument is written in style f or e (or style E in the case of a G conversion character), with the precision specifying the number of significant digits. The style used dependent on the value converted: style g is used only if the exponent resulting from the conversion is less than -4 or greater than or equal to the precision. Trailing zeros are removed from the result A radix character appears only if it is followed by a digit.
С	The integer argument is converted to an unsigned char and th resulting byte is written.
S	The argument is taken to be a string and bytes from the string ar written until the end of the string or the number of bytes indicate by the <i>precision</i> specification of the argument is reached. If the precision is omitted from the argument, it is taken to be infinite, all bytes up to the end of the string are written.
%	Write a % character; no argument is converted.
the result of a contain the cor	a non-existent or insufficient <i>field width</i> cause truncation of a field; if conversion is wider than the field width, the field is simply expanded version result. The term <i>field width</i> should not be confused with the used in the description of %s.

	One difference from the C function printf() is that the l and h conversion characters are not used. There is no differentiation between decimal values for type int, type long, or type short. The specifications %d or %i should be interpreted as an arbitrary length sequence of digits. Also, no distinction is made between single precision and double precision numbers (float or double in C). These are simply referred to as floating point numbers.
	Many of the output descriptions use the term line, such as:
	"%s", <input line=""/>
	Since the definition of line includes the trailing newline character already, there is no need to include a \n in the format; a double newline character would otherwise result.
EXAMPLES	EXAMPLE 1 To represent the output of a program that prints a date and time in the form Sunday, July 3, 10:02, where <i><weekday></weekday></i> and <i><month></month></i> are strings:
	"%s,/\%s/\%d,/\%d:%.2d\n", <weekday>,<month>,<day>,<hour>,<min></min></hour></day></month></weekday>
	EXAMPLE 2 To show pi written to 5 decimal places:
	"pi/\=/\%.5f\n", <value of="" pi=""></value>
	EXAMPLE 3 To show an input file format consisting of five colon-separated fields: "%s:%s:%s:%s:%s\n", <arg1>, <arg2>, <arg3>, <arg4>, <arg5></arg5></arg4></arg3></arg2></arg1>
SEE ALSO	<pre>awk(1), printf(1), printf(3C), scanf(3C)</pre>

NAME | fsattr – extended file attributes

DESCRIPTION

Attributes are logically supported as files within the file system. The file system is therefore augmented with an orthogonal name space of file attributes. Any file (including attribute files) can have an arbitrarily deep attribute tree associated with it. Attribute values are accessed by file descriptors obtained through a special attribute interface. This logical view of "attributes as files" allows the leveraging of existing file system interface functionality to support the construction, deletion, and manipulation of attributes.

The special files "." and ".." retain their accustomed semantics within the attribute hierarchy. The "." attribute file refers to the current directory and the ".." attribute file refers to the parent directory. The unnamed directory at the head of each attribute tree is considered the "child" of the file it is associated with and the ".." file refers to the associated file. For any non-directory file with attributes, the ".." entry in the unnamed directory refers to a file that is not a directory.

Conceptually, the attribute model is fully general. Extended attributes can be any type of file (doors, links, directories, and so forth) and can even have their own attributes (fully recursive). As a result, the attributes associated with a file could be an arbitrarily deep directory hierarchy where each attribute could have an equally complex attribute tree associated with it. Not all implementations are able to, or want to, support the full model. Implementation are therefore permitted to reject operations that are not supported. For example, the implementation for the UFS file system allows only regular files as attributes (for example, no sub-directories) and rejects attempts to place attributes.

The following list details the operations that are rejected in the current implementation:

link	Any attempt to create links between attribute and non-attribute space is rejected to prevent security-related or otherwise sensitive attributes from being exposed, and therefore manipulable, as regular files.
rename	Any attempt to rename between attribute and non-attribute space is rejected to prevent an already linked file from being renamed and thereby circumventing the link restriction above.
mkdir, symlink, mknod	Any attempt to create a "non-regular" file in attribute space is rejected to reduce the functionality, and therefore exposure and risk, of the initial implementation.

	The entire available name space has been allocated to "general use" to bring the implementation in line with the NFSv4 draft standard [NFSv4]. That standard defines "named attributes" (equivalent to Solaris Extended Attributes) with no naming restrictions. All Sun applications making use of opaque extended attributes will use the prefix "SUNW".
Shell-level API	The command interface for extended attributes is the set of applications provided by Solaris for the manipulation of attributes from the command line. This interface consists of a set of existing utilities that have been extended to be "attribute-aware", plus the runat utility designed to "expose" the extended attribute space so that extended attributes can be manipulated as regular files.
	The -@ option enable utilities to manipulate extended attributes. As a rule, this option enables the utility to enter into attribute space when the utility is performing a recursive traversal of file system space. This is a fully recursive concept. If the underlying file system supports recursive attributes and directory structures, the -@ option opens these spaces to the file tree-walking algorithms.
	The following utilities accommodate extended attributes (see the individual manual pages for details):
	cp By default, cp ignores attributes and copies only file data. This is intended to maintain the semantics implied by cp currently, where attributes (such as

The following utilities accomm pages for details):	nodate extended attributes (see the individual manual
ср	By default, cp ignores attributes and copies only file data. This is intended to maintain the semantics implied by cp currently, where attributes (such as owner and mode) are not copied unless the -p option is specified. With the -@ (or -p) option, cp attempts to copy all attributes along with the file data.
cpio	The -@ option informs cpio to archive attributes, but by default cpio ignores extended attributes. See Extended Archive Formats below for a description of the new archive records.
du	File sizes computed include the space allocated for any extended attributes present.
find	By default, find ignores attributes. The -xattr expression provides support for searches involving attribute space. It returns true if extended attributes are present on the current file.
fsck	The fsck utility manages extended attribute data on the disk. A file system with extended attributes can be mounted on versions of Solaris that are not attribute-aware (versions prior to Solaris 9), but the attributes will not be accessible and fsck will strip them from the files and place them in lost+found. Once the attributes have been stripped the file system is completely stable on Solaris versions that are not attribute-aware, but would now be considered

	corrupted on attribute-aware versions of Solaris. The attribute-aware fsck utility should be run to stabilize the file system before using it in an attribute-aware environment.
fsdb	This fsdb utility is able to find the inode for the "hidden" extended attribute directory.
ls	The ls -@ command displays an "@" following the mode information when extended attributes are present. More precisely, the output line for a given file contains an "@" character following the mode characters if the pathconf(2) variable XATTR_EXISTS is set to true. See the pathconf() section below. The -@ option uses the same general output format as the -l option.
mv	When a file is moved, all attributes are carried along with the file rename. When a file is moved across a file system boundary, the copy command invoked is similar to the cp -p variant described above and extended attributes are "moved". If the extended file attributes cannot be replicated, the move operation fails and the source file is not removed.
pax	The -@ option informs pax to archive attributes, but by default pax ignores extended attributes. The pax(1) utility is a generic replacement for both tar(1) and cpio(1) and is able to produce either output format in its archive. See Extended Archive Formats below for a description of the new archive records.
tar	In the default case, tar does not attempt to place attributes in the archive. If the -@ option is specified, however, tar traverses into the attribute space of all files being placed in the archive and attempts to add the attributes to the archive. A new record type has been introduced for extended attribute entries in tar archive files (the same is true for pax and cpio archives) similar to the way ACLs records were defined. See Extended Archive Formats below for a description of the new archive records.
modified in a manner similar performing chmod on a file w	mod, chown, chgrp) that one might expect to be to those listed above. For example, one might expect that ould not only affect the file itself but would also affect at rectory if not any existing extended attribute files. This is

modified in a manner similar to those listed above. For example, one might expect that performing chmod on a file would not only affect the file itself but would also affect at least the extended attribute directory if not any existing extended attribute files. This is not the case. The model chosen for extended attributes implies that the attribute directory and the attributes themselves are all file objects in their own right, and can therefore have independent file status attributes associated with them (a given

implementation cannot support this, for example, for intrinsic attributes). The relationship is left undefined and a fine-grained control mechanism (runat(1)) is provided to allow manipulation of extended attribute status attributes as necessary. The runat utility has the following syntax: runat filename [command] The runat utility executes the supplied command in the context of the "attribute space" associated with the indicated file. If no command argument is supplied, a shell is invoked. See runat(1) for details. **Application-level** The primary interface required to access extended attributes at the programmatic level API is the openat(2) function. Once a file descriptor has been obtained for an attribute file by an openat () call, all normal file system semantics apply. There is no attempt to place special semantics on read(2), write(2), ftruncate(3C), or other functions when applied to attribute file descriptors relative to "normal" file descriptors. The set of existing attributes can be browsed by calling openat () with "." as the file name and the O XATTR flag set, resulting in a file descriptor for the attribute directory. The list of attributes is obtained by calls to getdents(2) on the returned file descriptor. If the target file did not previously have any attributes associated with it, an empty top-level attribute directory is created for the file and subsequent getdents () calls will return only "." and "..". While the owner of the parent file owns the extended attribute directory, it is not charged against its quota if the directory is empty. Attribute files themselves, however, are charged against the user quota as any other regular file. Additional system calls have been provided as convenience functions. These include the fchownat(2), fstatat(2), futimesat(2), renameat(2), unlinkat(2). These new functions, along with openat(), provide a mechanism to access files relative to an arbitrary point in the file system, rather than only the current working directory. This mechanism is particularly useful in situations when a file descriptor is available with no path. The openat () function, in particular, can be used in many contexts where chdir() or fchdir() is currently required. See chdir(2). Open a file relative to a file descriptor int openat (int fd, const char *path, int oflag [, mode t mode]) The openat(2) function behaves exactly as open(2) except when given a relative path. Where open () resolves a relative path from the current working directory, openat () resolves the path based on the vnode indicated by the supplied file descriptor. When oflag is O XATTR, openat() interprets the *path* argument as an extended attribute reference. The following code fragment uses openat() to examine the attributes of some already opened file: dfd = openat(fd, ".", O RDONLY | O XATTR); (void)getdents(dfd, buf, nbytes);

If openat () is passed the special value AT_FDCWD as its first (*fd*) argument, its behavior is identical to open() and the relative path arguments are interpreted relative to the current working directory. If the O_XATTR flag is provided to openat() or to open(), the supplied path is interpreted as a reference to an extended attribute on the current working directory.

Unlink a file relative to a directory file descriptor

int unlinkat (int *dirfd*, const char *pathflag, int flagflag)

The unlinkat(2) function deletes an entry from a directory. The *path* argument indicates the name of the entry to remove. If *path* an absolute path, the *dirfd* argument is ignored. If it is a relative path, it is interpreted relative to the directory indicated by the *dirfd* argument. If *dirfd* does not refer to a valid directory, the function returns ENOTDIR. If the special value AT_FDCWD is specified for *dirfd*, a relative path argument is resolved relative to the current working directory. If the *flag* argument is 0, all other semantics of this function are equivalent to unlink(2). If *flag* is set to AT_REMOVEDIR, all other semantics of this function are equivalent to rmdir(2).

Rename a file relative to directories

int renameat (int *fromfd*, const char *old, int tofd, const char *new)

The renameat(2) function renames an entry in a directory, possibly moving the entry into a different directory. The *old* argument indicates the name of the entry to rename. If this argument is a relative path, it is interpreted relative to the directory indicated by the *fd* argument. If it is an absolute path, the *fromfd* argument is ignored. The *new* argument indicates the new name for the entry. If this argument is a relative path, it is interpreted by the *tofd* argument. If it is an absolute path, the *tofd* argument. If it is an absolute path, the *tofd* argument is a relative path, it is interpreted relative to the directory indicated by the *tofd* argument. If it is an absolute path, the *tofd* argument is ignored.

In the relative path cases, if the directory file descriptor arguments do not refer to a valid directory, the function returns ENOTDIR. All other semantics of this function are equivalent to rename(2).

If a special value AT_FDCWD is specified for either the *fromfd* or *tofd* arguments, their associated path arguments (*old* and *new*) are interpreted relative to the current working directory if they are not specified as absolute paths. Any attempt to use renameat () to move a file that is not an extended attribute into an extended attribute directory (so that it becomes an extended attribute) will fail. The same is true for an attempt to move a file that is an extended attribute into a directory that is not an extended attribute into a directory that is not an extended attribute directory.

Obtain information about a file

int fstatat (int fd, const char *path, struct stat* buf, int flag)

The fstatat(2) function obtains information about a file. If the *path* argument is relative, it is resolved relative to the *fd* argument file descriptor, otherwise the *fd* argument is ignored. If the *fd* argument is a special value AT_FDCWD the path is resolved relative to the current working directory. If the *path* argument is a null

pointer, the function returns information about the file referenced by the *fd* argument. In all other relative path cases, if the *fd* argument does not refer to a valid directory, the function returns ENOTDIR. If the *flag* argument is set to AT_SYMLINK_NOFOLLOW, the function will not automatically traverse a symbolic link at the position of the path. The fstatat() function is a multi-purpose function that can be used in place of stat(), lstat(), or fstat(). See stat(2).

The function call stat (*path*, *buf*) is identical to fstatat (AT_FDCWD, *path*, *buf*, 0).

The function call lstat (*path*, *buf*) is identical to fstatat (AT_FDCWD, *path*, *buf*, AT_SYMLINK_NOFOLLOW)

The function call fstat (fildes, buf) is identical to fstatat (fildes, NULL, buf, 0).

Set owner and group ID

int fchownat (int fd, const char *path, uid_t owner, gid_t group, int flag)

The fchownat(2) function sets the owner ID and group ID for a file. If the *path* argument is relative, it is resolved relative to the *fd* argument file descriptor, otherwise the *fd* argument is ignored. If the *fd* argument is a special value AT_FDCWD the path is resolved relative to the current working directory. If the path argument is a null pointer, the function sets the owner and group ID of the file referenced by the *fd* argument. In all other relative path cases, if the *fd* argument does not refer to a valid directory, the function returns ENOTDIR. If the *flag* argument is set to AT_SYMLINK_NOFOLLOW, the function will not automatically traverse a symbolic link at the position of the path. The fchownat() function is a multi-purpose function that can be used in place of chown(), lchown(), or fchown(). See chown(2).

The function call chown (*path*, *owner*, *group*) is equivalent to fchownat (AT_FDCWD, *path*, *owner*, *group*, 0).

The function call lchown (*path*, *owner*, *group*) is equivalent to fchownat (AT FDCWD, *path*, *owner*, *group*, AT SYMLINK NOFOLLOW).

Set file access and modification times

int futimesat (int fd, const char *path, const struct timeval times[2])

The futimesat(2) function sets the access and modification times for a file. If the *path* argument is relative, it is resolved relative to the *fd* argument file descriptor; otherwise the *fd* argument is ignored. If the *fd* argument is the special value AT_FDCWD , the path is resolved relative to the current working directory. If the *path* argument is a null pointer, the function sets the access and modification times of the file referenced by the *fd* argument. In all other relative path cases, if the *fd* argument does not refer to a valid directory, the function returns ENOTDIR. The futimesat() function can be used in place of utimes(2).

The function call utimes (*path*, *times*) is equivalent to futimesat (AT_FDCWD, *path*, *times*).

New pathconf() functionality

long int pathconf(const char *path, int name)

Two variables have been added to pathconf(2) to provide enhanced support for extended attribute manipulation. The XATTR_ENABLED variable allows an application to determine if attribute support is currently enabled for the file in question. The XATTR_EXISTS variable allows an application to determine whether there are any extended attributes associated with the supplied path.

Open/Create an attribute file

int attropen (const char *path, const char *attrpath, int oflag [, mode_t mode])

The attropen(3C) function returns a file descriptor for the named attribute, *attrpath*, of the file indicated by *path*. The *oflag* and *mode* arguments are identical to the open(2) arguments and are applied to the open operation on the attribute file (for example, using the O_CREAT flag creates a new attribute). Once opened, all normal file system operations can be used on the attribute file descriptor. The attropen() function is a convenience function and is equivalent to the following sequence of operations:

```
fd = open (path, O_RDONLY);
attrfd = openat(fd, attrpath, oflag|O_XATTR, mode);
close(fd);
```

The set of existing attributes can be browsed by calling attropen() with "." as the attribute name. The list of attributes is obtained by calling getdents(2) (or fdopendir(3C) followed by readdir(3C), see below) on the returned file descriptor.

Convert an open file descriptor for a directory into a directory descriptor

DIR * fdopendir (const int fd)

The fdopendir(3C) function promotes a file descriptor for a directory to a directory pointer suitable for use with the readdir(3C) function. The originating file descriptor should not be used again following the call to fdopendir(). The directory pointer should be closed with a call to closedir(3C). If the provided file descriptor does not reference a directory, the function returns ENOTDIR. This function is useful in circumstances where the only available handle on a directory is a file descriptor. See attropen(3C) and openat(2).

Using the API

The following examples demonstrate how the API might be used to perform basic operations on extended attributes:

EXAMPLE 1 List extended attributes on a file.

```
attrdirfd = attropen("test", ".", O_RDONLY);
dirp = fdopendir(attrdirfd);
while (dp = readdir(dirp)) {
...
```

```
EXAMPLE 2 Open an extended attribute.
                    attrfd = attropen("test", dp->d name, O RDONLY);
                    or
                    attrfd = openat(attrdirfd, dp->d_name, O_RDONLY);
                    EXAMPLE 3 Read from an extended attribute.
                    while (read(attrfd, buf, 512) > 0) {
                     . . .
                    EXAMPLE 4 Create an extended attribute.
                    newfd = attropen("test", "attr", O_CREAT|O_RDWR);
                    or
                    newfd = openat(attrdirfd, "attr", O_CREAT|O_RDWR);
                    EXAMPLE 5 Write to an extended attribute.
                    count = write(newfd, buf, length);
                    EXAMPLE 6 Delete an extended attribute.
                    error = unlinkat(attrdirfd, "attr");
                    Applications intending to access the interfaces defined here as well as the POSIX and
                    X/Open specification-conforming interfaces should define the macro
                     ATFILE SOURCE to be 1 and set whichever feature test macros are appropriate to
                    obtain the desired environment. See standards(5).
Extended Archive
                    As noted above in the description of command utilities modified to provide support
         Formats
                    for extended attributes, the archive formats for tar(1) and cpio(1) have been
                    extended to provide support for archiving extended attributes. This section describes
                    the specifics of the archive format extensions.
                    Extended tar format
                    The tar archive is made up of a series of 512 byte blocks. Each archived file is
                    represented by a header block and zero or more data blocks containing the file
                    contents. The header block is structured as shown in the following table.
                     Field Name
                                                  Length (in Octets)
                                                                               Description
                                                  100
                     Name
                                                                               File name string
                     Mode
                                                  8
                                                                               12 file mode bits
```

Standards, Environments, and Macros 103

Field Name	Length (in Octets)	Description
Uid	8	User ID of file owner
Gid	8	Group ID of file owner
Size	12	Size of file
Mtime	12	File modification time
Chksum	8	File contents checksum
Typeflag	1	File type flag
Linkname	100	Link target name if file linked
Magic	6	"ustar"
Version	2	"00"
Uname	32	User name of file owner
Gname	32	Group name of file owner
Devmajor	8	Major device ID if special file
Devminor	8	Minor device ID if special file
Prefix	155	Path prefix string for file

The extended attribute project extends the above header format by defining a new header type (for the Typeflag field). The type 'E' is defined to be used for all extended attribute files. Attribute files are stored in the tar archive as a sequence of two <header ,data> pairs. The first file contains the data necessary to locate and name the extended attribute in the file system. The second file contains the actual attribute file data. Both files use an 'E' type header. The prefix and name fields in extended attribute headers are ignored, though they should be set to meaningful values for the benefit of archivers that do not process these headers. Solaris archivers set the prefix field to "/dev/null" to prevent archivers that do not understand the type 'E' header from trying to restore extended attribute files in inappropriate places.

Extended cpio format

The cpio archive format is octet-oriented rather than block-oriented. Each file entry in the archive includes a header that describes the file, followed by the file name, followed by the contents of the file. These data are arranged as described in the following table.

Field Name	Length (in Octets)	Description
c_magic	6	70707
c_dev	6	First half of unique file ID

Field Name	Length (in Octets)	Description
c_ino	6	Second half of unique file ID
c_mode	6	File mode bits
c_uid	6	User ID of file owner
c_gid	6	Group ID of file owner
c_nlink	6	Number of links referencing file
c_rdev	6	Information for special files
c_mtime	11	Modification time of file
c_namesize	6	Length of file pathname
c_filesize	11	Length of file content
c_name	c_namesize	File pathname
c_filedata	c_filesize	File content

The basic archive file structure is not changed for extended attributes. The file type bits stored in the c_mode field for an attribute file are set to 0xB000. As with the tar archive format, extended attributes are stored in cpio archives as two consecutive file entries. The first file describes the location/name for the extended attribute. The second file contains the actual attribute file content. The c_name field in extended attribute headers is ignored, though it should be set to a meaningful value for the benefit of archivers that do not process these headers. Solaris archivers start the pathname with "/dev/null/" to prevent archivers that do not understand the type 'E' header from trying to restore extended attribute files in inappropriate places.

Attribute identification data format

Both the tar and cpio archive formats can contain the special files described above, always paired with the extended attribute data record, for identifying the precise location of the extended attribute. These special data files are necessary because there is no simple naming mechanism for extended attribute files. Extended attributes are not visible in the file system name space. The extended attribute name space must be "tunneled into" using the openat() function. The attribute identification data must support not only the flat naming structure for extended attributes, but also the possibility of future extensions allowing for attribute directory hierarchies and recursive attributes. The data file is therefore composed of a sequence of records. It begins with a fixed length header describing the content. The following table describes the format of this data file.

Field Name	Length (in Octets)	Description
h_version	7	Name file version
h_size	10	Length of data file
h_component_len	10	Total length of all path segments
h_link_comp_len	10	Total length of all link segments
path	h_component_len	Complex path
link_path	h_link_comp_len	Complex link path

As demonstrated above, the header is followed by a record describing the "path" to the attribute file. This path is composed of two or more path segments separated by a null character. Each segment describes a path rooted at the hidden extended attribute directory of the leaf file of the previous segment, making it possible to name attributes on attributes. The first segment is always the path to the parent file that roots the entire sequence in the normal name space. The following table describes the format of each segment.

Field Name	Length (in Octets)	Description
h_namesz	7	Length of segment path
h_typeflag	1	Actual file type of attribute file
h_names	h_namesz	Parent path + segment path

If the attribute file is linked to another file, the path record is followed by a second record describing the location of the referencing file. The structure of this record is identical to the record described above.

106 man pages section 5: Standards, Environments, and Macros • Last Revised 1 Aug 2001

NAME | iconv_1250 – code set conversion tables for MS 1250 (Windows Latin 2)

DESCRIPTION

The following code set conversions are supported:

	Code Set Conversions Supported			
Code	Symbol	Target Code	Symbol	Target Output
MS 1250	win2	ISO 8859-2	iso2	ISO Latin 2
MS 1250	win2	MS 852	dos2	MS-DOS Latin 2
MS 1250	win2	Mazovia	maz	Mazovia
MS 1250	win2	DHN	dhn	Dom Handlowy Nauki

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

MS 1250 to ISO 8859-2 For the conversion of MS 1250 to ISO 8859-2, all characters not in the following table are mapped unchanged.

Conversions Performed			
MS 1250	ISO 8859-2	MS 1250	ISO 8859-2
24-211	40	235	273
212	251	236	276
213	40	237	274
214	246	241	267
215	253	245	241
216	256	246-267	40
217	254	271	261
221-231	40	273	40
232	271	274	245
233	40	276	265
234	266	247	365

MS 1250 to MS 852 For the conversion of MS 1250 to MS 852, all characters not in the following table are mapped unchanged.

iconv_1250(5)

MS 1250	MS 852	MS 1250	MS 852
200-211	40	311	220
212	346	312	250
213	40	313	323
214	227	314	267
215	233	315	326
216	246	316	327
217	215	317	322
220-231	40	320	321
232	347	321	343
233	40	322	325
234	230	323	340
235	234	324	342
236	247	325	212
237	253	326	231
240	377	327	236
241	363	330	374
242	364	331	336
243	235	332	351
244	317	333	353
245	244	334	232
246	40	335	355
247	365	336	335
250	371	337	341
251	40	340	352
252	270	341	240
253	256	342	203
254	252	343	307
255	360	344	204

108 man pages section 5: Standards, Environments, and Macros • Last Revised 18 Apr 1997

MS 1250	MS 852	MS 1250	MS 852
56	40	345	222
257	275	346	206
260	370	347	207
261	40	350	237
262	362	351	202
263	210	352	251
264	357	353	211
265-267	40	354	330
270	367	355	241
271	245	356	214
272	255	357	324
273	257	360	320
274	225	361	344
275	361	362	345
276	226	363	242
277	276	364	223
300	350	365	213
301	265	366	224
302	266	367	366
303	306	370	375
304	216	371	205
305	221	372	243
306	217	374	201
307	200	375	354
310	254	376	356

MS 1250 to Mazovia

For the conversion of MS 1250 to Mazovia, all characters not in the following table are mapped unchanged.

MS 1250	Mazovia	MS 1250	Mazovia
200-213	40	310-311	40
214	230	312	220
215-216	40	313-320	40
217	240	321	245
220-233	40	322	40
234	236	323	243
235-236	40	324-325	40
237	246	326	231
240	377	327-333	40
241-242	40	334	232
243	234	335-336	40
244	40	337	341
245	217	340-341	40
246-252	40	342	203
253	256	343	40
254	252	344	204
255-256	40	345	40
257	241	346	215
260	370	347	207
261	361	350	40
262	40	351	202
263	222	352	221
264	40	353	211
265	346	354-355	40
266	40	356	214
267	372	357-360	40
270	40	361	244
271	206	362	40

Conversions Performed					
MS 1250	MS 1250 Mazovia MS 1250				
272	40	363	242		
273	257	364	223		
274-276	40	365	40		
277	247	366	224		
300-303	40	367	366		
304	216	370-373	40		
305	40	374	201		
306	225	375-376	40		
307	200				

MS 1250 to DHN For the conversion of MS 1250 to DHN, all characters not in the following table are mapped unchanged.

Conversions Performed					
MS 1250	DHN	MS 1250	DHN		
200-213	40	306	201		
214	206	307-311	40		
215-216	40	312	202		
217	207	313-320	40		
220-233	40	321	204		
234	217	322	40		
235-236	40	323	205		
237	220	324-325	40		
240	377	326	231		
241-242	40	327-333	40		
243	203	334	232		
244	40	335-336	40		
245	200	337	341		
246-252	40	340	40		

		Conversion	ns Performed	
	MS 1250	DHN	MS 1250	DHN
	253	256	341	240
	254	252	342-345	40
	255-256	40	346	212
	257	210	347-351	40
	260	370	352	213
	261	361	353-354	40
	262	40	355	241
	263	214	356-360	40
	264	40	361	215
	265	346	362	40
	266	40	363	216
	267	372	364	223
	270	40	365	40
	271	211	366	224
	272	40	367	366
	273	257	370-371	40
	274-276	40	372	243
	277	221	373-376	40
	300-305	40		
FILES	/usr/lib/iconv/*		conversion modules	
	/usr/lib/iconv/*		conversion tables	
	/usr/lib/iconv/i		list of conversions sup tables	ported by conversion
SEE ALSO	iconv(1), iconv(3C)	,iconv(5)		

NAME | iconv_1251 – code set conversion tables for MS 1251 (Windows Cyrillic)

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported					
Code	Symbol	Target Code	Symbol	Target Output	
MS 1251	win5	ISO 8859-5	iso5	ISO 8859-5 Cyrillic	
MS 1251	win5	KOI8-R	koi8	KOI8-R	
MS 1251	win5	PC Cyrillic	alt	Alternative PC Cyrillic	
MS 1251	win5	Mac Cyrillic	mac	Macintosh Cyrillic	

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

MS 1251 to ISO 8859-5 For the conversion of MS 1251 to ISO 8859-5, all characters not in the following table are mapped unchanged.

Conversions Performed				
MS 1251	ISO 8859-5	MS 1251	ISO 8859-5	
24	4	310	270	
200	242	311	271	
201	243	312	272	
202	40	313	273	
203	363	314	274	
204-207	40	315	275	
210	255	316	276	
211	40	317	277	
212	251	320	300	
213	40	321	301	
214	252	322	302	
215	254	323	303	
216	253	324	304	
217	257	325	305	

MS 1251 ISO 8859-5 MS 1251 ISO 88					
20	362	326	306		
221-227	40	327	307		
230	255	330	310		
231	40	331	311		
232	371	332	312		
233	40	333	313		
234	372	334	314		
235	374	335	315		
236	373	336	316		
237	377	337	317		
241	256	340	320		
242	376	341	321		
243	250	342	322		
244-247	40	343	323		
250	241	344	324		
251	40	345	325		
252	244	346	326		
253-254	40	347	327		
255	55	350	330		
256	40	351	331		
257	247	352	332		
260-261	40	353	333		
262	246	354	334		
263	366	355	335		
264-267	40	356	336		
270	361	357	337		
271	360	360	340		
272	364	361	341		

Conversions Performed					
MS 1251	ISO 8859-5	MS 1251	ISO 8859-5		
273	40	362	342		
274	370	363	343		
275	245	364	344		
276	365	365	345		
277	367	366	346		
300	260	367	347		
301	261	370	350		
302	262	371	351		
303	263	372	352		
304	264	373	353		
305	265	374	354		
306	266	375	355		
307	267	376	356		

MS 1251 to KOI8-R

For the conversion of MS 1251 to KOI8-R , all characters not in the following table are mapped unchanged.

Conversions Performed					
MS 1251	KOI8-R	MS 1251	KOI8-R		
24	4	310	351		
200	261	311	352		
201	262	312	353		
202	40	313	354		
203	242	314	355		
204-207	40	315	356		
210	255	316	357		
211	40	317	360		
212	271	320	362		
213	40	321	363		

MS 1251 KOI8-R MS 1251 KOI8-R					
214	272	322	364		
215	274	323	365		
216	273	324	346		
217	277	325	350		
220	241	326	343		
221-227	40	327	376		
230	255	330	373		
231	40	331	375		
232	251	332	377		
233	40	333	371		
234	252	334	370		
235	254	335	374		
236	253	336	340		
237	257	337	361		
241	276	340	301		
242	256	341	302		
243	270	342	327		
244-247	40	343	307		
250	263	344	304		
251	40	345	305		
252	264	346	326		
253-254	40	347	332		
255	55	350	311		
256	40	351	312		
257	267	352	313		
260-261	40	353	314		
262	266	354	315		
263	246	355	316		

Conversions Performed				
MS 1251	KOI8-R	MS 1251	KOI8-R	
264-267	40	356	317	
270	243	357	320	
271	260	360	322	
272	244	361	323	
273	40	362	324	
274	250	363	325	
275	265	364	306	
276	245	365	310	
277	247	366	303	
300	341	367	336	
301	342	370	333	
302	367	371	335	
303	347	372	337	
304	344	373	331	
305	345	374	330	
306	366	375	334	
307	372	376	300	

MS 1251 to PC Cyrillic

For the conversion of MS 1251 to PC Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed				
MS 1251 PC Cyrillic MS 1251 PC Cy				
24	4	332	232	
200-207	40	333	233	
210	260	334	234	
211-227	40	335	235	
230	260	336	236	
231-247	40	337	237	

MS 1251	PC Cyrillic	MS 1251	PC Cyrillic
50	360	340	240
51-254	40	341	241
255	55	342	242
256-267	40	343	243
270	361	344	244
271-277	40	345	245
300	200	346	246
301	201	347	247
302	202	350	250
303	203	351	251
304	204	352	252
305	205	353	253
306	206	354	254
307	207	355	255
310	210	356	256
311	211	357	257
312	212	360	340
313	213	361	341
314	214	362	342
315	215	363	343
316	216	364	344
317	217	365	345
320	220	366	346
321	221	367	347
322	222	370	350
323	223	371	351
324	224	372	352
325	225	373	353

Conversions Performed			
MS 1251 PC Cyrillic MS 1251 PC Cyrillic			
326	226	374	354
327	227	375	355
330	230	376	356
331	231		

MS 1251 to Mac Cyrillic

For the conversion of MS 1251 to Mac Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed			
MS 1251	Mac Cyrillic	MS 1251	Mac Cyrillic
24	4	260	241
200	253	262	247
201	256	263	264
202	40	264	266
203	257	266	246
204	327	267	245
205	311	270	336
206	240	271	334
207-211	40	272	271
212	274	273	310
213	40	274	300
214	276	275	301
215	315	276	317
216	40	277	273
217	332	300	200
220	254	301	201
221	324	302	202
222	325	303	203
223	322	304	204

MS 1251 Mac Cyrillic MS 1251 Mac Cy			
224	323	305	205
225	40	306	206
226	320	307	207
227	321	310	210
230	40	311	211
231	252	312	212
232	275	313	213
233	40	314	214
234	277	315	215
235	316	316	216
236	40	317	217
237	333	320	220
240	312	321	221
241	330	322	222
242	331	323	223
243	267	324	224
244	377	325	225
245	242	326	226
246	40	327	227
247	244	330	230
250	335	331	231
252	270	332	232
253	307	333	233
254	302	334	234
255	55	335	235
256	250	336	236
257	272	337	237
355	316		

FILES	/usr/lib/iconv/*.so	conversion modules
	/usr/lib/iconv/*.t	conversion tables
	/usr/lib/iconv/iconv_data	list of conversions supported by conversion tables
SEE ALSO	iconv(1), iconv(3C), iconv(5)	

NAME | iconv – code set conversion tables

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported

Code	Symbol	Target Code	Symbol	Target Output
ISO 646	646	ISO 8859-1	8859	US ASCII
ISO 646de	646de	ISO 8859-1	8859	German
ISO 646da	646da	ISO 8859-1	8859	Danish
ISO 646en	646en	ISO 8859-1	8859	English ASCII
ISO 646es	646es	ISO 8859-1	8859	Spanish
ISO 646fr	646fr	ISO 8859-1	8859	French
ISO 646it	646it	ISO 8859-1	8859	Italian
ISO 646sv	646sv	ISO 8859-1	8859	Swedish
ISO 8859-1	8859	ISO 646	646	7 bit ASCII
ISO 8859-1	8859	ISO 646de	646de	German
ISO 8859-1	8859	ISO 646da	646da	Danish
ISO 8859-1	8859	ISO 646en	646en	English ASCII
ISO 8859-1	8859	ISO 646es	646es	Spanish
ISO 8859-1	8859	ISO 646fr	646fr	French
ISO 8859-1	8859	ISO 646it	646it	Italian
ISO 8859-1	8859	ISO 646sv	646sv	Swedish
ISO 8859-16	iso16	ISO 8859-2	iso2	ISO Latin 2
ISO 8859-2	iso2	ISO 8859-16	iso16	ISO Latin 10
ISO 8859-16	iso16	IBM 850	ibm850	IBM 850 code page
ISO 8859-16	iso16	IBM 870	ibm870	IBM 870 code page
ISO 8859-2	iso2	MS 1250	win2	Windows Latin 2
ISO 8859-2	iso2	MS 852	dos2	MS-DOS Latin 2
ISO 8859-2	iso2	Mazovia	maz	Mazovia
IBM 850	ibm850	ISO 8859-16	iso16	ISO Latin 10
IBM 870	ibm870	ISO 8859-16	iso16	ISO Latin 10
MS 1250	win2	DHN	dhn	Dom Handlowy Nauki
MS 852	dos2	ISO 8859-2	iso2	ISO Latin 2
MS 852	dos2	MS 1250	win2	Windows Latin 2
MS 852	dos2	Mazovia	maz	Mazovia
MS 852	dos2	DHN	dhn	Dom Handlowy Nauki
Mazovia	maz	ISO 8859-2	iso2	ISO Latin 2
Mazovia	maz	MS 1250	win2	Windows Latin 2
Mazovia	maz	MS 852	dos2	MS-DOS Latin 2
Mazovia	maz	DHN	dhn	Dom Handlowy Nauki
DHN	dhn	ISO 8859-2	iso2	ISO Latin 2
DHN	dhn	MS 1250	win2	Windows Latin 2
DHN	dhn	MS 852	dos2	MS-DOS Latin 2
DHN	dhn	Mazovia	maz	Mazovia
ISO 8859-5	iso5	KOI8-R	koi8	KOI8-R
ISO 8859-5	iso5	PC Cyrillic	alt	Alternative PC Cyrillic
ISO 8859-5	iso5	MS 1251	win5	Windows Cyrillic
ISO 8859-5	iso5	Mac Cyrillic	mac	Macintosh Cyrillic
KOI8-R	koi8	ISO 8859-5	iso5	ISO 8859-5 Cyrillic
KOI8-R	koi8 koi8	PC Cyrillic	alt	Alternative PC Cyrillic
KOI8-R	koi8 koi8	MS 1251	win5	Windows Cyrillic
KOI8-R	koi8	Mac Cyrillic	mac	Macintosh Cyrillic
PC Cyrillic	alt alt	ISO 8859-5	iso5 koi8	ISO 8859-5 Cyrillic KOI8-R
PC Cyrillic		KOI8-R MS 1251		
PC Cyrillic	alt	PIC 1201	win5	Windows Cyrillic

PC Cyrillic MS 1251 MS 1251 MS 1251 MS 1251 Mac Cyrillic Mac Cyrillic Mac Cyrillic	alt win5 win5 win5 mac mac mac	Mac Cyrillic ISO 8859-5 KOI8-R PC Cyrillic Mac Cyrillic ISO 8859-5 KOI8-R PC Cyrillic	mac iso5 koi8 alt mac iso5 koi8 alt	Macintosh Cyrillic ISO 8859-5 Cyrillic KOI8-R Alternative PC Cyrillic Macintosh Cyrillic ISO 8859-5 Cyrillic KOI8-R Alternative PC Cyrillic
-	mac	PC Cyrillic	alt	Alternative PC Cyrillic
Mac Cyrillic	mac	MS 1251	win5	Windows Cyrillic

CONVERSIONS

The conversions are performed according to the tables contained in the manual pages cross-referenced in the Index of Conversion Code Tables below.

Index of Conversion (Index of Conversion Code Tables			
Code	Target Code	See Manual Page		
ISO 646	ISO 8859-1	iconv_646 (5)		
ISO 646de	ISO 8859-1			
ISO 646da	ISO 8859-1			
ISO 646en	ISO 8859-1			
ISO 646es	ISO 8859-1			
ISO 646fr	ISO 8859-1			
ISO 646it	ISO 8859-1			
ISO 646sv	ISO 8859-1			
ISO 8859-1	ISO 646	iconv_8859-1 (5)		
ISO 8859-1	ISO 646de			
ISO 8859-1	ISO 646da			
ISO 8859-1	ISO 646en			
ISO 8859-1	ISO 646es			
ISO 8859-1	ISO 646fr			
ISO 8859-1	ISO 646it			
ISO 8859-1	ISO 646sv			
ISO 8859-2	MS 1250	iconv_8859-2 (5)		
ISO 8859-2	MS 852			
ISO 8859-2	Mazovia			
ISO 8859-2	DHN			

Index of Conversion Code Tables		
MS 1250	ISO 8859-2	iconv_1250 (5)
MS 1250	MS 852	
MS 1250	Mazovia	
MS 1250	DHN	
MS 852	ISO 8859-2	iconv_852 (5)
MS 852	MS 1250	
MS 852	Mazovia	
MS 852	DHN	
Mazovia	ISO 8859-2	iconv_maz (5)
Mazovia	MS 1250	
Mazovia	MS 852	
Mazovia	DHN	

Index of Conversion Code Tables		
Code	Target Code	See Manual Page
DHN	ISO 8859-2	iconv_dhn (5)
DHN	MS 1250	
DHN	MS 852	
DHN	Mazovia	
ISO 8859-5	KOI8-R	iconv_8859-5 (5)
ISO 8859-5	PC Cyrillic	
ISO 8859-5	MS 1251	
ISO 8859-5	Mac Cyrillic	
KOI8-R	ISO 8859-5	iconv_koi8-r (5)
KOI8-R	PC Cyrillic	
KOI8-R	MS 1251	
KOI8-R	Mac Cyrillic	
PC Cyrillic	ISO 8859-5	iconv_pc_cyr (5)
PC Cyrillic	KOI8-R	

¹²⁴ man pages section 5: Standards, Environments, and Macros • Last Revised 5 Dec 2001

Index of Conversion Code Tables		
PC Cyrillic	MS 1251	
PC Cyrillic	Mac Cyrillic	
MS 1251	ISO 8859-5	iconv_1251 (5)
MS 1251	KOI8-R	
MS 1251	PC Cyrillic	
MS 1251	Mac Cyrillic	
Mac Cyrillic	ISO 8859-5	iconv_mac_cyr (5)
Mac Cyrillic	KOI8-R	
Mac Cyrillic	PC Cyrillic	
Mac Cyrillic	MS 1251	

FILES /usr/lib/iconv/*.so conversion modules

/usr/lib/iconv/*.t Conversion tables.

/usr/lib/iconv/geniconvtbl/binarytables/*.bt Conversion binary tables.

/usr/lib/iconv/iconv_data
List of conversions supported by conversion tables.

SEE ALSO iconv(1), iconv(3C), iconv_1250(5), iconv_1251(5), iconv_646(5), iconv_852(5), iconv_8859-1(5), iconv_8859-2(5), iconv_8859-5(5), iconv_dhn(5), iconv_koi8-r(5), iconv_mac_cyr(5), iconv_maz(5), iconv_pc_cyr(5), iconv_unicode(5)

iconv_646(5)

NAME | iconv_646 - code set conversion tables for ISO 646

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported				
Code	Symbol	Target Code	Symbol	Target Output
ISO 646	646	ISO 8859-1	8859	US ASCII
ISO 646de	646de	ISO 8859-1	8859	German
ISO 646da	646da	ISO 8859-1	8859	Danish
ISO 646en	646en	ISO 8859-1	8859	English ASCII
SO 646es	646es	ISO 8859-1	8859	Spanish
ISO 646fr	646fr	ISO 8859-1	8859	French
ISO 646it	646it	ISO 8859-1	8859	Italian
SO 546sv	646sv	ISO 8859-1	8859	Swedish

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

ISO 646 (US ASCII) to ISO 8859-1 ISO 646de (GERMAN) to ISO 8859-1 For the conversion of ISO 646 to ISO 8859-1, all characters in ISO 646 can be mapped unchanged to ISO 8859-1

For the conversion of ISO 646de to ISO 8859-1, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 646de	ISO 8859-1	ISO 646de	ISO 8859-1	
100	247	173	344	
133	304	174	366	
134	326	175	374	
135	334	176	337	

ISO 646da (DANISH) to ISO 8859-1 For the conversion of ISO 646da to ISO 8859-1, all characters not in the following table are mapped unchanged.

iconv_646(5)

Conversions Performed					
ISO 646da	ISO 8859-1	ISO 646da	ISO 8859-1		
133	306	173	346		
134	330	174	370		
135	305	175	345		

ISO 646en (ENGLISH ASCII) to ISO 8859-1 For the conversion of ISO 646en to ISO 8859-1, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 646en ISO 8859-1				
043	243			

ISO 646es (SPANISH) to ISO 8859-1

For the conversion of ISO 646es to ISO 8859-1, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 646es	ISO 8859-1	ISO 646es	ISO 8859-1	
100	247	173	260	
133	241	174	361	
134	321	175	347	
135	277			

ISO 646fr (FRENCH) to ISO 8859-1

For the conversion of ISO 646fr to ISO 8859-1, all characters not in the following table are mapped unchanged.

Conversions Performed					
ISO 646fr	ISO 8859-1	ISO 646fr	ISO 8859-1		
043	243	173	351		
100	340	174	371		
133	260	175	350		
134	347	176	250		

iconv_646(5)

Conversions Performed					
ISO 646fr ISO 8859-1 ISO 646fr ISO 8859-1					
135	247				

ISO 646it (ITALIAN) to ISO 8859-1

For the conversion of ISO 646it to ISO 8859-1, all characters not in the following table are mapped unchanged.

Conversions Performed					
ISO 646it	ISO 8859-1	ISO 646it	ISO 8859-1		
043	243	140	371		
100	247	173	340		
133	260	174	362		
134	347	175	350		
135	351	176	354		

ISO 646sv (SWEDISH) to ISO 8859-1 For the conversion of ISO 646sv to ISO 8859-1, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 646sv	ISO 8859-1	ISO 646sv	ISO 8859-1	
100	311	140	351	
133	304	173	344	
134	326	174	366	
135	305	175	345	
136	334	176	374	

FILES	/usr/lib/iconv/*.so	conversion modules
	/usr/lib/iconv/*.t	conversion tables
	/usr/lib/iconv/iconv_data	list of conversions supported by conversion tables
SEE ALSO	<pre>iconv(1), iconv(3C), iconv(5)</pre>	

NAME | iconv_852 – code set conversion tables for MS 852 (MS-DOS Latin 2)

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported					
Code	Symbol	Target Code	Symbol	Target Output	
MS 852	dos2	ISO 8859-2	iso2	ISO Latin 2	
MS 852	dos2	MS 1250	win2	Windows Latin 2	
MS 852	dos2	Mazovia	maz	Mazovia	
MS 852	dos2	DHN	dhn	Dom Handlowy Nauki	

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

MS 852 to ISO 8859-2 For the conversion of MS 852 to ISO 8859-2, all characters not in the following table are mapped unchanged.

Conversions Performed				
MS 852	ISO 8859-2	MS 852	ISO 8859-2	
24-177	40	271-274	40	
200	307	275	257	
201	374	276	277	
202	351	277-305	40	
203	342	306	303	
204	344	307	343	
205	371	310-316	40	
206	346	317	244	
207	347	320	360	
210	263	321	320	
211	353	322	317	
212	325	323	313	
213	365	324	357	
214	356	325	322	

MS 852	ISO 8859-2	MS 852	ISO 8859-2
215	254	326	315
216	304	327	316
217	306	330	354
220	311	331-334	40
221	305	335	336
222	345	336	331
223	364	337	40
224	366	340	323
225	245	341	337
226	265	342	324
227	246	343	321
230	266	344	361
231	326	345	362
232	334	346	251
233	253	347	271
234	273	350	300
235	243	351	332
236	327	352	340
237	350	353	333
240	341	354	375
241	355	355	335
242	363	356	376
243	372	357	264
244	241	360	255
245	261	361	275
246	256	362	262
247	276	363	267
250	312	364	242

Conversions Performed					
MS 852	MS 852 ISO 8859-2		ISO 8859-2		
251	352	365	247		
252	40	366	367		
253	274	367	270		
254	310	370	260		
255	272	371	250		
256-264	40	372	377		
265	301	374	330		
266	302	375	370		
267	314	376	40		
270	252				

MS 852 to MS 1250 For the conversion of MS 852 to MS 1250, all characters not in the following table are mapped unchanged.

	Convers	sions Performed	
MS 852	MS 1250	MS 852	MS 1250
200	307	270	252
201	374	271-274	40
202	351	275	257
203	342	276	277
204	344	277-305	40
205	371	306	303
206	346	307	343
207	347	310-316	40
210	263	317	244
211	353	320	360
212	325	321	320
213	365	322	317
214	356	323	313

MS 852	MS 1250	MS 852	MS 1250
:15	217	324	357
16	304	325	322
217	306	326	315
220	311	327	316
221	305	330	354
222	345	331-334	40
223	364	335	336
224	366	336	331
225	274	337	40
226	276	340	323
227	214	341	337
230	234	342	324
231	326	343	321
232	334	344	361
233	215	345	362
234	235	346	212
235	243	347	232
236	327	350	300
237	350	351	332
240	341	352	340
241	355	353	333
242	363	354	375
243	372	355	335
244	245	356	376
245	271	357	264
246	216	360	255
247	236	361	275
250	312	362	262

	Convers	sions Performed	
MS 852	MS 1250	MS 852	MS 1250
251	352	363	241
252	254	364	242
253	237	365	247
254	310	366	367
255	272	367	270
256	253	370	260
257	273	371	250
260-264	40	372	377
265	301	374	330
266	302	375	370
267	314	376	40

MS 852 to Mazovia For the conversion of MS 852 to Mazovia, all characters not in the following table are mapped unchanged.

	Convers	sions Performed	
MS 852	Mazovia	MS 852	Mazovia
205	40	246-247	40
206	215	250	220
210	222	251	221
212-213	40	253	246
215	240	254-270	40
217	225	275	241
220-226	40	276	247
227	230	306-336	40
230	236	340	243
233-234	40	342	40
235	234	343	245
236-243	40	344	244

Conversions Performed						
MS 852	Mazovia	MS 852	Mazovia			
244	217	345-375	40			
245	206					

MS 852 to DHN For the conversion of MS 852 to DHN, all characters not in the following table are mapped unchanged.

	Conversions Performed					
MS 852	DHN	MS 852	DHN			
200-205	40	244	200			
206	212	245	211			
207	40	246-247	40			
210	214	250	202			
211-214	40	251	213			
215	207	253	220			
216	40	254-270	40			
217	201	275	210			
220-226	40	276	221			
227	206	306-336	40			
230	217	340	205			
233-234	40	342	40			
235	203	343	204			
236-237	40	344	215			
242	216	345-375	40			
252	254					
/usr/lib/ic	0777/* 50	conversion module				
/usr/lib/ic			.0			
	onv/iconv_data	conversion tables list of conversions supported by conve tables				

SEE ALSO | iconv(1), iconv(3C), iconv(5)

NAME | iconv_8859-1 – code set conversion tables for ISO 8859-1 (Latin 1)

DESCRIPTION

The following code set conversions are supported:

		Code Set Conversion	s Supported	
Code	Symbol	Target Code	Symbol	Target Output
ISO 8859-1	8859	ISO 646	646	7 bit ASCII
ISO 8859-1	8859	ISO 646de	646de	German
ISO 8859-1	8859	ISO 646da	646da	Danish
ISO 8859-1	8859	ISO 646en	646en	English ASCII
ISO 8859-1	8859	ISO 646es	646es	Spanish
ISO 8859-1	8859	ISO 646fr	646fr	French
ISO 8859-1	8859	ISO 646it	646it	Italian
ISO 8859-1	8859	ISO 646sv	646sv	Swedish

CONVERSIONS The conversions are performed according to the following tables. All values in the tables are given in octal.

ISO 8859-1 to ISO 646 (7-bit ASCII) For the conversion of ISO 8859-1 to ISO 646, all characters not in the following table are mapped unchanged.

	Conv	verte	ed to	o Uno	derso	core	'_'	(137)	
	200	201	202	203	204	205	206	207	
	210	211	212	213	214	215	216	217	
	220	221	222	223	224	225	226	227	
	230	231	232	233	234	235	236	237	
	240	241	242	243	244	245	246	247	
	250	251	252	253	254	255	256	257	
	260	261	262	263	264	265	266	267	
	270	271	272	273	274	275	276	277	
	300	301	302	303	304	305	306	307	
	310	311	312	313	314	315	316	317	
	320	321	322	323	324	325	326	327	
	330	331	332	333	334	335	336	337	
	340	341	342	343	344	345	346	347	
	350	351	352	353	354	355	356	357	
	360	361	362	363	364	365	366	367	
	370	371	372	373	374	375	376	377	
9-1 to ISO	For	the	conv	versi	on o	f ISC) 885	9-1 to I	S
					-	-		-	

ISO 8859-1 to ISO
646de (GERMAN)For the conversion of ISO 8859-1 to ISO 646de, all characters not in the following
tables are mapped unchanged.

iconv_8859-1(5)

Conversions Performed						
ISO 8859-1	ISO 8859-1 ISO 646de ISO 8859-1 ISO 646					
247	100	337	176			
304	133	344	173			
326	134	366	174			
334	135	374	175			
133 134 135 1 200 201 202 203 2 210 211 212 213 2 220 221 222 223 2 230 231 232 233 2 240 241 242 243 2 250 251 252 253 2 260 261 262 263 2 270 271 272 273 2	204 205 206 207 214 215 216 217 224 225 226 227 234 235 236 237 244 245 246 254 255 256 257 264 265 266 267					
300 301 302 303 310 311 312 313 3 320 321 322 323 3 340 311 332 333 3 340 341 342 343 350 351 352 353 3	324 325 327 335 336 337 345 346 347					
60 361 362 363 3 70 371 372 373						

ISO 8859-1 to ISO 646da (DANISH)

For the conversion of ISO 8859-1 to ISO 646da, all characters not in the following tables are mapped unchanged.

Conversions Performed							
ISO 8859-1	ISO 646da	ISO 8859-1	ISO 646da				
305	135	345	175				
306	133	346	173				
330	134	370	174				
Converted to Underscore '_' (137)							
133 134 135 173 174	 175						
200 201 202 203 204 210 211 212 213 214							

200	201	202	203	204	205	206	207	
210	211	212	213	214	215	216	217	
220	221	222	223	224	225	226	227	
230	231	232	233	234	235	236	237	
240	241	242	243	244	245	246	247	
250	251	252	253	254	255	256	257	
260	261	262	263	264	265	266	267	
270	271	272	273	274	275	276	277	
300	301	302	303	304			307	

 310
 311
 312
 313
 314
 315
 316
 317

 320
 321
 322
 323
 324
 325
 326
 327

 331
 332
 333
 334
 335
 336
 337

 340
 341
 342
 343
 344

 347

 350
 351
 352
 353
 354
 355
 356

 360
 361
 362
 363
 364
 365
 366

 371
 372
 373
 374
 374
 376
 377

ISO 8859-1 to ISO 646en (ENGLISH ASCII) For the conversion of ISO 8859-1 to ISO 646en, all characters not in the following tables are mapped unchanged.

							Conversio	ons Performed	l	
				ISO	8859-	1			ISO 646en	
243								043		
Conv	verte	ed to	o Uno	ders	core	'_'	(137)			
043	201	202	202	204	205	200	207			
					205					
					215					
					235					
	241		200		245					
			253		255					
260	261	262	263	264	265	266	267			
270	271	272	273	274	275	276	277			
300	301	302	303	304	305	306	307			
310	311	312	313	314	315	316	317			
320	321	322	323	324	325	326	327			
330	331	332	333	334	335	336	337			
340	341	342	343	344	345	346	347			
350	351	352	353	354	355	356	357			
360	361	362	363	364	365	366	367			
370	371	372	373	374	375	376	377			

ISO 8859-1 to ISO 646fr (FRENCH) For the conversion of ISO 8859-1 to ISO 646fr, all characters not in the following tables are mapped unchanged.

	Conversion	s Performed	
ISO 8859-1	ISO 646fr	ISO 8859-1	ISO 646fr
243	043	347	134
247	135	350	175
250	176	351	173
260	133	371	174
340	100		

iconv_8859-1(5)

Conv	verte	ed to	o Uno	lers	core	'_'	(137)
043							
100	133	134	135	173	174	175	176
200	201	202	203	204	205	206	207
210	211	212	213	214	215	216	217
220	221	222	223	224	225	226	227
230	231	232	233	234	235	236	237
240	241	242		244	245	246	
	251	252	253	254	255	256	257
	261	262	263	264	265	266	267
270	271	272	273	274	275	276	277
300	301	302	303	304	305	306	307
310	311	312	313	314	315	316	317
320	321	322	323	324	325	326	327
330	331	332	333	334	335	336	337
	341	342	343	344	345	346	
		352	353	354	355	356	357
360	361	362	363	364	365	366	367
370		372	373	374	375	376	377

ISO 8859-1 to ISO 646it (ITALIAN)

For the conversion of ISO 8859-1 to ISO 646it, all characters not in the following tables are mapped unchanged.

	Conversion	s Performed	
ISO 8859-1	ISO 646it	ISO 8859-1	ISO 646it
243	043	350	175
247	100	351	135
260	133	354	176
340	173	362	174
347	134	371	140
Converted to Undersco	ore '_' (137)		
043			
100 133 134 135 173	174 175 176		
200 201 202 203 204	205 206 207		
210 211 212 213 214	215 216 217		
220 221 222 223 224	225 226 227		
230 231 232 233 234	235 236 237		
240 241 242 244	245 246		
250 251 252 253 254	255 256 257		

 210
 211
 212
 213
 214
 215
 216
 217

 220
 221
 222
 223
 224
 225
 226
 227

 230
 231
 232
 233
 234
 235
 236
 237

 240
 241
 242
 244
 245
 246

 250
 251
 252
 253
 254
 255
 256
 257

 261
 262
 263
 264
 265
 266
 267

 270
 271
 272
 273
 374
 275
 276
 277

 300
 301
 302
 303
 304
 305
 306
 307

 310
 311
 312
 313
 314
 315
 316
 317

 320
 321
 322
 323
 324
 325
 326
 327

 330
 331
 332
 333
 334
 335
 336
 337

</tabu/>

341 342 343 344 345 346 352 353 354 355 356 357

ISO 8859-1 to ISO 646es (SPANISH) For the conversion of ISO 8859-1 to ISO 646es, all characters not in the following tables are mapped unchanged.

Conversions Performed				
ISO 8859-1	ISO 646es	ISO 8859-1	ISO 646es	
241	133	321	134	
247	100	347	175	
260	173	361	174	
277	135			

Conv	verte	ed to	o Uno	derso	core	'_'	(137)
100	133	134	135	173	174	175	
200	201	202	203	204	205	206	207
210	211	212	213	214	215	216	217
220	221	222	223	224	225	226	227
230	231	232	233	234	235	236	237
240		242	243	244	245	246	
250	251	252	253	254	255	256	257
	261	262	263	264	265	266	267
270	271	272	273	274	275	276	
300	301	302	303	304	305	306	307
310	311	312	313	314	315	316	317
320		322	323	324	325	326	327
330	331	332	333	334	335	336	337
340	341	342	343	344	345	346	
350	351	352	353	354	355	356	357
360		362	363	364	365	366	367
370	371	372	373	374	375	376	377

ISO 8859-1 to ISO 646sv (SWEDISH) For the conversion of ISO 8859-1 to ISO 646sv, all characters not in the following tables are mapped unchanged.

	Convers	ions Performed	
ISO 8859-1	ISO 646sv	ISO 8859-1	ISO 646sv
304	133	344	173
305	135	345	175
311	100	351	140
326	134	366	174
334	136	374	176

iconv_8859-1(5)

_8859-1(5)		
	Converted to Underscore '_' (137)	
	100 133 134 135 136 140 173 174 175 176	
	200 201 202 203 204 205 206 207	
	210 211 212 213 214 215 216 217	
	220 221 222 223 224 225 226 227	
	230 231 232 233 234 235 236 237	
	240 241 242 243 244 245 246 247	
	250 251 252 253 254 255 256 257	
	260 261 262 263 264 265 266 267	
	270 271 272 273 274 275 276 277	
	300 301 302 303 306 307	
	310 312 313 314 315 316 317	
	320 321 322 323 324 325 327	
	330 331 332 333 335 336 337	
	340 341 342 343 346 347 250 250 251 255 255 255	
	350 352 353 354 355 356 357 360 361 362 363 364 365 367	
	360 361 362 363 364 365 367 370 371 372 373 375 376 377	
FILES	/usr/lib/iconv/*.so	conversion modules
	/usr/lib/iconv/*.t	conversion tables
	/usr/lib/iconv/iconv_data	list of conversions supported by conversion tables
SEE ALSO	<pre>iconv(1), iconv(3C), iconv(5)</pre>	
SEE ALSO	100110(1), 100110(3C), 100110(3)	

NAME | iconv_8859-2 - code set conversion tables for ISO 8859-2 (Latin 2)

DESCRIPTION

The following code set conversions are supported:

	-	Code Set Conversions	Supported	
Code	Symbol	Target Code	Symbol	Target Output
ISO 8859-2	iso2	MS 1250	win2	Windows Latin 2
ISO 8859-2	iso2	MS 852	dos2	MS-DOS Latin 2
ISO 8859-2	iso2	Mazovia	maz	Mazovia
ISO 8859-2	iso2	DHN	dhn	Dom Handlowy Nauki

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

ISO 8859-2 to MS 1250

For the conversion of ISO 8859-2 to MS 1250, all characters not in the following table are mapped unchanged.

	Conversion	s Performed	
ISO 8859-2	MS 1250	ISO 8859-2	MS 1250
24	4	261	271
177-237	40	265	276
241	245	266	234
245	274	267	241
246	214	271	232
251	212	273	235
253	215	274	237
254	217	276	236
256	216	266	236

ISO 8859-2 to MS 852

For the conversion of ISO 8859-2 to MS 852, all characters not in the following table are mapped unchanged.

iconv_8859-2(5)

ISO 8859-2	MS 852	ISO 8859-2	MS 852
24	4	316	327
177-237	40	317	322
240	377	320	321
241	244	321	343
242	364	322	325
243	235	323	340
244	317	324	342
245	225	325	212
246	227	326	231
247	365	327	236
250	371	330	374
251	346	331	336
252	270	332	351
253	233	333	353
254	215	334	232
255	360	335	355
256	246	336	335
257	275	337	341
260	370	340	352
261	245	341	240
262	362	342	203
263	210	343	307
264	357	344	204
265	226	345	222
266	230	346	206
267	363	347	207
270	367	350	237
271	347	351	202

iconv_8859-2(5)

Conversions Performed					
ISO 8859-2	MS 852	ISO 8859-2	MS 852		
272	255	352	251		
273	234	353	211		
274	253	354	330		
275	361	355	241		
276	247	356	214		
277	276	357	324		
300	350	360	320		
301	265	361	344		
302	266	362	345		
303	306	363	242		
304	216	364	223		
305	221	365	213		
306	217	366	224		
307	200	367	366		
310	254	370	375		
311	220	371	205		
312	250	372	243		
313	323	374	201		
314	267	375	354		
315	326	376	356		
366	367				

ISO 8859-2 to Mazovia

For the conversion of ISO 8859-2 to Mazovia, all characters not in the following table are mapped unchanged.

	Conversion	s Performed	
ISO 8859-2	Mazovia	ISO 8859-2	Mazovia
24	4	323	243
177-237	40	324-325	40

iconv_8859-2(5)

ISO 8859-2	Mazovia	ISO 8859-2	Mazovia
240	377	326	231
241	217	327-333	40
242	40	334	232
243	234	335-336	40
244-245	40	337	341
246	230	340-341	40
247-253	40	342	203
254	240	343	40
255-256	40	344	204
257	241	345	40
260	370	346	215
261	206	347	207
262	40	350	40
263	222	351	202
264-265	40	352	221
266	236	353	211
267-273	40	354-355	40
274	246	356	214
275-276	40	357-360	40
277	247	361	244
300-303	40	362	40
304	216	363	242
305	40	364	223
306	225	365	40
307	200	366	224
310-311	40	367	366
312	220	370-373	40
313-320	40	374	201

Conversions Performed				
ISO 8859-2 Mazovia ISO 8859-2 Mazovia				
321	245	375-376	40	
322	40			

ISO 8859-2 to DHN

For the conversion of ISO 8859-2 to DHN, all characters not in the following table are mapped unchanged.

	Conve	rsions Performed	
ISO 8859-2	DHN	ISO 8859-2	DHN
24	4	322	40
177-237	40	323	205
240	377	324-325	40
241	200	326	231
242	40	327-333	40
243	203	334	232
244-245	40	335-336	40
246	206	337	341
247-253	40	340	40
254	207	341	240
255-256	40	342-345	40
257	210	346	212
260	370	347-351	40
261	211	352	213
262	40	353-354	40
263	214	355	241
264-265	40	356-360	40
266	217	361	215
267-273	40	362	40
274	220	363	216
275-276	40	364	223

	Conversions Performed				
	ISO 8859-2	DHN	ISO 8859-2	DHN	
	277	221	365	40	
	300-305	40	366	224	
	306	201	367	366	
	307-311	40	370-371	40	
	312	202	372	243	
	313-320	40	373-376	40	
	321	204			
			•	·	
FILES	/usr/lib/iconv/*		conversion modules		
	/usr/lib/iconv/*		conversion tables		
	/usr/lib/iconv/i		list of conversions sup tables	ported by conversion	
SEE ALSO	iconv(1), iconv(3C)	,iconv(5)			
		, (-)			

NAME | iconv_8859-5 - code set conversion tables for ISO 8859-5 (Cyrillic)

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported						
Code Symbol Target Code Symbol Target Output						
ISO 8859-5	iso5	KOI8-R	koi8	KOI8-R		
ISO 8859-5	iso5	PC Cyrillic	alt	Alternative PC Cyrillic		
ISO 8859-5	iso5	MS 1251	win5	Windows Cyrillic		
ISO 8859-5	iso5	Mac Cyrillic	mac	Macintosh Cyrillic		

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

ISO 8859-5 to KOI8-R

For the conversion of ISO 8859-5 to KOI8-R, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 8859-5	KOI8-R	ISO 8859-5	KOI8-R	
24	4	320	301	
241	263	321	302	
242	261	322	327	
243	262	323	307	
244	264	324	304	
245	265	325	305	
246	266	327	332	
247	267	330	311	
250	270	331	312	
251	271	332	313	
252	272	333	314	
253	273	334	315	
254	274	335	316	
256	276	336	317	

ISO 8859-5 KOI8-R ISO 8859-5 KOI8-R					
.57	277	337	320		
260	341	340	322		
261	342	341	323		
262	367	342	324		
263	347	343	325		
264	344	344	306		
265	345	345	310		
266	366	346	303		
267	372	347	336		
270	351	350	333		
271	352	351	335		
72	353	352	337		
73	354	353	331		
274	355	354	330		
75	356	355	334		
276	357	356	300		
277	360	357	321		
600	362	360	260		
01	363	361	243		
802	364	362	241		
603	365	363	242		
304	346	364	244		
305	350	365	245		
306	343	366	246		
307	376	367	247		
310	373	370	250		
311	375	371	251		
312	377	372	252		

Conversions Performed				
ISO 8859-5	KOI8-R	ISO 8859-5	KOI8-R	
313	371	373	253	
314	370	374	254	
315	374	375	255	
316	340	376	256	
317	361			

ISO 8859-5 to PC Cyrillic

For the conversion of ISO 8859-5 to PC Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 8859-5	PC Cyrillic	ISO 8859-5	PC Cyrillic	
24	4	307	227	
200-240	40	310	230	
241	360	311	231	
242-254	40	312	232	
255	260	313	233	
256-257	40	314	234	
260	200	315	235	
261	201	316	236	
262	202	317	237	
263	203	320	240	
264	204	321	241	
265	205	322	242	
266	206	323	243	
267	207	324	244	
270	210	325	245	
271	211	326	246	
272	212	327	247	
273	213	330	250	

Conversions Performed					
ISO 8859-5	PC Cyrillic	ISO 8859-5	PC Cyrillic		
274	214	331	251		
275	215	332	252		
276	216	333	253		
277	217	334	254		
300	220	335	255		
301	221	336	256		
302	222	337	257		
303	223	360-374	40		
304	224	375	260		
305	225	376	40		
306	226	365	40		

ISO 8859-5 to MS 1251

For the conversion of ISO 8859-5 to MS 1251, all characters not in the following table are mapped unchanged.

Conversions Performed				
ISO 8859-5	MS 1251	ISO 8859-5	MS 1251	
24	4	317	337	
200-237	40	320	340	
241	250	321	341	
242	200	322	342	
243	201	323	343	
244	252	324	344	
245	275	325	345	
246	262	326	346	
247	257	327	347	
250	243	330	350	
251	212	331	351	
252	214	332	352	

ISO 8859-5 MS 1251 ISO 8859-5 MS 1251					
253	216	333	353		
254	215	334	354		
255	210	335	355		
256	241	336	356		
257	217	337	357		
260	300	340	360		
261	301	341	361		
262	302	342	362		
263	303	343	363		
264	304	344	364		
265	305	345	365		
266	306	346	366		
267	307	347	367		
270	310	350	370		
271	311	351	371		
272	312	352	372		
273	313	353	373		
274	314	354	374		
275	315	355	375		
276	316	356	376		
277	317	357	377		
300	320	360	271		
301	321	361	270		
302	322	362	220		
303	323	363	203		
304	324	364	272		
305	325	365	276		
306	326	366	263		

Conversions Performed				
ISO 8859-5	MS 1251	ISO 8859-5	MS 1251	
307	327	367	277	
310	330	370	274	
311	331	371	232	
312	332	372	234	
313	333	373	236	
314	334	374	235	
315	335	375	210	
316	336	376	242	
376	331			

ISO 8859-5 to Mac Cyrillic

For the conversion of ISO 8859-5 to Mac Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed			
ISO 8859-5	Mac Cyrillic	ISO 8859-5	Mac Cyrillic
24	4	317	237
200-237	40	320	340
240	312	321	341
241	335	322	342
242	253	323	343
243	256	324	344
244	270	325	345
245	301	326	346
246	247	327	347
247	272	330	350
250	267	331	351
251	274	332	352
252	276	333	353
253	40	334	354

ISO 8859-5 Mac Cyrillic ISO 8859-5 Mac Cyrillic					
254	315	335	355		
255	40	336	356		
256	330	337	357		
257	332	340	360		
260	200	341	361		
261	201	342	362		
262	202	343	363		
263	203	344	364		
264	204	345	365		
265	205	346	366		
266	206	347	367		
267	207	350	370		
270	210	351	371		
271	211	352	372		
272	212	353	373		
273	213	354	374		
274	214	355	375		
275	215	356	376		
276	216	357	337		
277	217	360	334		
300	220	361	336		
301	221	362	254		
302	222	363	257		
303	223	364	271		
304	224	365	317		
305	225	366	264		
306	226	367	273		
307	227	370	300		

	Conversions Performed					
	ISO 8859-5	Mac Cyrillic	ISO 8859-5	Mac Cyrillic		
	310	230	371	275		
	311	231	372	277		
	312	232	373	40		
	313	233	374	316		
	314	234	375	40		
	315	235	376	331		
	316	236				
	/ / / / / / / / / / / / /					
FILES	/usr/lib/iconv/* /usr/lib/iconv/*		conversion modules conversion tables			
	/usr/lib/iconv/i		list of conversions sup	ported by conversion		
	/usr/11D/1CONV/1	CONV_data	tables	ported by conversion		
SEE ALSO	iconv(1), iconv(3C)), iconv(5)				

NAME | iconv_dhn - code set conversion tables for DHN (Dom Handlowy Nauki)

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported					
Code Symbol Target Code Symbol Target Output					
DHN	dhn	ISO 8859-2	iso2	ISO Latin 2	
DHN	dhn	MS 1250	win2	Windows Latin 2	
DHN	dhn	MS 852	dos2	MS-DOS Latin 2	
DHN	dhn	Mazovia	maz	Mazovia	

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

DHN to ISO 8859-2

For the conversion of DHN to ISO 8859-2, all characters not in the following table are mapped unchanged.

Conversions Performed				
DHN	ISO 8859-2	DHN	ISO 8859-2	
24-177	40	222	40	
200	241	223	364	
201	306	224	366	
202	312	225-230	40	
203	243	231	326	
204	321	232	334	
205	323	233-237	40	
206	246	240	341	
207	254	241	355	
210	257	242	363	
211	261	243	372	
212	346	244-340	40	
213	352	341	337	
214	263	342-365	40	

iconv_dhn(5)

Conversions Performed					
DHN	DHN ISO 8859-2 DHN ISO 8859-				
215	361	366	367		
216	363	367	40		
217	266	370	260		
220	274	371-376	40		
221	277				

DHN to MS 1250

D For the conversion of DHN to MS 1250, all characters not in the following table are mapped unchanged.

	Conversions Performed				
DHN	MS 1250	DHN	MS 1250		
200	245	233-237	40		
201	306	240	341		
202	312	241	355		
203	243	242	363		
204	321	243	372		
205	323	244-251	40		
206	214	252	254		
207	217	253-255	40		
210	257	256	253		
211	271	257	273		
212	346	260-340	40		
213	352	341	337		
214	263	342-345	40		
215	361	346	265		
216	363	347-360	40		
217	234	361	261		
220	237	362-365	40		
221	277	366	367		

iconv_dhn(5)

Conversions Performed				
DHN MS 1250 DHN MS				
222	40	367	40	
223	364	370	260	
224	366	371	40	
225-230	40	372	267	
231	326	373-376	40	
232	334			

DHN to MS 852 For the conversion of DHN to MS 852, all characters not in the following table are mapped unchanged.

Conversions Performed					
DHN	DHN MS 852 DHN				
200	244	212	206		
201	217	213	251		
202	250	214	210		
203	235	215	344		
204	343	216	242		
205	340	217	230		
206	227	220	253		
207	215	221	276		
210	275	222-375	40		
211	245				

DHN to Mazovia For the conversion of DHN to Mazovia, all characters not in the following table are mapped unchanged.

Conversions Performed				
DHN Mazovia DHN Mazovia				
200	217	212	215	

iconv_dhn(5)

	Conversions Performed				
	DHN	Mazovia	DHN	Mazovia	
	201	225	213	221	
	202	220	214	222	
	203	234	215	244	
	204	245	216	242	
	205	243	217	236	
	206	230	220	246	
	207	240	221	247	
	210	241	222-247	40	
	211	206			
		1		1	
FILES	/usr/lib/iconv/*		conversion modules		
	/usr/lib/iconv/*		conversion tables		
	/usr/lib/iconv/i		ist of conversions sup ables	ported by conversion	
SEE ALSO	iconv(1), iconv(3C)	,1conv(5)			

NAME | iconv_koi8-r - code set conversion tables for KOI8-R

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported					
Code	Symbol	Target Code	Symbol	Target Output	
KOI8-R	koi8	ISO 8859-5	iso5	ISO 8859-5 Cyrillic	
KOI8-R	koi8	PC Cyrillic	alt	Alternative PC Cyrillic	
KOI8-R	koi8	MS 1251	win5	Windows Cyrillic	
KOI8-R	koi8	Mac Cyrillic	mac	Macintosh Cyrillic	

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

KOI8-R to ISO 8859-5 For the conversion of KOI8-R to ISO 8859-5, all characters not in the following table are mapped unchanged.

Conversions Performed				
KOI8-R	ISO 8859-5	KOI8-R	ISO 8859-5	
24	4	320	337	
241	362	321	357	
242	363	322	340	
243	361	323	341	
244	364	324	342	
245	365	325	343	
246	366	327	322	
247	367	330	354	
250	370	331	353	
251	371	332	327	
252	372	333	350	
253	373	334	355	
254	374	335	351	
256	376	336	347	

KOI8-R	ISO 8859-5	KOI8-R	ISO 8859-5
	377	337	352
257			
260	360	340	316
261	242	341	260
262	243	342	261
263	241	343	306
264	244	344	264
265	245	345	265
266	246	346	304
267	247	347	263
270	250	350	305
271	251	351	270
272	252	352	271
273	253	353	272
274	254	354	273
275	255	355	274
276	256	356	275
277	257	357	276
300	356	360	277
301	320	361	317
302	321	362	300
303	346	363	301
304	324	364	302
305	325	365	303
306	344	366	266
307	323	367	262
310	345	370	314
311	330	370	313
312	331	372	267

Conversions Performed				
KOI8-R	ISO 8859-5	KOI8-R	ISO 8859-5	
313	332	373	310	
314	333	374	315	
315	334	375	311	
316	335	376	307	
317	336			

KOI8-R to PC Cyrillic

For the conversion of KOI8-R to PC Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed				
KOI8-R	PC Cyrillic	KOI8-R	PC Cyrillic	
24	4	333	350	
200-242	40	334	355	
243	361	335	351	
244-254	40	336	347	
255	260	337	352	
256-262	40	340	236	
263	360	341	200	
264-274	40	342	201	
275	260	343	226	
276-277	40	344	204	
300	356	345	205	
301	240	346	224	
302	241	347	203	
303	346	350	225	
304	244	351	210	
305	245	352	211	
306	344	353	212	
307	243	354	213	

Conversions Performed				
KOI8-R	PC Cyrillic	KOI8-R	PC Cyrillic	
310	345	355	214	
311	250	356	215	
312	251	357	216	
313	252	360	217	
314	253	361	237	
315	254	362	220	
316	255	363	221	
317	256	364	222	
320	257	365	223	
321	357	366	206	
322	340	367	202	
323	341	370	234	
324	342	371	233	
325	343	372	207	
326	246	373	230	
327	242	374	235	
330	354	375	231	
331	353	376	227	
332	247			

KOI8-R to MS 1251

For the conversion of KOI8-R to MS 1251, all characters not in the following table are mapped unchanged.

Conversions Performed					
KOI8-R	MS 1251	KOI8-R	MS 1251		
24	4	317	356		
200-237	40	320	357		
241	220	321	377		
242	203	322	360		

¹⁶² man pages section 5: Standards, Environments, and Macros • Last Revised 18 Apr 1997

KOI8-R MS 1251 KOI8-R N					
243	270	323	361		
244	272	324	362		
245	276	325	363		
246	263	326	346		
247	277	327	342		
250	274	330	374		
51	232	331	373		
52	234	332	347		
253	236	333	370		
254	235	334	375		
255	210	335	371		
.56	242	336	367		
57	237	337	372		
260	271	340	336		
61	200	341	300		
262	201	342	301		
.63	250	343	326		
264	252	344	304		
65	275	345	305		
266	262	346	324		
67	257	347	303		
270	243	350	325		
71	212	351	310		
.72	214	352	311		
73	216	353	312		
274	215	354	313		
275	210	355	314		
276	241	356	315		

Conversions Performed				
KOI8-R	MS 1251	KOI8-R	KOI8-R MS 1251	
277	217	357	316	
300	376	360	317	
301	340	361	337	
302	341	362	320	
303	366	363	321	
304	344	364	322	
305	345	365	323	
306	364	366	306	
307	343	367	302	
310	365	370	334	
311	350	371	333	
312	351	372	307	
313	352	373	330	
314	353	374	335	
315	354	375	331	
316	355	376	327	
376	227			

KOI8-R to Mac For the conversion of KOI8-R to Mac Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed				
KOI8-R	Mac Cyrillic			
24	4	317	356	
200-237	40	320	357	
240	312	321	337	
241	254	322	360	
242	257	323	361	
243	336	324	362	

man pages section 5: Standards, Environments, and Macros • Last Revised 18 Apr 1997 164

KOI8-R Mac Cyrillic KOI8-R Mac Cyrillic				
KOI8-R		KOI8-R	Mac Cyrillic	
244	271	325	363	
245	317	326	346	
246	264	327	342	
247	273	330	374	
250	300	331	373	
251	275	332	347	
252	277	333	370	
253	40	334	375	
254	316	335	371	
255	40	336	367	
256	331	337	372	
257	333	340	236	
260	334	341	200	
261	253	342	201	
262	256	343	226	
263	335	344	204	
264	270	345	205	
265	301	346	224	
266	247	347	203	
267	272	350	225	
270	267	351	210	
271	274	352	211	
272	276	353	212	
273	40	354	213	
274	315	355	214	
275	40	356	215	
276	330	357	216	
277	332	360	217	

		Convers	sions Performed	
	KOI8-R	Mac Cyrillic	KOI8-R	Mac Cyrillic
	300	376	361	237
	301	340	362	220
	302	341	363	221
	303	366	364	222
	304	344	365	223
	305	345	366	206
	306	364	367	202
	307	343	370	234
	310	365	371	233
	311	350	372	207
	312	351	373	230
	313	352	374	235
	314	353	375	231
	315	354	376	227
	316	355		
FILES	/usr/lib/iconv/*	.so	conversion modules	5
	/usr/lib/iconv/*	t.t	conversion tables	
	/usr/lib/iconv/i	.conv_data	list of conversions st tables	upported by conversion
SEE ALSO	iconv(1), iconv(3C)),iconv(5)		

NAME | iconv_mac_cyr - code set conversion tables for Macintosh Cyrillic

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported					
Code	Symbol	Target Code	Symbol	Target Output	
Mac Cyrillic	mac	ISO 8859-5	iso5	ISO 8859-5 Cyrillic	
Mac Cyrillic	mac	KOI8-R	koi8	KOI8-R	
Mac Cyrillic	mac	PC Cyrillic	alt	Alternative PC Cyrillic	
Mac Cyrillic	mac	MS 1251	win5	Windows Cyrillic	

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

Mac Cyrillic to ISO 8859-5 For the conversion of Mac Cyrillic to ISO 8859-5, all characters not in the following table are mapped unchanged.

Conversions Performed				
Mac Cyrillic	ISO 8859-5	Mac Cyrillic	ISO 8859-5	
24	4	276	252	
200	260	277	372	
201	261	300	370	
202	262	301	245	
203	263	302-311	40	
204	264	312	240	
205	265	313	242	
206	266	314	362	
207	267	315	254	
210	270	316	374	
211	271	317	365	
212	272	320-327	40	
213	273	330	256	
214	274	331	376	

Mac Cyrillic ISO 8859-5 Mac Cyrillic ISO 8859-					
215	275	332	257		
216	276	333	377		
217	277	334	360		
220	300	335	241		
221	301	336	361		
222	302	337	357		
223	303	340	320		
224	304	341	321		
225	305	342	322		
226	306	343	323		
227	307	344	324		
230	310	345	325		
231	311	346	326		
232	312	347	327		
233	313	350	330		
234	314	351	331		
235	315	352	332		
236	316	353	333		
237	317	354	334		
240-246	40	355	335		
247	246	356	336		
250-252	40	357	337		
253	242	360	340		
254	362	361	341		
255	40	362	342		
256	243	363	343		
257	363	364	344		
260-263	40	365	345		

Conversions Performed			
Mac Cyrillic	ISO 8859-5	Mac Cyrillic	ISO 8859-5
264	366	366	346
265-266	40	367	347
267	250	370	350
270	244	371	351
271	364	372	352
272	247	373	353
273	367	374	354
274	251	375	355
275	371	376	356
375	370		

Mac Cyrillic to KOI8-R

For the conversion of Mac Cyrillic to KOI8-R, all characters not in the following table are mapped unchanged.

Conversions Performed			
Mac Cyrillic	KOI8-R	Mac Cyrillic	KOI8-R
24	4	276	272
200	341	277	252
201	342	300	250
202	367	301	265
203	347	302-311	40
204	344	312	240
205	345	313	261
206	366	314	241
207	372	315	274
210	351	316	254
211	352	317	245
212	353	320-327	40
213	354	330	276

Mac Cyrillic KOI8-R Mac Cyrillic KOI8-					
214	355	331	256		
215	356	332	277		
216	357	333	257		
217	360	334	260		
220	362	335	263		
221	363	336	243		
222	364	337	321		
223	365	340	301		
224	346	341	302		
225	350	342	327		
226	343	343	307		
227	376	344	304		
230	373	345	305		
231	375	346	326		
232	377	347	332		
233	371	350	311		
234	370	351	312		
235	374	352	313		
236	340	353	314		
237	361	354	315		
240-246	40	355	316		
247	266	356	317		
250-252	40	357	320		
253	261	360	322		
254	241	361	323		
255	40	362	324		
256	262	363	325		
257	242	364	306		

Conversions Performed			
Mac Cyrillic	KOI8-R	Mac Cyrillic	KOI8-R
260-263	40	365	310
264	246	366	303
265-266	40	367	336
267	270	370	333
270	264	371	335
271	244	372	337
272	267	373	331
273	247	374	330
274	271	375	334
275	251	376	300
375	370		

Mac Cyrillic to PC Cyrillic

For the conversion of Mac Cyrillic to PC Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed				
Mac Cyrillic	PC Cyrillic	Mac Cyrillic	PC Cyrillic	
24	4	355	255	
240-334	40	356	256	
335	360	357	257	
336	361	360	340	
337	357	361	341	
340	240	362	342	
341	241	363	343	
342	242	364	344	
343	243	365	345	
344	244	366	346	
345	245	367	347	
346	246	370	350	

Conversions Performed					
Mac Cyrillic	PC Cyrillic	Mac Cyrillic	PC Cyrillic		
347	247	371	351		
350	250	372	352		
351	251	373	353		
352	252	374	354		
353	253	375	355		
354	254	376	356		
303	366				

Mac Cyrillic to MS 1251 For the conversion of Mac Cyrillic to MS 1251, all characters not in the following table are mapped unchanged.

	Convers	sions Performed	
Mac Cyrillic	MS 1251	MS 1251 Mac Cyrillic	
24	4	255	40
200	300	256	201
201	301	257	203
202	302	260-263	40
203	303	264	263
204	304	266	264
205	305	267	243
206	306	270	252
207	307	271	272
210	310	272	257
211	311	273	277
212	312	274	212
213	313	275	232
214	314	276	214
215	315	277	234
216	316	300	274

Mac Cyrillic	MS 1251	Mac Cyrillic	MS 1251
.17	317	301	275
220	320	302	254
221	321	303-306	40
222	322	307	253
23	323	310	273
24	324	311	205
25	325	312	240
26	326	313	200
27	327	314	220
30	330	315	215
31	331	316	235
32	332	317	276
33	333	320	226
34	334	321	227
35	335	322	223
36	336	323	224
37	337	324	221
40	206	325	222
41	260	326	40
42	245	327	204
43	40	330	241
44	247	331	242
45	267	332	217
46	266	333	237
47	262	334	271
50	256	335	250
52	231	336	270
53	200	337	377

	Conversions Performed					
	Mac Cyrillic	MS 1251	Mac Cyrillic	MS 1251		
	254	220	362	324		
FILES	/usr/lib/iconv/*		conversion modules			
	/usr/lib/iconv/*		conversion tables			
	/usr/lib/iconv/i	.conv_data	list of conversions sup tables	ported by conversion		
SEE ALSO	iconv(1), iconv(3C)	i conv(5)				
JEE ALSO	100110(1), 100110(30)	, 100110(3)				

NAME | iconv_maz – code set conversion tables for Mazovia

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported						
Code	Symbol	Target Code	Symbol	Target Output		
Mazovia	maz	ISO 8859-2	iso2	ISO Latin 2		
Mazovia	maz	MS 1250	win2	Windows Latin 2		
Mazovia	maz	MS 852	dos2	MS-DOS Latin 2		
Mazovia	maz	DHN	dhn	Dom Hanlowy Nauki		

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

Mazovia to ISO 8859-2 For the conversion of Mazovia to ISO 8859-2, all characters not in the following table are mapped unchanged.

Conversions Performed			
Mazovia	ISO 8859-2	Mazovia	ISO 8859-2
24–177	40	230	246
200	307	231	326
201	374	232	334
202	351	233	40
203	342	234	243
204	344	235	40
205	40	236	266
206	261	237	40
207	347	240	254
210	40	241	257
211	353	242	363
212-213	40	243	323
214	356	244	361
215	346	245	321

iconv_maz(5)

Conversions Performed			
Mazovia	ISO 8859-2	Mazovia	ISO 8859-2
216	304	246	274
217	241	247	277
220	312	250-340	40
221	352	341	337
222	263	342-365	40
223	364	366	367
224	366	367	40
225	306	370	260
226-227	40	371-376	40
256	201		

Mazovia to MS 1250

For the conversion of Mazovia to MS 1250, all characters not in the following table are mapped unchanged.

Mazovia	MS 1250	Mazovia	MS 1250
200	307	236	234
201	374	237	40
202	351	240	217
203	342	241	257
204	344	242	363
205	40	243	323
206	271	244	361
207	347	245	321
210	40	246	237
211	353	247	277
212-213	40	250-251	40
214	356	252	254
215	346	253-255	40
216	304	256	253

iconv_maz(5)

Mazovia	MS 1250	Mazovia	MS 1250
217	245	257	273
220	312	260-340	40
221	352	341	337
222	263	342-345	40
223	364	346	265
224	366	347-360	40
225	306	361	261
226-227	40	362-365	0
230	214	366	367
231	326	367	40
232	334	370	260
233	40	371	40
234	243	372	267
235	40	373-376	40
274	212		

Mazovia to MS 852

For the conversion of Mazovia to MS 852, all characters not in the following table are mapped unchanged.

Conversions Performed			
Mazovia	MS 852	Mazovia	MS 852
205	40	234	235
206	245	235	40
210-213	40	236	230
215	206	237	40
217	244	240	215
220	250	241	275
221	251	243	340
222	210	244	344
225	217	245	343

iconv_maz(5)

Conversions Performed			
Mazovia	MS 852	Mazovia	MS 852
226-227	40	246	253
230	227	247	276
233	40	250-375	40
227	327		

Mazovia to DHN For the conversion of Mazovia to DHN, all characters not in the following table are mapped unchanged.

Conversions Performed			
Mazovia	DHN	Mazovia	DHN
200-205	40	234	203
206	211	236	217
207-214	40	240	207
215	212	241	210
216	40	242	216
217	200	243	205
220	202	244	215
221	214	246	220
225	201	247	221
230	206		

FILES	/usr/lib/iconv/*.so	conversion modules
	/usr/lib/iconv/*.t	conversion tables
	/usr/lib/iconv/iconv_data	list of conversions supported by conversion tables
SEE ALSO	<pre>iconv(1), iconv(3C), iconv(5)</pre>	

NAME | iconv_pc_cyr – code set conversion tables for Alternative PC Cyrillic

DESCRIPTION

The following code set conversions are supported:

Code Set Conversions Supported				
Code	Symbol	Target Code	Symbol	Target Output
PC Cyrillic	alt	ISO 8859-5	iso5	ISO 8859-5 Cyrillic
PC Cyrillic	alt	KOI8-R	koi8	KOI8-R
PC Cyrillic	alt	MS 1251	win5	Windows Cyrillic
PC Cyrillic	alt	Mac Cyrillic	mac	Macintosh Cyrillic

CONVERSIONS

The conversions are performed according to the following tables. All values in the tables are given in octal.

PC Cyrillic to ISO 8859-5

For the conversion of PC Cyrillic to ISO 8859-5, all characters not in the following table are mapped unchanged.

	Conversions Performed		
PC Cyrillic	ISO 8859-5	PC Cyrillic	ISO 8859-5
24	4	231	311
200	260	232	312
201	261	233	313
202	262	234	314
203	263	235	315
204	264	236	316
205	265	237	317
206	266	240	320
207	267	241	321
210	270	242	322
211	271	243	323
212	272	244	324
213	273	245	325
214	274	246	326

Conversions Performed			
PC Cyrillic	ISO 8859-5	PC Cyrillic	ISO 8859-5
215	275	247	327
216	276	250	330
217	277	251	331
220	300	252	332
221	301	253	333
222	302	254	334
223	303	255	335
224	304	256	336
225	305	257	337
226	306	260-337	255
227	307	360	241
230	310	362-376	255

PC Cyrillic to KOI8-R

For the conversion of PC Cyrillic to KOI8-R, all characters not in the following table are mapped unchanged.

Conversions Performed				
PC Cyrillic	KOI8-R	PC Cyrillic	KOI8-R	
24	4	242	327	
200	341	243	307	
201	342	244	304	
202	367	245	305	
203	347	246	326	
204	344	247	332	
205	345	250	311	
206	366	251	312	
207	372	252	313	
210	351	253	314	
211	352	254	315	

¹⁸⁰ man pages section 5: Standards, Environments, and Macros • Last Revised 18 Apr 1997

Conversions Performed			
PC Cyrillic	KOI8-R	PC Cyrillic	KOI8-R
212	353	255	316
213	354	256	317
214	355	257	320
215	356	260-337	255
216	357	340	322
217	360	341	323
220	362	342	324
221	363	343	325
222	364	344	306
223	365	345	310
224	346	346	303
225	350	347	336
226	343	350	333
227	376	351	335
230	373	352	337
231	375	353	331
232	377	354	330
233	371	355	334
234	370	356	300
235	374	357	321
236	340	360	263
237	361	361	243
240	301	362-376	255
241	302		

PC Cyrillic to MS 1251

For the conversion of PC Cyrillic to MS 1251, all characters not in the following table are mapped unchanged.

PC Cyrillic	MS 1251	PC Cyrillic	MS 1251
24	4	242	342
200	300	243	343
201	301	244	344
202	302	245	345
203	303	246	346
204	304	247	347
205	305	250	350
206	306	251	351
207	307	252	352
210	310	253	353
211	311	254	354
212	312	255	355
213	313	256	356
214	314	257	357
215	315	260-337	210
216	316	340	360
217	317	341	361
220	320	342	362
221	321	343	363
222	322	344	364
223	323	345	365
224	324	346	366
225	325	347	367
226	326	350	370
227	327	351	371
230	330	352	372
231	331	353	373
232	332	354	374

182 man pages section 5: Standards, Environments, and Macros • Last Revised 18 Apr 1997

Conversions Performed				
PC Cyrillic	MS 1251	PC Cyrillic	MS 1251	
233	333	355	375	
234	334	356	376	
235	335	357	377	
236	336	360	250	
237	337	361	270	
240	340	362-376	210	
241	341			

PC Cyrillic to Mac Cyrillic

For the conversion of PC Cyrillic to Mac Cyrillic, all characters not in the following table are mapped unchanged.

Conversions Performed				
PC Cyrillic	Mac Cyrillic	PC Cyrillic	Mac Cyrillic	
24	4	341	361	
240	340	342	362	
241	341	343	363	
242	342	344	364	
243	343	345	365	
244	344	346	366	
245	345	347	367	
246	346	350	370	
247	347	351	371	
250	350	352	372	
251	351	353	373	
252	352	354	374	
253	353	355	375	
254	354	356	376	
255	355	357	337	
256	356	360	335	

	Conversions Performed				
	PC Cyrillic	Mac Cyrillic	PC Cyrillic	Mac Cyrillic	
	257	357	361	336	
	260-337	40	362-376	40	
	340	360			
FILES	/usr/lib/iconv/	* 50	conversion modules		
FILLS	/usr/lib/iconv/		conversion tables		
	/usr/lib/iconv/			ported by conversion	
SEE ALSO	iconv(1), iconv(3C), iconv(5)			

NAME

E | iconv_unicode – code set conversion tables for Unicode

DESCRIPTION

The following code set conversions are supported:

CODE SET CONVERSIONS SUPPORTED

FROM Code Set			TO Code	Set	
Code	FROM	Tar	get Code	то	
	Filename			Filename	
	Element			Element	
ISO 8859-1 (Latin			UTF-8	UTF-8	
ISO 8859-2 (Latin ISO 8859-3 (Latin	2) 8859-2		UTF-8	UTF-8	
ISO 8859-3 (Latin	3) 8859-3		UTF-8	UTF-8	
ISO 8859-4 (Latin	4) 8859-4		UTF-8	UTF-8	
ISO 8859-4 (Latin ISO 8859-5 (Cyril)	ic) 8859-5		UTF-8	UTF-8	
ISO 8859-6 (Arabic	2) 8859-6		UTF-8	UTF-8	
ISO 8859-7 (Greek)	8859-7		UTF-8	UTF-8	
ISO 8859-8 (Hebrew	ı) 8859-8		UTF-8	UTF-8	
ISO 8859-9 (Latin ISO 8859-10 (Latir	5) 8859-9		UTF-8	UTF-8	
ISO 8859-10 (Latir	n 6) 8859-10		UTF-8	UTF-8	
Japanese EUC	eucJP		UTF-8	UTF-8	
Chinese/PRC EUC					
(GB 2312-1980)	gb2312		UTF-8	UTF-8	
ISO-2022	iso2022		UTF-8	UTF-8	
Korean EUC			Korean UTF-		
ISO-2022-KR	ko_KR-is	02022-7	Korean UTF-	8 ko_KR_UTF-8	
Korean Johap					
(KS C 5601-1987)	ko_KR-jo	hap	Korean UTF-	8 ko_KR-UTF-8	
Korean Johap					
(KS C 5601-1992)	ko_KR-jo		Korean UTF-		
Korean UTF-8	ko_KR-UT	F-8	Korean EUC Korean Joha	ko_KR-euc	
Korean UTF-8	ko_KR-UT				
			(KS C 5601-		
Korean UTF-8	ko_KR-UT	'F - 8	Korean Joha		3
			(KS C 5601-		
KOI8-R (Cyrillic)	KO18-R		UCS-2	UCS-2	
KOI8-R (Cyrillic) PC Kanji (SJIS)	KOI8-R		UTF-8	UTF-8	
PC Kanji (SJIS)	PCK		UTF-8	UTF-8	
PC Kanji (SJIS)	SJIS		UTF-8	UTF-8	
UCS-2	UCS-2		-	illic) KOI8-R	
UCS-2	UCS-2		UCS-4	UCS-4	
	CODE SET CON				
FROM Code Set			TO Code		
Code	FROM	Targ	get Code	TO	
	Filename			Filename	
	Element			Element	
UCS-2	UCS-2	UTF-7		UTF-7	
UCS-2	UCS-2	UTF-8		UTF-8	
UCS-4	UCS-4	UCS-2		UCS-2	
UCS-4	UCS-4	UTF-16		UTF-16	
UCS-4	UCS-4	UTF - 7		UTF - 7	
UCS-4	UCS-4	UTF-8		UTF-8	
UTF-16	UTF-16	UCS-4		UCS-4	

iconv_unicode(5)

_ ()				
	UTF-16	UTF-16	UTF-8	UTF-8
	UTF-7	UTF-7	UCS-2	UCS-2
	UTF-7	UTF-7	UCS-4	UCS-4
	UTF-7	UTF-7	UTF-8	UTF-8
	UTF-8	UTF-8	ISO 8859-1 (Latin 1)	8859-1
	UTF-8	UTF-8	ISO 8859-2 (Latin 2)	8859-2
	UTF-8	UTF-8	ISO 8859-3 (Latin 3)	8859-3
	UTF-8	UTF-8	ISO 8859-4 (Latin 4)	8859-4
	UTF-8	UTF-8	ISO 8859-5 (Cyrillic)	8859-5
	UTF-8	UTF-8	ISO 8859-6 (Arabic)	8859-6
	UTF-8	UTF-8	ISO 8859-7 (Greek)	8859-7
	UTF-8	UTF-8	ISO 8859-8 (Hebrew)	8859-8
	UTF-8	UTF-8	ISO 8859-9 (Latin 5)	8859-9
	UTF-8	UTF-8	ISO 8859-10 (Latin 6)	8859-10
	UTF-8	UTF-8	Japanese EUC	eucJP
	UTF-8	UTF-8	Chinese/PRC EUC	gb2312
			(GB 2312-1980)	
	UTF-8	UTF-8	ISO-2022	iso2022
	UTF-8	UTF-8	KOI8-R (Cyrillic)	KOI8-R
	UTF-8	UTF-8	PC Kanji (SJIS)	PCK
	UTF-8	UTF-8	PC Kanji (SJIS)	SJIS
	UTF-8	UTF-8	UCS-2	UCS-2
	UTF-8	UTF-8	UCS-4	UCS-4
	UTF-8	UTF-8	UTF-16	UTF-16
	UTF-8	UTF-8	UTF-7	UTF-7
	UTF-8	UTF-8	Chinese/PRC EUC	zh_CN.euc
			(GB 2312-1980)	
		CODE SET CONVE	RSIONS SUPPORTED	
	FROM Code Set		TO Code Set	
	Code	FROM	Target Code	ТО
		Filename		Filename
		Element		Element
	UTF-8	UTF-8	ISO 2022-CN	zh CN.iso2022-7
	UTF-8	UTF-8	Chinese/Taiwan Big	_
	UTF-8	UTF-8	Chinese/Taiwan EU	
			(CNS 11643-1992)	—
	UTF-8	UTF-8	ISO 2022-TW	zh TW-iso2022-7
	Chinese/PRC EUC	zh CN.euc	UTF-8	
	(GB 2312-1980)	-		
	ISO 2022-CN	zh_CN.iso2022	2-7 UTF-8	UTF-8
	Chinese/Taiwan Big!	5 zh_TW-big5	UTF-8	UTF-8
	Chinese/Taiwan EU	C zh_TW-euc	UTF-8	UTF-8
	(CNS 11643-1992)	-		
	ISO 2022-TW	zh_TW-iso2022	2-7 UTF-8	UTF-8
EXAMPLES	EXAMPLE 1 The library	y module filename	2	
		,		

In the conversion library, /usr/lib/iconv (see iconv(3C)), the library module filename is composed of two symbolic elements separated by the percent sign (%). The first symbol specifies the code set that is being converted; the second symbol specifies the *target code*, that is, the code set to which the first one is being converted.

¹⁸⁶ man pages section 5: Standards, Environments, and Macros • Last Revised 18 Apr 1997

EXAMPLE 1 The library module filename (<i>Continued</i>)					
In the conversion table above, the first symbol is termed the "FROM Filename Element". The second symbol, representing the target code set, is the "TO Filename Element".					
For example, the library module filename to convert from the <i>Korean EUC</i> code set to the <i>Korean UTF-8</i> code set is					
ko_KR-euc%ko_KR-UTF-8					
/usr/lib/iconv/*.so conversion modules					
<pre>iconv(1), iconv(3C), iconv(5)</pre>					
Chernov, A., <i>Registration of a Cyrillic Character Set</i> , RFC 1489, RELCOM Development Team, July 1993.					
Chon, K., H. Je Park, and U. Choi, <i>Korean Character Encoding for Internet Messages</i> , RFC 1557, Solvit Chosun Media, December 1993.					
Goldsmith, D., and M. Davis, <i>UTF-7 – A Mail-Safe Transformation Format of Unicode</i> , RFC 1642, Taligent, Inc., July 1994.					
Lee, F., <i>HZ – A Data Format for Exchanging Files of Arbitrarily Mixed Chinese and ASCII characters</i> , RFC 1843, Stanford University, August 1995.					
Murai, J., M. Crispin, and E. van der Poel, <i>Japanese Character Encoding for Internet Messages</i> , RFC 1468, Keio University, Panda Programming, June 1993.					
Nussbacher, H., and Y. Bourvine, <i>Hebrew Character Encoding for Internet Messages</i> , RFC 1555, Israeli Inter-University, Hebrew University, December 1993.					
Ohta, M., <i>Character Sets ISO-10646 and ISO-10646-J-1</i> , RFC 1815, Tokyo Institute of Technology, July 1995.					
Ohta, M., and K. Handa, ISO-2022-JP-2: Multilingual Extension of ISO-2022-JP, RFC 1554, Tokyo Institute of Technology, December 1993.					
Reynolds, J., and J. Postel, ASSIGNED NUMBERS, RFC 1700, University of Southern California/Information Sciences Institute, October 1994.					
Simonson, K., Character Mnemonics & Character Sets, RFC 1345, Rationel Almen Planlaegning, June 1992.					
Spinellis, D., Greek Character Encoding for Electronic Mail Messages, RFC 1947, SENA S.A., May 1996.					

iconv_unicode(5)						
		The Unicode Consortium, <i>The Unicode Standard</i> , Version 2.0, Addison Wesley Developers Press, July 1996.				
	Wei, Y., Y. Zhang, J. Li, J. Ding, and Y. Jiang, <i>ASCII Printable Characters-Based Chinese Character Encoding for Internet Messages</i> , RFC 1842, AsiaInfo Services Inc., Harvard University, Rice University, University of Maryland, August 1995.					
	Yergeau, F., UTF-8, a transformation format of Unicode and ISO 10646, RFC 2044, Alis Technologies, October 1996.					
	for Internet Messages, RFC	1922, Tsinghua Üniversi Committee (CITS), Inst	A. Crispin, <i>Chinese Character Encoding</i> ty, China Information Technology itute for Information Industry (III),			
NOTES	ISO 8859 character sets us	ing Latin alphabetic cha	racters are distinguished as follows:			
	ISO 8859-1 (Latin 1) For most West European languages, including:					
	Albanian	Finnish	Italian			
	Catalan	French	Norwegian			
	Danish	German	Portuguese			
	Dutch	Galician	Spanish			
	English	Irish	Swedish			
	Faeroese	Icelandic				
	ISO 8859-2 (Latin 2) For most Latin-written Slavic and Central European languages:					
	Czech	Polish	Slovak			
	German	Rumanian	Slovene			
	Hungarian	Croatian				
	ISO 8859-3 (Latin 3) Popularly used for Esperanto, Galician, Maltese, and Turkish.					
	ISO 8859-4 (Latin 4) Introduces letters for Estonian, Latvian, and Lithuanian. It is an incomplete predecessor of ISO 8859-10 (Latin 6).					

iconv_unicode(5)

ISO 8859-9 (Latin 5)

Replaces the rarely needed Icelandic letters in ISO 8859-1 (Latin 1) with the Turkish ones.

ISO 8859-10 (Latin 6)

Adds the last Inuit (Greenlandic) and Sami (Lappish) letters that were not included in ISO 8859-4 (Latin 4) to complete coverage of the Nordic area.

isalist(5)

NAME	isalist – the native instruction sets known to Solaris software				
DESCRIPTION	The possible instruction set names returned by isalist(1) and the SI_ISALIST command of sysinfo(2) are listed here.				
	The list is ordered within an instruction set family in the sense that later name generally faster then earlier names; note that this is in the reverse order than l isalist(1) and sysinfo(2). In the following list of values, numbered entrie generally represent increasing performance; lettered entries are either mutual exclusive or cannot be ordered.				
SPARC Platforms		Where appropriate, correspondence with a given value of the –xarch option of Sun's C .0 compiler is indicated. Other compilers may have similar options.			
	1a. sparc	Indicates the SPARC V8 instruction set, as defined in The SPARC Architecture Manual, Version 8, Prentice-Hall, Inc., 1992. Some instructions (such as integer multiply and divide, FSMULD, and all floating point operations on quad operands) may be emulated by the kernel on certain systems.			
	1b. sparcv7	Same as sparc. This corresponds to code produced with the -xarch=v7 option of Sun's C 4.0 compiler.			
	2. sparcv8–fsmuld	Like sparc, except that integer multiply and divide must be executed in hardware. This corresponds to code produced with the –xarch=v8a option of Sun's C 4.0 compiler.			
	3. sparcv8	Like sparcv8–fsmuld, except that FSMULD must also be executed in hardware. This corresponds to code produced with the –xarch=v8 option of Sun's C 4.0 compiler.			
	4. sparcv8plus	Indicates the SPARC V8 instruction set plus those instructions in the SPARC V9 instruction set, as defined in The SPARC Architecture Manual, Version 9, Prentice-Hall, 1994, that can be used according to The V8+ Technical Specification. This corresponds to code produced with the –xarch=v8plus option of Sun's C 4.0 compiler.			
	5a. sparcv8plus+vis	Like sparcv8plus, with the addition of those UltraSPARC I Visualization Instructions that can be used according to The V8+ Technical Specification. This corresponds to code produced with the -xarch=v8plusa option of Sun's C 4.0 compiler.			
	5b. sparcv8plus+fmuladd	Like sparcv8plus, with the addition of the Hal SPARC64 floating multiply-add and multiply-subtract instructions.			

	6. sparcv9	Indicates the SPARC V9 instruction set, as defined in The SPARC Architecture Manual, Version 9, Prentice-Hall, 1994.
	7a. sparcv9+vis	Like sparcv9, with the addition of the UltraSPARC I Visualization Instructions.
	7b. sparcv9+fmuladd	Like sparcv9, with the addition of the Hal SPARC64 floating multiply-add and multiply-subtract instructions.
x86 Platforms	1. i386	The Intel 80386 instruction set, as described in The i386 Microprocessor Programmer's Reference Manual.
	2. i486	The Intel 80486 instruction set, as described in The i486 Microprocessor Programmer's Reference Manual. (This is effectively i386, plus the CMPXCHG, BSWAP, and XADD instructions.)
	3. pentium	The Intel Pentium instruction set, as described in The Pentium Processor User's Manual. (This is effectively i486, plus the CPU_ID instruction, and any features that the CPU_ID instruction indicates are present.)
	4. pentium+mmx	Like pentium, with the MMX instructions guaranteed present.
	5. pentium_pro	The Intel PentiumPro instruction set, as described in The PentiumPro Family Developer's Manual. (This is effectively pentium, with the CMOVcc, FCMOVcc, FCOMI, and RDPMC instructions guaranteed present.)
	6. pentium_pro+mmx	Like pentium_pro, with the MMX instructions guaranteed present.
SEE ALSO	<pre>isalist(1), sysinfo(2)</pre>	

largefile(5)

largefile – large file status of utilities NAME |

DESCRIPTION

A *large file* is a regular file whose size is greater than or equal to 2 Gbyte (2^{31} bytes). A *small file* is a regular file whose size is less than 2 Gbyte.

Large file aware A utility is called *large file aware* if it can process large files in the same manner as it utilities does small files. A utility that is large file aware is able to handle large files as input and generate as output large files that are being processed. The exception is where additional files are used as system configuration files or support files that can augment the processing. For example, the file utility supports the -m option for an alternative "magic" file and the -f option for a support file that can contain a list of file names. It is unspecified whether a utility that is large file aware will accept configuration or support files that are large files. If a large file aware utility does not accept configuration or support files that are large files, it will cause no data loss or corruption upon encountering such files and will return an appropriate error.

The following /usr/bin utilities are large file aware:

adb	awk	bdiff	cat	chgrp	
chmod	chown	cksum	cmp	compress	
ср	csh	csplit	cut	dd	
dircmp	du	egrep	fgrep	file	
find	ftp	getconf	grep	head	
join	jsh	ksh	ln	ls	
mdb	mkdir	mkfifo	more	mv	
nawk	page	paste	pathchk	pg	
rcp	remsh	rksh	rm	rmdir	
rsh	sed	sh	sort	split	
sum	tail	tar	tee	test	
touch	tr	uncompress	uudecode	uuencode	
WC	zcat				
The following /usr/xpg4/bin utilities are large file aware:					
awk	ср	du	egrep	fgrep	

	_			
grep	ln	ls	more	mv
rm	sed	sh	sort	tail

largefile(5)

	tr						
	The following /usr/sbin utilities are large file aware:						
	install	mkfile	mknod	mv	dir	swap	
	See the USAGE section of the swap(1M) manual page for limitations of swap on block devices greater than 2 Gbyte on a 32-bit operating system.						
	The following /usr/ucb utilities are large file aware:						
	chown	from	ln	ls		sed	
	sum	touch					
	The /usr/bin/ archive a file wh				large file	aware, but can	not
The /usr/bin/truss utilities has been modified to read a dump file and o information relevant to large files, such as offsets.				p file and disp	lay		
cachefs file systems The following /usr/bin utilities are large file aware for cachefs file systems							
	cachefspack	cachefss	tat				
	The following /1	usr/sbin utilit	ies are larg	ge file aware	for cache	fs file system	s:
	cachefslog	cachefsw	ssize	cfsadmin		fsck	
	mount	umount					
file systems The following utilities are large file aware for nfs file systems:							
	/usr/lib/auto	fs/automountd		/usr/sbin	/mount		
	/usr/lib/nfs/	rquotad					
file systems The following /usr/bin utility is large file aware for ufs file systems:				stems:			
	df						

nfs

ufs

largefile(5	5)
-------------	----

	The following /usr/lib/nfs utility is large file aware for ufs file systems:					
	rquotad					
	The following /usr/xpg4/bin utility is large file aware for ufs file systems:					
	df					
	The following /usr/sbin utilities are large file aware for ufs file systems:					
	clri	dcopy	edquota	ff	fsck	
	fsdb	fsirand	fstyp	labelit	lockfs	
	mkfs	mount	ncheck	newfs	quot	
	quota	quotacheck	quotaoff	quotaon	repquota	
	tunefs	ufsdump	ufsrestore	umount		
Large file safe utilities	A utility is called <i>large file safe</i> if it causes no data loss or corruption when it encounters a large file. A utility that is large file safe is unable to process properly a large file, but returns an appropriate error. The following /usr/bin utilities are large file safe:					
	audioconvert	audioplay	audiorecor	d comm	diff	
	diff3	diffmk	ed	lp	mail	
	mailcompat	mailstats	mailx	pack	pcat	
	red	rmail	sdiff	unpack	vi	
	view					
	The following /usr/xpg4/bin utilities are large file safe:					
	ed	vi	view			
	The following /usr/sbin utilities are large file safe:					
lpfilter lpforms						
	The following /usr/ucb utilities are large file safe:					

194 man pages section 5: Standards, Environments, and Macros • Last Revised 16 Jul 2003

largefile(5)

	Mail	lpr
	The following /u	sr/lib utility is large file safe:
	sendmail	
SEE ALSO	lf64(5),lfcomp	ile(5),lfcompile64(5)

NAME	lf64 – transitional interfaces for 64-bit file offsets			
DESCRIPTION	The data types, interfaces, and macros described on this page provide explicit access to 64-bit file offsets. They are accessible through the transitional compilation environment described on the lfcompile64(5) manual page. The function prototype and semantics of a transitional interface are equivalent to those of the standard version of the call, except that relevant data types are 64-bit entities.			
Data Types	The following tables list the standard data or struct types in the left-hand column and their corresponding explicit 64-bit file offset types in the right-hand column, grouped by header. The absence of an entry in the left-hand column indicates that there is no existing explicit 32-bit type that corresponds to the 64–bit type listed in the right—hand column. Note that in a 64-bit application, the standard definition is equivalent to the 64-bit file offset definition.			
	<aio.h></aio.h>			
	struct a joch			
	struct aiocb struct aiocb64 off t aio_offset; off64 t aio_offset;			
	<sys dirent.h=""></sys>			
	struct dirent	struct dirent64		
	ino_t d_ino;	ino64_t d_ino;		
	off_t d_off;			
	<sys fcntl.h=""></sys>			
	struct flock	struct flock64		
	off_t l_start;	off64_t l_start;		
	off_t l_len;	off64_t l_len;		
	F_SETLK	F_SETLK64		
	F_SETLKW F_SETLKW64			
	F_GETLK F_GETLK64			
	F_FREESP F_FREESP64			
	O_LARGEFILE			
	<sys stdio.h=""></sys>			
I				

fpos64_t

<sys/resource.h>

fpos_t

rlim_t	rlim64_t
struct rlimit	struct rlimit64
<pre>rlim_t rlim_cur;</pre>	rlim64_t rlim_cur;
<pre>rlim_t rlim_max;</pre>	rlim64_t rlim_max;
RLIM_INFINITY	RLIM64_INFINITY
RLIM_SAVED_MAX	RLIM64_SAVED_MAX
RLIM_SAVED_CUR	RLIM64_SAVED_CUR

<sys/stat.h>

struct stat	struct stat64
<pre>ino_t st_ino;</pre>	<pre>ino64_t st_ino;</pre>
off_t st_size;	off64_t st_size;
blkcnt_t st_blocks;	blkcnt64_t st_blocks;

<sys/statvfs.h>

struct statvfs	struct statvfs64
fsblkcnt_t f_blocks;	<pre>fsblkcnt64_t f_blocks;</pre>
fsblkcnt_t f_bfree;	fsblkcnt64_t f_bfree;
fsblkcnt_t f_bavial;	<pre>fsblkcnt64_t f_bavial;</pre>
fsfilcnt_t f_files;	fsfilcnt64_t f_files;
<pre>fsfilcnt_t f_ffree;</pre>	fsfilcnt64_t f_ffree;
fsfilcnt_t f_favail;	<pre>fsfilcnt64_t f_favail;</pre>

<sys/types.h>

off_t;

off64_t;

Standards, Environments, and Macros 197

lf64(5)

	ino_t;	<pre>ino64_t;</pre>
	blkcnt_t;	blkcnt64_t;
	fsblkcnt_t;	fsblkcnt64_t;
	fsfilcnt_t;	fsfilcnt64_t;
	<unistd.h></unistd.h>	
		_LFS64_LARGEFILE
		_LFS64_STDIO
	<sys unistd.h=""></sys>	
		_CS_LFS64_CFLAGS
		_CS_LFS64_LDFLAGS
		_CS_LFS64_LIBS
		_CS_LFS64_LINTFLAGS
System Interfaces	The following tables display the standard interfaces for 64-bit file offsets. The interfa name and the affected data types are displ	ces are grouped by header. The interface
System Interfaces	interfaces for 64-bit file offsets. The interfa	ces are grouped by header. The interface
System Interfaces	interfaces for 64-bit file offsets. The interfa name and the affected data types are displ	ces are grouped by header. The interface
System Interfaces	interfaces for 64-bit file offsets. The interfa name and the affected data types are displ <aio.h></aio.h>	ces are grouped by header. The interface ayed in courier font
System Interfaces	<pre>interfaces for 64-bit file offsets. The interfaces and the affected data types are displ <aio.h> int aio_cancel(,</aio.h></pre>	ces are grouped by header. The interface ayed in courier font int aio_cancel64(,
System Interfaces	<pre>interfaces for 64-bit file offsets. The interfaces name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *);</aio.h></pre>	ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *);
System Interfaces	<pre>interfaces for 64-bit file offsets. The interfaces name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *); int aio_error(</aio.h></pre>	<pre>ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *); int aio_error64(</pre>
System Interfaces	<pre>interfaces for 64-bit file offsets. The interfa name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *); int aio_error(const struct aiocb *);</aio.h></pre>	<pre>ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *); int aio_error64(const struct aiocb64 *);</pre>
System Interfaces	<pre>interfaces for 64-bit file offsets. The interfaces name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *); int aio_error(const struct aiocb *); int aio_fsync(,</aio.h></pre>	<pre>ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *); int aio_error64(const struct aiocb64 *); int aio_fsync64(,</pre>
System Interfaces	<pre>interfaces for 64-bit file offsets. The interface name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *); int aio_error(const struct aiocb *); int aio_fsync(, struct aiocb *);</aio.h></pre>	<pre>ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *); int aio_error64(const struct aiocb64 *); int aio_fsync64(, struct aiocb64 *);</pre>
System Interfaces	<pre>interfaces for 64-bit file offsets. The interface name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *); int aio_error(const struct aiocb *); int aio_fsync(, struct aiocb *); int aio_read(struct aiocb *);</aio.h></pre>	<pre>ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *); int aio_error64(const struct aiocb64 *); int aio_fsync64(, struct aiocb64 *); int aio_read64(struct aiocb64 *);</pre>
System Interfaces	<pre>interfaces for 64-bit file offsets. The interface name and the affected data types are displ <aio.h> int aio_cancel(, struct aiocb *); int aio_error(const struct aiocb *); int aio_fsync(, struct aiocb *); int aio_read(struct aiocb *); int aio_return(struct aiocb *);</aio.h></pre>	<pre>ces are grouped by header. The interface ayed in courier font int aio_cancel64(, struct aiocb64 *); int aio_error64(const struct aiocb64 *); int aio_fsync64(, struct aiocb64 *); int aio_read64(struct aiocb64 *); int aio_return64(struct aiocb64 *);</pre>

<pre>int aio_write(struct aiocb *);</pre>	<pre>int aio_write64(struct aiocb64 *);</pre>
<pre>int lio_listio(,</pre>	<pre>intlio_listio64(,</pre>
<pre>const struct aiocb *,);</pre>	const struct aiocb64 *,);

<dirent.h>

<pre>struct dirent *readdir();</pre>	<pre>struct dirent64 *readdir64();</pre>
<pre>struct dirent *readdir_r();</pre>	<pre>struct dirent64 *readdir64_r();</pre>

<fcntl.h>

int attropen64();
<pre>int creat64();</pre>
int open64();
int openat64();

<ftw.h>

int ftw(,	int ftw64(,
const struct stat *,	<pre>const struct stat64 *,</pre>
););
int nftw(int nftw64(,
const struct stat *,	<pre>const struct stat64 *,</pre>
););

<libgen.h>

char *copylist(..., off_t);

<stdio.h>

int fgetpos();

int fgetpos64();

off64_t);

char*copylist64(...,

<pre>FILE *fopen();</pre>	FILE *fopen64();	
<pre>FILE *freopen();</pre>	<pre>FILE *freopen64();</pre>	
int fseeko(,	int fseeko64(,	
off_t,);	off64_t,);	
int fsetpos(,	int fsetpos64(,	
<pre>const fpos_t *);</pre>	<pre>const fpos64_t *);</pre>	
off_t ftello();	off64_t ftello64();	
<pre>FILE *tmpfile();</pre>	<pre>FILE *tmpfile64();</pre>	
<stdlib.h></stdlib.h>		
<pre>int mkstemp();</pre>	<pre>int mkstemp64();</pre>	
<sys async.h=""></sys>		
<pre>int aioread(, off_t,</pre>	<pre>int aioread64(, off64_t,</pre>	
););	
<pre>int aiowrite(, off_t,</pre>	intaiowrite64(,	
);	off64_t,);	
<ucbinclude dir.h="" sys=""></ucbinclude>		
int alphasort(int alphasort64(
<pre>struct direct **,</pre>	<pre>struct direct64 **,</pre>	
<pre>struct direct **);</pre>	<pre>struct direct64 **);</pre>	

struct direct64 *readdir64();

int scandir64(...,

. . .);

struct direct64 *(*[]);,

struct direct *readdir()

int scandir(...,

<sys/dirent.h>

. . .);

struct direct *(*[]);,

int getdents(..., dirent);

int getdents64(...,
dirent64);

<sys/mman.h>

void mmap(..., off_t);

void mmap64(..., off64_t);

<sys/resource.h>

<pre>intgetrlimit(,</pre>	<pre>intgetrlimit64(,</pre>
<pre>struct rlimit *);</pre>	<pre>struct rlimit64 *);</pre>
<pre>int setrlimit(,</pre>	int setrlimit64(,
<pre>const struct rlimit *);</pre>	<pre>const struct rlimit64 *);</pre>

<sys/stat.h>

int fstat(,	int fstat64(,
<pre>struct stat *);</pre>	<pre>struct stat64 *);</pre>
int fstatat(,	int fstatat64(,
<pre>struct stat *, int);</pre>	<pre>struct stat64 *, int);</pre>
<pre>intlstat(,</pre>	intlstat64(,
<pre>struct stat *);</pre>	<pre>struct stat64 *);</pre>
int stat(,	int stat64(,
<pre>struct stat *);</pre>	<pre>struct stat64 *);</pre>

<sys/statvfs.h>

int statvfs(,	int statvfs64(,
<pre>struct statvfs *);</pre>	<pre>struct statvfs64 *);</pre>
int fstatvfs(,	int fstatvfs64(,
<pre>struct statvfs *);</pre>	<pre>struct statvfs64 *);</pre>

<unistd.h>

	<pre>int lockf(, off_t);</pre>	intlockf64(,
		off64_t);
	off_t lseek(, off_t,	off64_t lseek64(,
);	off64_t,);
	<pre>int ftruncate(, off_t);</pre>	int ftruncate64(,
		off64_t);
	<pre>ssize_t pread(, off_t);</pre>	ssize_t pread64(,
		off64_t);
	<pre>ssize_t pwrite(, off_t);</pre>	ssize_t pwrite64(,
		off64_t);
	<pre>inttruncate(, off_t);</pre>	int truncate64(,
		off64_t);
SEE ALSO	lfcompile(5), lfcompile64(5)	

NAME	lfcompile –	large file com	pilation e	environment	for 32-bit applications

DESCRIPTION

All 64-bit applications can manipulate large files by default. The methods described on this page allow 32-bit applications to manipulate large files.

In the large file compilation environment, source interfaces are bound to appropriate 64-bit functions, structures, and types. Compiling in this environment allows 32-bit applications to access files whose size is greater than or equal to 2 Gbyte (2^{31} bytes).

Each interface named xxx() that needs to access 64-bit entities to access large files maps to a xxx64() call in the resulting binary. All relevant data types are defined to be of correct size (for example, off_t has a typedef definition for a 64-bit entity).

An application compiled in this environment is able to use the xxx() source interfaces to access both large and small files, rather than having to explicitly utilize the transitional xxx64() interface calls to access large files. See the lfcompile64(5) manual page for information regarding the transitional compilation environment.

Applications can be compiled in the large file compilation environment by using the following methods:

 Use the getconf(1) utility with one or more of the arguments listed in the table below. This method is recommended for portable applications.

argument	purpose	
LFS_CFLAGS	obtain compilation flags necessary to enable the large file compilation environment	
LFS_LDFLAGS	obtain link editor options	
LFS_LIBS	obtain link library names	
LFS_LINTFLAGS	obtain lint options	

 Set the compile-time flag _FILE_OFFSET_BITS to 64 before including any headers. Applications may combine objects produced in the large file compilation environment with objects produced in the transitional compilation environment, but must be careful with respect to interoperability between those objects. Applications should not declare global variables of types whose sizes change between compilation environments.

Access to Additional Large File Interfaces The fseek() and ftell() functions *do not* map to functions named fseek64() and ftell64(); rather, the large file additions fseeko() and ftello(), have functionality identical to fseek() and ftell() and *do* map to the 64-bit functions fseeko64() and ftello64(). Applications wishing to access large files should use fseeko() and ftello() in place of fseek() and ftell(). See the fseek(3C) and ftell(3C) manual pages for information about fseeko() and ftello().

lfcompile(5)

Applications wishing to access fseeko() and ftello() as well as the POSIX and X/Open specification-conforming interfaces should define the macro _LARGEFILE_SOURCE to be 1 and set whichever feature test macros are appropriate to obtain the desired environment (see standards(5)).

EXAMPLES In the following examples, the large file compilation environment is accessed by invoking the getconf utility with one of the arguments listed in the table above. The additional large file interfaces are accessed by specifying -D_LARGEFILE_SOURCE.

The examples that use the form of command substitution specifying the command within parentheses preceded by a dollar sign can be executed only in a POSIX-conforming shell such as the Korn Shell (see ksh(1)). In a shell that is not POSIX-conforming, such as the Bourne Shell (see sh(1)) and the C Shell (see csh(1)), the getconf calls must be enclosed within grave accent marks, as shown in the second example.

 $\mbox{EXAMPLE 1}$ An example of compiling a program with a "large" off_t, and that uses $\mbox{fseeko(),ftello(),and}\ \mbox{yacc}(1)$

```
$ c89 -D_LARGEFILE_SOURCE \
    -D_FILE_OFFSET_BITS=64 -o foo \
    $(getconf LFS_CFLAGS) y.tab.c b.o \
    $(getconf LFS_LDFLAGS) \
    -ly $(getconf LFS_LIBS)
```

EXAMPLE 2 An example of compiling a program with a "large" off_t that does not use fseeko() and ftello() and has no application specific libraries:

%	c89	-D FILE OFFSET BITS=64	
		`getconf LFS_CFLAGS` a.c	\
		`getconf LFS_LDFLAGS`	\
		`getconf LFS LIBS`	\

EXAMPLE 3 An example of compiling a program with a "default" off_t and that uses fseeko() and ftello():

```
$ c89 -D_LARGEFILE_SOURCE a.c
```

- SEE ALSO csh(1), getconf(1), ksh(1), lint(1B), sh(1), fseek(3C), ftell(3C), lf64(5), lfcompile64(5), standards(5)
 - **NOTES** Certain system-specific or non-portable interfaces are not usable in the large file compilation environment. Known cases are:
 - Kernel data structures read from /dev/kmem.
 - Interfaces in the kernel virtual memory library, -lkvm.
 - Interfaces in the ELF access library, -lelf.
 - Interfaces to /proc defined in <procfs.h>.
 - The ustat(2) system call.

Programs that use these interfaces should not be compiled in the large file compilation environment. As a partial safeguard against making this mistake, including either of the <libelf.h> or <sys/procfs.h> header files will induce a compilation error when the large file compilation environment is enabled.

204 man pages section 5: Standards, Environments, and Macros • Last Revised 23 Jul 2001

In general, caution should be exercised when using any separately-compiled library whose interfaces include data items of type off_t or the other redefined types either directly or indirectly, such as with 'struct stat'. (The redefined types are off_t, rlim_t, ino_t, blkcnt_t, fsblkcnt_t, and fsfilcnt_t.) For the large file compilation environment to work correctly with such a library, the library interfaces must include the appropriate xxx64 () binary entry points and must have them mapped to the corresponding primary functions when _FILE_OFFSET_BITS is set to 64.

Care should be exercised using any of the printf() or scanf() routines on variables of the types mentioned above. In the large file compilation environment, these variables should be printed or scanned using long long formats.

BUGS

The lint(1B) utility will generate spurious error messages when __FILE_OFFSET_BITS is set to 64. This is because the binary libc lint library, /usr/lib/llib-lc.ln, is compiled only for the standard interfaces, not with __FILE_OFFSET_BITS set to 64. This deficiency hampers static error-checking for programs compiled in the large file compilation environment.

Symbolic formats analogous to those found in <sys/int_fmtio.h> do not exist for printing or scanning variables of the types that are redefined in the large file compilation environment.

lfcompile64(5)			
NAME	lfcompile64 – transitional compilation environment		
DESCRIPTION	All 64-bit applications can manipulate large files by default. The transitional interfaces described on this page can be used by 32-bit and 64-bit applications to manipulate large files.		
	In the transitional compilation environment, explicit 64-bit functions, structures, and types are added to the API. Compiling in this environment allows both 32-bit and 64-bit applications to access files whose size is greater than or equal to 2 Gbyte (2 ³¹ bytes).		
	The transitional compilation environment exports all the explicit 64-bit functions $(xxx64())$ and types in addition to all the regular functions $(xxx())$ and types. Both $xxx()$ and $xxx64()$ functions are available to the program source. A 32-bit application must use the $xxx64()$ functions in order to access large files. See the lf64(5) manual page for a complete listing of the 64-bit transitional interfaces.		
	The transitional compilation environment differs from the large file compilation environment, wherein the underlying interfaces are bound to 64-bit functions, structures, and types. An application compiled in the large file compilation environment is able to use the $xxx()$ source interfaces to access both large and small files, rather than having to explicitly utilize the transitional $xxx64()$ interface calls to access large files. See the lfcompile(5) manual page for more information regarding the large file compilation environment.		
	Applications may combine objects produced in the large file compilation environment with objects produced in the transitional compilation environment, but must be careful with respect to interoperability between those objects. Applications should not declare global variables of types whose sizes change between compilation environments.		
	For applications that do not wish to conform to the POSIX or X/Open specifications, the 64-bit transitional interfaces are available by default. No compile-time flags need to be set.		
Access to Additional Large File Interfaces	$\frac{1}{2}$ X/Open specification-conforming interfaces should use the following compilation		
	 Set the compile-time flag _LARGEFILE64_SOURCE to 1 before including any bonders 		
	 Use the getconf(1) command with one or more of the following arguments: 		
	compilation environment		
	 Set the compile-time flag _LARGEFILE64_SOURCE to 1 before including any headers. Use the getconf(1) command with one or more of the following arguments: argument purpose LFS64_CFLAGS obtain compilation flags necessary to enable the transitional 		

lfcompile64(5)

	argument	purpose
	LFS64_LDFLAGS	obtain link editor options
	LFS64_LIBS	obtain link library names
	LFS64_LINTFLAGS	obtain lint options
EXAMPLES	In the following examples, the transitional compilation environment is accessed by invoking the getconf utility with one of the arguments listed in the table above. The additional large file interfaces are accessed either by specifying -D_LARGEFILE64_SOURCE or by invoking the getconf utility with the arguments listed above.	
	The example that uses the form of command substitution specifying the command within parentheses preceded by a dollar sign can be executed only in a POSIX-conforming shell such as the Korn Shell (see ksh(1)). In a shell that is not POSIX-conforming, such as the Bourne Shell (see sh(1)) and the C Shell (see csh(1)), the command must be enclosed within grave accent marks.	
	EXAMPLE 1 An example of compiling a program using transitional interfaces such as lseek64() and fopen64():	
	<pre>\$ c89 -D_LARGEFILE64_SOURCE \ \$ (getconf LFS64_CFLAGS) a.c \ \$ (getconf LFS64_LDFLAGS) \ \$ (getconf LFS64_LIBS)</pre>	
	EXAMPLE 2 An example of running lint on a program using transitional interfaces:	
	<pre>% lint -D_LARGEFILE64_SOURCE \ `getconf LFS64_LINTFLAGS` \ `getconf LFS64_LIBS`</pre>	
SEE ALSO	<pre>getconf(1), lseek(2), fog</pre>	pen(3C), 1f64(5), standards(5)

live	upg	rade	(5)

NAME | live_upgrade – overview of Live Upgrade feature

DESCRIPTION

The Live Upgrade feature of the Solaris operating environment enables you to maintain multiple operating system images on a single system. An image—called a boot environment, or BE—represents a set of operating system and application software packages. The BEs might contain different operating system and/or application versions.

On a system with the Solaris Live Upgrade software, your currently booted OS environment is referred to as your active, or current BE. You have one active, or current BE; all others are inactive. You can perform any number of modifications to inactive BEs on the same system, then boot from one of those BEs. If there is a failure or some undesired behavior in the newly booted BE, Live Upgrade software makes it easy for you to fall back to the previously running BE.

Live Upgrade software includes a full suite of commands, listed below and described in individual man pages, which implement all of the Live Upgrade features and functions. The software also includes a Forms and Menu Language Interpreter-based user interface named lu(1M). (See fmli(1) for a description of the Forms and Menu Language Interpreter.) The FMLI interface implements a subset of Live Upgrade functions. Unlike the command-line interfaces, output from the FMLI interface is not internationalizable.

The following are some of the tasks you can perform with Live Upgrade software:

- You can make one or more copies of the currently running system.
- You can upgrade to a new OS version on a second boot environment, then boot from that environment. If you choose, you can then fall back to your original boot environment or boot from yet another environment.
- You can install application or OS packages to a boot environment, then boot from that environment.
- You can install OS patches to a boot environment, then boot from that environment.
- From a flash archive, you can install an OS to a boot environment, then boot from that environment. See flar(1M) for information on administering flash archives.
- You can split and rejoin file systems in a new BE. For example, you can separate /usr, /var, and /opt from /, putting them on their own partitions. Conversely, you could join these file systems on a single partition under /.
- You can mount any or all of the filesystems of a BE that is not active, compare the files in any pair of BEs, delete or rename a BE, and perform other administrative tasks.

The Live Upgrade software supports upgrade from any valid Solaris installation medium, including a CD-ROM, an NFS or UFS directory, or a flash archive. (See flash_archive(4) for a description of the flash archive feature.)

In simplest terms, a BE, for Live Upgrade, consists of the disk slice containing a root file system and the file system/device (usually disk) slice entries specified in vfstab(4). This set of slices is not limited to a single disk. This means that you can have multiple BEs on a single device, or have a BE spread across slices on multiple devices.

The minimal requirement for a Live Upgrade BE is the same as for any Solaris boot environment: you must have root (/) and usr filesystems (which might both reside on /). All filesystems except for /, /usr, /var, and /opt can be shared among multiple BEs, if you choose.

Each BE must have a unique copy of the file systems that contain the OS—/, /usr, /var, and /opt. For Live Upgrade purposes, these are referred to as non-shareable (sometimes referred to as *critical*) file systems. With other file systems, such as /export or /home, you have the option of copying the files to a new BE or, the default, sharing them among BEs. These are referred to as shareable file systems. A BE is made up of a unique copy of one or more non-shareable file systems and zero or more copies of shareable file systems.

Live Upgrade commands support an option (-X) that enables XML output. Characteristics of the XML are specified in a DTD shipped with the product. XML output enables programmatic parsing of portions of the command output.

Live Upgrade supports the notion of a BE description, an optional attribute of a BE. A BE description can be of any length and format. It might be a text string or a binary file. See ludesc(1M) for details.

Below is an example set of steps that you might follow in the use of Live Upgrade software. These steps specify the use of commands rather than lu(1M), the FMLI interface. Many Live Upgrade functions are accessible through lu. Except where lu does not support a function, the choice between lu and Live Upgrade commands is a matter of your requirements and preferences. The following example is by no means exhaustive of the possibilities of the use of the Live Upgrade software.

1. You create a new BE, using lucreate(1M). The first time you create a BE on a given system, you must designate the current Solaris operating environment as a BE (give it a name). You then specify a name and a set of device (disk) slices you want to use for the new BE. The lucreate command copies the contents of the current Solaris operating environment (now a BE) to the new BE.

After you have created additional BEs, you can use a BE other than the current BE as the source for a new BE. Also, you can create an empty BE onto which you can later install a flash archive.

- 2. Using luupgrade(1M), you upgrade the OS version on your new BE (or on yet another BE you created with lucreate). The luupgrade enables you to upgrade an OS (from any valid Solaris installation medium, including a flash archive), add or remove packages (OS or application), and add or remove patches.
- 3. You use luactivate(1M) to make the new BE bootable. The next time you reboot your system, you will come up in the new BE.

live_upgrade(5)

- 4. Using lucompare(1M), you compare the system files on two different BEs. This utility gives you a comprehensive list of the files that have differences.
- 5. Using lumount(1M), you mount the filesystems of a BE that is not active, enabling you to make changes. When you are finished with the changes, use luumount(1M) to unmount the BE's file systems.
- Upon booting a new BE, you discover a failure or some other undesirable behavior. Using the procedure specified in luactivate, you can fall back to the previous BE.
- 7. Using ludelete then lucreate, you reassign file systems on the now-deleted BE to different disk slices. You separate /opt and /var from / on the new BE. Also, you specify that swap be spread over slices on multiple disks.

The following is a summary of Live Upgrade commands. All commands require root privileges.

lu

FMLI-based interface for creating and administering BEs.

```
luactivate
```

Designate a BE as the BE to boot from upon the next reboot of the system.

lucancel

Cancel a previously scheduled operation.

lucompare

Compare the contents of two BEs.

lucreate Create a BE.

lucurr

Display the name of the current BE.

ludelete

Delete a BE.

ludesc

Add or change BE descriptions.

lufslist

List the file systems on a specified BE.

lumake

Re-create a BE based on the active BE.

lumount, luumount

Mount, unmount file systems of a specified BE.

```
lurename
Rename a BE.
```

live_upgrade(5)

	ive_upgruue(0)
	lustatus For all BEs on a system, report on whether a BE is active, active upon the next reboot, in the midst of a copy operation, and whether a copy operation is scheduled for it.
	luupgrade Upgrade an OS and install application software on a BE. Such software includes flash archives, complete OS installations, OS and application packages, and OS patches.
FILES	/etc/lutab list of BEs on the system
SEE ALSO	<pre>lu(1M), luactivate(1M), lucancel(1M), lucompare(1M), lucreate(1M), lucurr(1M), ludelete(1M), ludesc(1M), lufslist(1M), lumake(1M), lumount(1M), lurename(1M), lustatus(1M), luupgrade(1M), lutab(4)</pre>
NOTES	As described in the following paragraph, Solaris Live Upgrade software is designed to install and run on multiple versions of the Solaris operating environment. Correct operation of Solaris Live Upgrade requires a certain level of patch cluster for a given OS version. Consult http://www.sunsolve.sun.com for the correct revision level for a patch cluster for your OS version.
	Live Upgrade supports the release it is distributed on and up to three marketing releases back. For example, if you obtained Live Upgrade with Solaris 9 (including a Solaris 9 upgrade), that version of Live Upgrade supports Solaris versions 2.6, Solaris 7, and Solaris 8, in addition to Solaris 9. No version of Live Upgrade supports a Solaris version prior to Solaris 2.6.

locale(5)

NAME	locale – subset of a user's environment that depends on language and cultural conventions	
DESCRIPTION	A locale is the definition of the subset of a user's environment that depends on language and cultural conventions. It is made up from one or more categories. Each category is identified by its name and controls specific aspects of the behavior of components of the system. Category names correspond to the following environment variable names:	
	LC_CTYPE	Character classification and case conversion.
	LC_COLLATE	Collation order.
	LC_TIME	Date and time formats.
	LC_NUMERIC	Numeric formatting.
	LC_MONETARY	Monetary formatting.
	LC_MESSAGES	Formats of informative and diagnostic messages and interactive responses.
	The standard utilities base their behavior on the current locale, as defined in the ENVIRONMENT section for each utility. The behavior of some of the C-language functions will also be modified based on the current locale, as defined by the last call to setlocale(3C). Locales other than those supplied by the implementation can be created by the application via the localedef(1) utility. The value that is used to specify a locale when using environment variables will be the string specified as the <i>name</i> operand to localedef when the locale was created. The strings "C" and "POSIX" are reserved as identifiers for the POSIX locale.	
	Applications can select the desired locale by invoking the setlocale() function with the appropriate value. If the function is invoked with an empty string, such as:	
	<pre>setlocale(LC_ALL, "");</pre>	
	the value of the corresponding environment variable is used. If the environment variable is unset or is set to the empty string, the setlocale() function sets the appropriate environment.	
Locale Definition	ition Locales can be described with the file format accepted by the localedef utility.	
	The locale definition file must contain one or more locale category source definitions, and must not contain more than one definition for the same locale category.	
	A category source definition consists of a category header, a category body and a category trailer. A category header consists of the character string naming of the category, beginning with the characters LC The category trailer consists of the string END, followed by one or more blank characters and the string used in the corresponding category header.	

The category body consists of one or more lines of text. Each line contains an identifier, optionally followed by one or more operands. Identifiers are either keywords, identifying a particular locale element, or collating elements. Each keyword within a locale must have a unique name (that is, two categories cannot have a commonly-named keyword); no keyword can start with the characters LC_. Identifiers must be separated from the operands by one or more blank characters.

Operands must be characters, collating elements or strings of characters. Strings must be enclosed in double-quotes. Literal double-quotes within strings must be preceded by the *<escape character>*, described below. When a keyword is followed by more than one operand, the operands must be separated by semicolons; blank characters are allowed both before and after a semicolon.

The first category header in the file can be preceded by a line modifying the comment character. It has the following format, starting in column 1:

"comment_char %c\n", <comment character>

The comment character defaults to the number sign (#). Blank lines and lines containing the *<comment character>* in the first position are ignored.

The first category header in the file can be preceded by a line modifying the escape character to be used in the file. It has the following format, starting in column 1:

"escape_char %c\n", <escape character>

The escape character defaults to backslash.

A line can be continued by placing an escape character as the last character on the line; this continuation character will be discarded from the input. Although the implementation need not accept any one portion of a continued line with a length exceeding {LINE_MAX} bytes, it places no limits on the accumulated length of the continued line. Comment lines cannot be continued on a subsequent line using an escaped newline character.

Individual characters, characters in strings, and collating elements must be represented using symbolic names, as defined below. In addition, characters can be represented using the characters themselves or as octal, hexadecimal or decimal constants. When non-symbolic notation is used, the resultant locale definitions will in many cases not be portable between systems. The left angle bracket (<) is a reserved symbol, denoting the start of a symbolic name; when used to represent itself it must be preceded by the escape character. The following rules apply to character representation:

 A character can be represented via a symbolic name, enclosed within angle brackets < and >. The symbolic name, including the angle brackets, must exactly match a symbolic name defined in the charmap file specified via the localedef -f option, and will be replaced by a character value determined from the value associated with the symbolic name in the charmap file. The use of a symbolic name not found in the charmap file constitutes an error, unless the category is LC_CTYPE

locale(5)

or LC_COLLATE, in which case it constitutes a warning condition (see localedef(1) for a description of action resulting from errors and warnings). The specification of a symbolic name in a collating-element or collating-symbol section that duplicates a symbolic name in the charmap file (if present) is an error. Use of the escape character or a right angle bracket within a symbolic name is invalid unless the character is preceded by the escape character. Example:

<c>;<c-cedilla>"<M><a><y>"

2. A character can be represented by the character itself, in which case the value of the character is implementation-dependent. Within a string, the double-quote character, the escape character and the right angle bracket character must be escaped (preceded by the escape character) to be interpreted as the character itself. Outside strings, the characters

, ; < > escape_char

must be escaped to be interpreted as the character itself.

For example:

c beta-char "May"

3. A character can be represented as an octal constant. An octal constant is specified as the escape character followed by two or more octal digits. Each constant represents a byte value. Multi-byte values can be represented by concatenated constants specified in byte order with the last constant specifying the least significant byte of the character.

For example:

\143;\347;\143\150 "\115\141\171"

4. A character can be represented as a hexadecimal constant. A hexadecimal constant is specified as the escape character followed by an x followed by two or more hexadecimal digits. Each constant represents a byte value. Multi-byte values can be represented by concatenated constants specified in byte order with the last constant specifying the least significant byte of the character.

For example:

\x63;\xe7;\x63\x68 "\x4d\x61\x79"

5. A character can be represented as a decimal constant. A decimal constant is specified as the escape character followed by a d followed by two or more decimal digits. Each constant represents a byte value. Multi-byte values can be represented by concatenated constants specified in byte order with the last constant specifying the least significant byte of the character.

For example:

\d99;\d231;\d99\d104 "\d77\d97\d121"

214 man pages section 5: Standards, Environments, and Macros • Last Revised 20 Dec 1996

Only characters existing in the character set for which the locale definition is created can be specified, whether using symbolic names, the characters themselves, or octal, decimal or hexadecimal constants. If a charmap file is present, only characters defined in the charmap can be specified using octal, decimal or hexadecimal constants. Symbolic names not present in the charmap file can be specified and will be ignored, as specified under item 1 above.

LC_CTYPE The LC_CTYPE category defines character classification, case conversion and other character attributes. In addition, a series of characters can be represented by three adjacent periods representing an ellipsis symbol (. . .). The ellipsis specification is interpreted as meaning that all values between the values preceding and following it represent valid characters. The ellipsis specification is valid only within a single encoded character set; that is, within a group of characters of the same size. An ellipsis is interpreted as including in the list all characters with an encoded value higher than the encoded value of the character preceding the ellipsis and lower than the encoded value of the character following the ellipsis.

For example:

\x30;. . .;\x39;

includes in the character class all characters with encoded values between the endpoints.

The following keywords are recognized. In the descriptions, the term "automatically included" means that it is not an error either to include or omit any of the referenced characters.

The character classes digit, xdigit, lower, upper, and space have a set of automatically included characters. These only need to be specified if the character values (that is, encoding) differ from the implementation default values.

cswidth	Moved to extensions file (see extensions(5)).
upper	Define characters to be classified as upper-case letters.
	In the POSIX locale, the 26 upper-case letters are included:
	A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
	In a locale definition file, no character specified for the keywords cntrl, digit, punct, or space can be specified. The upper-case letters A to Z are automatically included in this class.
lower	Define characters to be classified as lower-case letters. In the POSIX locale, the 26 lower-case letters are included:
	a b c d e f g h i j k l m n o p q r s t u v w x y z

locale(5)

	In a locale definition file, no character specified for the keywords cntrl, digit, punct, or space can be specified. The lower-case letters a to z of the portable character set are automatically included in this class.
alpha	Define characters to be classified as letters.
	In the POSIX locale, all characters in the classes upper and lower are included.
	In a locale definition file, no character specified for the keywords cntrl, digit, punct, or space can be specified. Characters classified as either upper or lower are automatically included in this class.
digit	Define the characters to be classified as numeric digits.
	In the POSIX locale, only
	0 1 2 3 4 5 6 7 8 9
	are included.
	In a locale definition file, only the digits 0, 1, 2, 3, 4, 5, 6, 7, 8, and 9 can be specified, and in contiguous ascending sequence by numerical value. The digits 0 to 9 of the portable character set are automatically included in this class.
	The definition of character class digit requires that only ten characters; the ones defining digits can be specified; alternative digits (for example, Hindi or Kanji) cannot be specified here.
space	Define characters to be classified as white-space characters.
	In the POSIX locale, at a minimum, the characters SPACE, FORMFEED, NEWLINE, CARRIAGE RETURN, TAB, and VERTICAL TAB are included.
	In a locale definition file, no character specified for the keywords upper, lower, alpha, digit, graph, or xdigit can be specified. The characters SPACE, FORMFEED, NEWLINE, CARRIAGE RETURN, TAB, and VERTICAL TAB of the portable character set, and any characters included in the class blank are automatically included in this class.
cntrl	Define characters to be classified as control characters.

	In the POSIX locale, no characters in classes alpha or print are included.
	In a locale definition file, no character specified for the keywords upper, lower, alpha, digit, punct, graph, print, or xdigit can be specified.
punct	Define characters to be classified as punctuation characters.
	In the POSIX locale, neither the space character nor any characters in classes alpha, digit, or cntrl are included.
	In a locale definition file, no character specified for the keywords upper, lower, alpha, digit, cntrl, xdigit or as the space character can be specified.
graph	Define characters to be classified as printable characters, not including the space character.
	In the POSIX locale, all characters in classes alpha, digit, and punct are included; no characters in class cntrl are included.
	In a locale definition file, characters specified for the keywords upper, lower, alpha, digit, xdigit, and punct are automatically included in this class. No character specified for the keyword cntrl can be specified.
print	Define characters to be classified as printable characters, including the space character.
	In the POSIX locale, all characters in class graph are included; no characters in class cntrl are included.
	In a locale definition file, characters specified for the keywords upper, lower, alpha, digit, xdigit, punct, and the space character are automatically included in this class. No character specified for the keyword cntrl can be specified.
xdigit	Define the characters to be classified as hexadecimal digits.
	In the POSIX locale, only:
	0 1 2 3 4 5 6 7 8 9 A B C D E F a b c d e f
	are included.

	In a locale definition file, only the characters defined for the class digit can be specified, in contiguous ascending sequence by numerical value, followed by one or more sets of six characters representing the hexadecimal digits 10 to 15 inclusive, with each set in ascending order (for example A, B, C, D, E, F, a, b, c, d, e, f). The digits 0 to 9, the upper-case letters A to F and the lower-case letters a to f of the portable character set are automatically included in this class.
	The definition of character class xdigit requires that the characters included in character class digit be included here also.
blank	Define characters to be classified as blank characters.
	In the POSIX locale, only the space and tab characters are included.
	In a locale definition file, the characters space and tab are automatically included in this class.
charclass	Define one or more locale-specific character class names as strings separated by semi-colons. Each named character class can then be defined subsequently in the LC_CTYPE definition. A character class name consists of at least one and at most {CHARCLASS_NAME_MAX} bytes of alphanumeric characters from the portable filename character set. The first character of a character class name cannot be a digit. The name cannot match any of the LC_CTYPE keywords defined in this document.
charclass-name	Define characters to be classified as belonging to the named locale-specific character class. In the POSIX locale, the locale-specific named character classes need not exist. If a class name is defined by a charclass keyword, but no characters are subsequently assigned to it, this is not an error; it represents a class without any characters belonging to it. The charclass-name can be used as the <i>property</i> argument to the wctype(3C) function, in regular expression and shell pattern-matching bracket expressions, and by the tr(1) command.
toupper	Define the mapping of lower-case letters to upper-case letters.
	In the POSIX locale, at a minimum, the 26 lower-case characters:

		locale(5)
		a b c d e f g h i j k l m n o p q r s t u v w x y z
		are mapped to the corresponding 26 upper-case characters:
		A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
		In a locale definition file, the operand consists of character pairs, separated by semicolons. The characters in each character pair are separated by a comma and the pair enclosed by parentheses. The first character in each pair is the lower-case letter, the second the corresponding upper-case letter. Only characters specified for the keywords lower and upper can be specified. The lower-case letters a to z, and their corresponding upper-case letters A to Z, of the portable character set are automatically included in this mapping, but only when the toupper keyword is omitted from the locale definition.
	tolower	Define the mapping of upper-case letters to lower-case letters.
		In the POSIX locale, at a minimum, the 26 upper-case characters:
		A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
		are mapped to the corresponding 26 lower-case characters:
		a b c d e f g h i j k l m n o p q r s t u v w x y z
		In a locale definition file, the operand consists of character pairs, separated by semicolons. The characters in each character pair are separated by a comma and the pair enclosed by parentheses. The first character in each pair is the upper-case letter, the second the corresponding lower-case letter. Only characters specified for the keywords lower and upper can be specified. If the tolower keyword is omitted from the locale definition, the mapping will be the reverse mapping of the one specified for toupper.
LC_COLLATE	utilities (such as sort(1), uni	ovides a collation sequence definition for numerous q(1), and so forth), regular expression matching (see 3C), strxfrm(3C), wcscoll(3C), and wcsxfrm(3C)

A collation sequence definition defines the relative order between collating elements (characters and multi-character collating elements) in the locale. This order is expressed in terms of collation values; that is, by assigning each element one or more collation values (also known as collation weights). At least the following capabilities are provided:

- 1. Multi-character collating elements. Specification of multi-character collating elements (that is, sequences of two or more characters to be collated as an entity).
- 2. User-defined ordering of collating elements. Each collating element is assigned a collation value defining its order in the character (or basic) collation sequence. This ordering is used by regular expressions and pattern matching and, unless collation weights are explicitly specified, also as the collation weight to be used in sorting.
- 3. Multiple weights and equivalence classes. Collating elements can be assigned one or more (up to the limit {COLL_WEIGHTS_MAX}) collating weights for use in sorting. The first weight is hereafter referred to as the primary weight.
- 4. One-to-Many mapping. A single character is mapped into a string of collating elements.
- 5. Equivalence class definition. Two or more collating elements have the same collation value (primary weight).
- 6. Ordering by weights. When two strings are compared to determine their relative order, the two strings are first broken up into a series of collating elements; the elements in each successive pair of elements are then compared according to the relative primary weights for the elements. If equal, and more than one weight has been assigned, then the pairs of collating elements are recompared according to the relative subsequent weights, until either a pair of collating elements compare unequal or the weights are exhausted. The following keywords are recognized in a collation sequence definition. They are described in detail in the following sections.
- 7. Define a collating-element symbol representing a multi-character collating element. This keyword is optional.
- 8. Define a collating symbol for use in collation order statements. This keyword is optional.
- 9. Define collation rules. This statement is followed by one or more collation order statements, assigning character collation values and collation weights to collating elements.
- 10. Specify the end of the collation-order statements.

collating-elementIn addition to the collating elements in the character set, the collating-elementkeywordkeyword is used to define multi-character collating elements. The syntax is:

"collating-element %s from \"%s\"\n", <collating-symbol>, <string>

220 man pages section 5: Standards, Environments, and Macros • Last Revised 20 Dec 1996

	The <i><collating-symbol></collating-symbol></i> operand is a symbolic name, enclosed between angle brackets (< and >), and must not duplicate any symbolic name in the current charmap file (if any), or any other symbolic name defined in this collation definition. The string operand is a string of two or more characters that collates as an entity. A <i><collating-element></collating-element></i> defined via this keyword is only recognized with the LC_COLLATE category. For example:
	<pre>collating-element <ch> from "<c><h>" collating-element <e-acute> from "<acute><e>" collating-element <ll> from "ll"</ll></e></acute></e-acute></h></c></ch></pre>
collating-symbol <i>keyword</i>	This keyword will be used to define symbols for use in collation sequence statements; that is, between the order_start and the order_end keywords. The syntax is:
	<pre>"collating-symbol %s\n",<collating-symbol></collating-symbol></pre>
	The <i><collating-symbol></collating-symbol></i> is a symbolic name, enclosed between angle brackets (<i><</i> and <i>></i>), and must not duplicate any symbolic name in the current charmap file (if any), or any other symbolic name defined in this collation definition.
	A collating-symbol defined via this keyword is only recognized with the LC_COLLATE category.
	For example:
	collating-symbol <upper_case> collating-symbol <high></high></upper_case>
	The collating-symbol keyword defines a symbolic name that can be associated with a relative position in the character order sequence. While such a symbolic name does not represent any collating element, it can be used as a weight.
order_start keyword	The order_start keyword must precede collation order entries and also defines the number of weights for this collation sequence definition and other collation rules.
	The syntax of the order_start keyword is:
	"order_start %s;%s;;%s\n",< <i>sort-rules></i> ,< <i>sort-rules></i>
	The operands to the order_start keyword are optional. If present, the operands define rules to be applied when strings are compared. The number of operands define how many weights each element is assigned; if no operands are present, one forward operand is assumed. If present, the first operand defines rules to be applied when comparing strings using the first (primary) weight; the second when comparing strings using the second weight, and so on. Operands are separated by semicolons (;).

Each operand consists of one or more collation directives, separated by commas (,). If the number of operands exceeds the {COLL_WEIGHTS MAX} limit, the utility will issue a warning message. The following directives will be supported: forward Specifies that comparison operations for the weight level proceed from start of string towards the end of string. backward Specifies that comparison operations for the weight level proceed from end of string towards the beginning of string. position Specifies that comparison operations for the weight level will consider the relative position of elements in the strings not subject to IGNORE. The string containing an element not subject to IGNORE after the fewest collating elements subject to IGNORE from the start of the compare will collate first. If both strings contain a character not subject to IGNORE in the same relative position, the collating values assigned to the elements will determine the ordering. In case of equality, subsequent characters not subject to IGNORE are considered in the same manner. The directives forward and backward are mutually exclusive. For example: order_start forward; backward If no operands are specified, a single forward operand is assumed. The character (and collating element) order is defined by the order in which characters and elements are specified between the order start and order end keywords. This character order is used in range expressions in regular expressions (see regex(5)). Weights assigned to the characters and elements define the collation sequence; in the absence of weights, the character order is also the collation sequence. The position keyword provides the capability to consider, in a compare, the relative position of characters not subject to IGNORE. As an example, consider the two strings "o-ring" and "or-ing". Assuming the hyphen is subject to IGNORE on the first pass, the two strings will compare equal, and the position of the hyphen is immaterial. On second pass, all characters except the hyphen are subject to IGNORE, and in the normal case the two strings would again compare equal. By taking position into account, the first collates before the second. Collation Order The order start keyword is followed by collating identifier entries. The syntax for the collating element entries is "%s %s;%s;...;%s\n"<collating-identifier>,<weight>,<weight>, ...

Each *collating-identifier* consists of either a character described in Locale Definition above, a *<collating-element>*, a *<collating-symbol>*, an ellipsis, or the special symbol UNDEFINED. The order in which collating elements are specified determines the character order sequence, such that each collating element compares less than the elements following it. The NUL character compares lower than any other character.

A *<collating-element>* is used to specify multi-character collating elements, and indicates that the character sequence specified via the *<collating-element>* is to be collated as a unit and in the relative order specified by its place.

A *<collating-symbol>* is used to define a position in the relative order for use in weights. No weights are specified with a *<collating-symbol>*.

The ellipsis symbol specifies that a sequence of characters will collate according to their encoded character values. It is interpreted as indicating that all characters with a coded character set value higher than the value of the character in the preceding line, and lower than the coded character set value for the character in the following line, in the current coded character set, will be placed in the character collation order between the previous and the following character in ascending order according to their coded character, set values. An initial ellipsis is interpreted as if the preceding line specified the NUL character, and a trailing ellipsis as if the following line specified the highest coded character set value in the current coded character set. An ellipsis is treated as invalid if the preceding or following lines do not specify characters in the current coded character set.

The symbol UNDEFINED is interpreted as including all coded character set values not specified explicitly or via the ellipsis symbol. Such characters are inserted in the character collation order at the point indicated by the symbol, and in ascending order according to their coded character set values. If no UNDEFINED symbol is specified, and the current coded character set contains characters not specified in this section, the utility will issue a warning message and place such characters at the end of the character collation order.

The optional operands for each collation-element are used to define the primary, secondary, or subsequent weights for the collating element. The first operand specifies the relative primary weight, the second the relative secondary weight, and so on. Two or more collation-elements can be assigned the same weight; they belong to the same *equivalence class* if they have the same primary weight. Collation behaves as if, for each weight level, elements subject to IGNORE are removed, unless the position collation directive is specified for the corresponding level with the order_start keyword. Then each successive pair of elements is compared according to the relative weights for the elements. If the two strings compare equal, the process is repeated for the next weight level, up to the limit {COLL WEIGHTS MAX}.

Weights are expressed as characters described in Locale Definition above, <*collating-symbol>s*, *<collating-element>s*, an ellipsis, or the special symbol IGNORE. A single character, a *<collating-symbol>* or a *<collating-element>* represent the relative

position in the character collating sequence of the character or symbol, rather than the character or characters themselves. Thus, rather than assigning absolute values to weights, a particular weight is expressed using the relative order value assigned to a collating element based on its order in the character collation sequence.

One-to-many mapping is indicated by specifying two or more concatenated characters or symbolic names. For example, if the character <eszet> is given the string "<s><s>" as a weight, comparisons are performed as if all occurrences of the character <eszet> are replaced by <s><s> (assuming that <s> has the collating weight <s>). If it is necessary to define <eszet> and <s><s> as an equivalence class, then a collating element must be defined for the string ss.

All characters specified via an ellipsis will by default be assigned unique weights, equal to the relative order of characters. Characters specified via an explicit or implicit UNDEFINED special symbol will by default be assigned the same primary weight (that is, belong to the same equivalence class). An ellipsis symbol as a weight is interpreted to mean that each character in the sequence has unique weights, equal to the relative order of their character in the character collation sequence. The use of the ellipsis as a weight is treated as an error if the collating element is neither an ellipsis nor the special symbol UNDEFINED.

The special keyword IGNORE as a weight indicates that when strings are compared using the weights at the level where IGNORE is specified, the collating element is ignored; that is, as if the string did not contain the collating element. In regular expressions and pattern matching, all characters that are subject to IGNORE in their primary weight form an equivalence class.

An empty operand is interpreted as the collating element itself.

For example, the order statement:

```
<a> <a>;<a>
```

is equal to:

<a>

An ellipsis can be used as an operand if the collating element was an ellipsis, and is interpreted as the value of each character defined by the ellipsis.

The collation order as defined in this section defines the interpretation of bracket expressions in regular expressions.

For example:

order_start

UNDEFINED

forward;backward
IGNORE;IGNORE

	<low></low>		
	<space></space>	<low>;<space></space></low>	
		<low>;</low>	
	<a>	<a>;<a>	
	<a-acute></a-acute>	<a>;<a-acute></a-acute>	
	<a-grave></a-grave>	<a>;<a-grave></a-grave>	
	<a>	<a>;<a>	
	<a-acute></a-acute>	<a>;<a-acute></a-acute>	
	<a-grave></a-grave>	<a>;<a-grave></a-grave>	
	<ch></ch>	<ch>;<ch></ch></ch>	
	<ch></ch>	<ch>;<ch></ch></ch>	
	<\$>	<\$>;<\$>	
	<eszet></eszet>	" <s><s>";"<eszet><eszet>"</eszet></eszet></s></s>	
	order_end		
	This example is interpreted as follows:		
	1. The UNDEFINED means that all characters not specified in this definition (explicitly or via the ellipsis) are ignored for collation purposes; for regular expression purposes they are ordered first.		
	2. All characters between <space> and <a ")="" for="" href="https://www.and.individual.secondary.weights-based-secondary.se</th><th>a> have the same primary equivalence class
d on their ordinal encoded values.</th></tr><tr><th></th><th>3. All characters based on the upper– or lo primary equivalence class.</th><th>ower–case character a belong to the same</th></tr><tr><th></th><th colspan=3>4. The multi-character collating element <ch> is represented by the collating symbol <ch> and belongs to the same primary equivalence class as the multi-character collating element <Ch>.</th></tr><tr><th>order_end keyword</th><th colspan=3>The collating order entries must be terminated with an order_end keyword.</th></tr><tr><th>LC_MONETARY</th><th colspan=3>The LC_MONETARY category defines the rules and symbols that are used to format monetary numeric information. This information is available through the localeconv(3C) function</th></tr><tr><th></th><th>The following items are defined in this cate
keywords recognized by the localedef(1
also similar to the member names of the lo
localeconv function returns {CHAR_MAX
empty string (" or="" size="" th="" unspecified="" ze<=""><th>) utility when defining a locale. They are conv structure defined in <locale.h>. The } for unspecified integer items and the</locale.h></th></space>) utility when defining a locale. They are conv structure defined in <locale.h>. The } for unspecified integer items and the</locale.h>	

In a locale definition file the operands are strings. For some keywords, the strings can contain only integers. Keywords that are not provided, string values set to the empty string (""), or integer keywords set to -1, are used to indicate that the value is not available in the locale.

int_curr_symbol	The international currency symbol. The operand is a four-character string, with the first three characters containing the alphabetic international currency symbol in accordance with those specified in the ISO 4217:1987 standard. The fourth character is the character used to separate the international currency symbol from the monetary quantity.
currency_symbol	The string used as the local currency symbol.
mon_decimal_point	The operand is a string containing the symbol that is used as the decimal delimiter (radix character) in monetary formatted quantities. In contexts where standards (such as the ISO C standard) limit the mon_decimal_point to a single byte, the result of specifying a multi-byte operand is unspecified.
mon_thousands_sep	The operand is a string containing the symbol that is used as a separator for groups of digits to the left of the decimal delimiter in formatted monetary quantities. In contexts where standards limit the mon_thousands_sep to a single byte, the result of specifying a multi-byte operand is unspecified.
mon_grouping	Define the size of each group of digits in formatted monetary quantities. The operand is a sequence of integers separated by semicolons. Each integer specifies the number of digits in each group, with the initial integer defining the size of the group immediately preceding the decimal delimiter, and the following integers defining the preceding groups. If the last integer is not –1, then the size of the previous group (if any) will be repeatedly used for the remainder of the digits. If the last integer is –1, then no further grouping will be performed.
	The following is an example of the interpretation of the mon_grouping keyword. Assuming that the value to be formatted is 123456789 and the mon_thousands_sep is ', then the following table shows the result. The third column shows the equivalent string in the ISO C standard that would be used by the localeconv function to accommodate this grouping.

mon_grouping	Formatted Value	ISO C String
3;-1	123456′789	"\3\177"
3	123′456′789	"\3"
3;2;-1	1234′56′789	"\3\2\177"
3;2	12'34'56'789	"\3\2"
-1	123456789	"\177"

	In these examples, the octal value of {CHAR_MAX} is 177.
positive_sign	A string used to indicate a non-negative-valued formatted monetary quantity.
negative_sign	A string used to indicate a negative-valued formatted monetary quantity.
int_frac_digits	An integer representing the number of fractional digits (those to the right of the decimal delimiter) to be written in a formatted monetary quantity using int_curr_symbol.
frac_digits	An integer representing the number of fractional digits (those to the right of the decimal delimiter) to be written in a formatted monetary quantity using currency_symbol.
p_cs_precedes	An integer set to 1 if the currency_symbol or int_curr_symbol precedes the value for a monetary quantity with a non-negative value, and set to 0 if the symbol succeeds the value.
p_sep_by_space	An integer set to 0 if no space separates the currency_symbol or int_curr_symbol from the value for a monetary quantity with a non-negative value, set to 1 if a space separates the symbol from the value, and set to 2 if a space separates the symbol and the sign string, if adjacent.
n_cs_precedes	An integer set to 1 if the currency_symbol or int_curr_symbol precedes the value for a monetary quantity with a negative value, and set to 0 if the symbol succeeds the value.
n_sep_by_space	An integer set to 0 if no space separates the currency_symbol or int_curr_symbol from the value for a monetary quantity with a negative value,

		pace separates tl if a space separa f adjacent.				
p_sign_posn	the position non-negative	An integer set to a value indicating the positioning of the positive_sign for a monetary quantity with a non-negative value. The following integer values are recognized for both p_sign_posn and n_sign_posn:				
		0 Parentheses enclose the quantity and the currency_symbol or int_curr_symbol				
	tl					
	tl					
		he sign string su urrency_symb		urr_symbol.		
n_sign_posn	the negative	An integer set to a value indicating the positioning of the negative_sign for a negative formatted monetary quantity.				
The following table show	ws the result of vari	ous combination	IS:			
		p_sep_by_	space			
		2	1	0		
p_cs_precedes=1	p_sign_posn=0	(\$1.25)	(\$1.25)	(\$1.25)		
	p_sign_posn=1	+\$1.25	+\$1.25	+\$1.25		
	p_sign_posn=2	\$1.25+	\$1.25+	\$1.25+		
	p_sign_posn=3	+\$1.25	+\$1.25	+\$1.25		
	p_sign_posn=4	\$+1.25	\$+1.25	\$+1.25		
p_cs_precedes=0	p_sign_posn=0	(1.25 \$)	(1.25 \$)	(1.25\$)		
	p_sign_posn=1	+1.25 \$	+1.25 \$	+1.25\$		
	p_sign_posn=2	1.25\$ +	1.25 \$+	1.25\$+		
	p_sign_posn=3	1.25+ \$	1.25 +\$	1.25+\$		

p_sign_posn=4 1.25\$ + 1.25\$+ 1.25\$+

The monetary formatting definitions for the POSIX locale follow; the code listing depicting the localedef(1) input, the table representing the same information with the addition of localeconv(3C) and nl_langinfo(3C) formats. All values are unspecified in the POSIX locale.

LC MONETARY

<pre># fills is the POSIX locale definition for # the LC_MONETARY category. #</pre>	L
int_curr_symbol	" "
currency_symbol	" "
mon_decimal_point	" "
mon_thousands_sep	" "
mon_grouping	-1
positive_sign	" "
negative_sign	" "
int_frac_digits	-1
p_cs_precedes	-1
p_sep_by_space	-1
n_cs_precedes	-1
n_sep_by_space	-1
p_sign_posn	-1
n_sign_posn	-1
# END LC_MONETARY	

This is the POSIX locale definition for

The entry n/a indicates that the value is not available in the POSIX locale.

LC_NUMERIC

The LC_NUMERIC category defines the rules and symbols that will be used to format non-monetary numeric information. This information is available through the localeconv(3C) function.

The following items are defined in this category of the locale. The item names are the keywords recognized by the localedef utility when defining a locale. They are also similar to the member names of the *lconv* structure defined in <locale.h>. The localeconv() function returns {CHAR_MAX} for unspecified integer items and the empty string ("") for unspecified or size zero string items.

In a locale definition file the operands are strings. For some keywords, the strings only can contain integers. Keywords that are not provided, string values set to the empty string (""), or integer keywords set to -1, will be used to indicate that the value is not available in the locale. The following keywords are recognized:

		0 5	0	
decimal_point		used as the decim numeric, non-mo keyword cannot b empty string. In c decimal_point	string containing th nal delimiter (radix on netary formatted qu be omitted and canr contexts where stand to a single byte, the i-byte operand is ur	character) in nantities. This not be set to the dards limit the e result of
thousands_sep)	used as a separate decimal delimiter monetary quantit the thousands_	string containing th or for groups of dig in numeric, non-m ies. In contexts whe sep to a single byte i-byte operand is ur	its to the left of the onetary formatted re standards limit , the result of
grouping		non-monetary qu integers separated the number of dia integer defining t preceding the dec integers defining integer is not -1 , any) will be repeat digits. If the last i will be performed formatting definit code listing depict	each group of digit antities. The operan d by semicolons. Ea- gits in each group, v he size of the group simal delimiter, and the preceding group then the size of the nteger is -1, then no l. The non-monetary tions for the POSIX ting the localede same information w ues and nl_langin	Id is a sequence of ch integer specifies with the initial immediately the following os. If the last previous group (if emainder of the p further grouping y numeric locale follow; the f input, the table ith the addition of
		LC_NUMERIC # This is the PO: # the LC_NUMERIC #	SIX locale definiti category.	lon for
		<pre># decimal_point thousands_sep grouping -1 # END LC_NUMERIC</pre>	" <period>" ""</period>	
	POSIX loca	le langinfo	localeconv()	localedef
Item	Value	Constant	Value	Value

decimal_point	"."	RADIXCHAR	"."	•
thousands_sep	n/a	THOUSEP		
grouping	n/a	-	""	-1

The entry n/a indicates that the value is not available in the POSIX locale.

- LC TIME The LC TIME category defines the interpretation of the field descriptors supported by date(1) and affects the behavior of the strftime(3C), wcsftime(3C), strptime(3C), and nl langinfo(3C) functions. Because the interfaces for C-language access and locale definition differ significantly, they are described separately. For locale definition, the following mandatory keywords are recognized: abday Define the abbreviated weekday names, corresponding to the %a field descriptor (conversion specification in the strftime(), wcsftime(), and strptime() functions). The operand consists of seven semicolon-separated strings, each surrounded by double-quotes. The first string is the abbreviated name of the day corresponding to Sunday, the second the abbreviated name of the day corresponding to Monday, and so on. day Define the full weekday names, corresponding to the %A field descriptor. The operand consists of seven semicolon-separated strings, each surrounded by double-quotes. The first string is the full name of the day corresponding to Sunday, the second the full name of the day corresponding to Monday, and so on. Define the abbreviated month names, corresponding to the %b abmon field descriptor. The operand consists of twelve semicolon-separated strings, each surrounded by double-quotes. The first string is the abbreviated name of the first month of the year (January), the second the abbreviated name of the second month, and so on. Define the full month names, corresponding to the %B field mon descriptor. The operand consists of twelve semicolon-separated strings, each surrounded by double-quotes. The first string is the full name of the first month of the year (January), the second the full name of the second month, and so on. d t fmt Define the appropriate date and time representation, corresponding to the %c field descriptor. The operand consists of a string, and can contain any combination of characters and field descriptors. In addition, the string can contain the escape sequences $\, \a, \b, \f, \n, \r, \t, \v.$
 - date_fmt Define the appropriate date and time representation, corresponding to the %C field descriptor. The operand consists of a string, and can contain any combination of characters and field

	descriptors. In addition, the string can contain the escape sequences \a,\b,\f,\n,\r,\t,\v .			
d_fmt	Define the appropriate date representation, corresponding to the x field descriptor. The operand consists of a string, and can contain any combination of characters and field descriptors. In addition, the string can contain the escape sequences $\langle , a, b, f, n, r, t, v.$			
t_fmt	Define the appropriate time representation, corresponding to the X field descriptor. The operand consists of a string, and can contain any combination of characters and field descriptors. In addition, the string can contain the escape sequences $\langle \rangle$, $\langle a, \rangle b$, $\langle f, \rangle n, \langle r, \rangle t, \langle v$.			
am_pm	Define the appropriate representation of the <i>ante meridiem</i> and <i>post meridiem</i> strings, corresponding to the %p field descriptor. The operand consists of two strings, separated by a semicolon, each surrounded by double-quotes. The first string represents the <i>ante meridiem</i> designation, the last string the <i>post meridiem</i> designation.			
t_fmt_ampm	Define the appropriate time representation in the 12-hour clock format with am_pm, corresponding to the %r field descriptor. The operand consists of a string and can contain any combination of characters and field descriptors. If the string is empty, the 12-hour format is not supported in the locale.			
era	locale. The operand	are counted and displayed for each era in a d consists of semicolon-separated strings. Each scription segment with the format:		
	direction:offset:start_	_date:end_date:era_name:era_format		
		efinitions below. There can be as many era nts as are necessary to describe the different		
	The start of an era might not be the earliest point For example, the Christian era B.C. starts on the day before January 1, A.D. 1, and increases with earlier time.			
	direction	Either a + or a – character. The + character indicates that years closer to the <i>start_date</i> have lower numbers than those closer to the <i>end_date</i> . The – character indicates that years closer to the <i>start_date</i> have higher numbers than those closer to the <i>end_date</i> .		
	offset	The number of the year closest to the <i>start_date</i> in the era, corresponding to the %Eg and %Ey field descriptors.		

			IOCale(5)
		start_date	A date in the form <i>yyy/mm</i> /dd, where <i>yyy,</i> <i>mm</i> , and dd are the year, month and day numbers respectively of the start of the era. Years prior to A.D. 1 are represented as negative numbers.
		end_date	The ending date of the era, in the same format as the <i>start_date</i> , or one of the two special values –* or +*. The value –* indicates that the ending date is the beginning of time. The value +* indicates that the ending date is the end of time.
		era_name	A string representing the name of the era, corresponding to the %EC field descriptor.
		era_format	A string for formatting the year in the era, corresponding to the %EG and %EY field descriptors.
	era_d_fmt		at of the date in alternative era notation, to the %Ex field descriptor.
	era_t_fmt		e's appropriate alternative time format, to the %EX field descriptor.
	era_d_t_fmt		e's appropriate alternative date and time format, to the %Ec field descriptor.
	alt_digits	descriptor modi strings, each sur alternative sym symbol correspo symbol strings of string correspon	ve symbols for digits, corresponding to the %O field ifier. The operand consists of semicolon-separated rrounded by double-quotes. The first string is the bol corresponding with zero, the second string the onding with one, and so on. Up to 100 alternative can be specified. The %O modifier indicates that the nding to the value specified by means of the field be used instead of the value.
LC_TIME <i>C-language</i> Access			ccessed. These correspond to constants defined in nents to the nl_langinfo(3C) function.
	ABDAY_X		abbreviated weekday names (for example Sun), re <i>x</i> is a number from 1 to 7.
	DAY_x		full weekday names (for example Sunday), where a number from 1 to 7.
	ABMON_x		abbreviated month names (for example Jan), where a number from 1 to 12.
	MON_x		full month names (for example January), where <i>x</i> number from 1 to 12.
	1		

D_T_FMT	The appropriate date and time representation.		
D_FMT	The appropriate date representation.		
T_FMT	The appropriate time representation.		
AM_STR	The appropriate ante-meridiem affix.		
PM_STR	The appropriate post-meridiem affix.		
T_FMT_AMPM	The appropriate time representation in the 12-hour clock format with AM_STR and PM_STR.		
ERA	The era description segments, which describe how years are counted and displayed for each era in a locale. Each era description segment has the format:		
	direction : offset : start_dat	e : end_date : era_name : era_format	
	many era descripti	efinitions below. There will be as on segments as are necessary to ent eras. Era description segments are colons.	
	example, the Chris	might not be the earliest point For tian era B.C. starts on the day before and increases with earlier time.	
	direction	Either a + or a – character. The + character indicates that years closer to the start_date have lower numbers than those closer to the <i>end_date</i> . The – character indicates that years closer to the start_date have higher numbers than those closer to the <i>end_date</i> .	
	offset	The number of the year closest to the start_date in the era.	
	start_date	A date in the form <i>yyyy/mm/dd</i> , where <i>yyyy, mm</i> , and dd are the year, month and day numbers respectively of the start of the era. Years prior to AD 1 are represented as negative numbers.	
	end_date	The ending date of the era, in the same format as the <i>start_date</i> , or one of the two special values –* or +*. The value –* indicates that the ending date is the beginning of	

		time. The value +* indicates that the ending date is the end of time.	
	era_name	The era, corresponding to the %EC conversion specification.	
	era_format	The format of the year in the era, corresponding to the %EY and %EY conversion specifications.	
ERA_D_FMT	The era date forma	at.	
ERA_T_FMT	The locale's appropriate alternative time format, corresponding to the %EX field descriptor.		
ERA_D_T_FMT	The locale's appropriate alternative date and time format, corresponding to the %Ec field descriptor.		
ALT_DIGITS	The alternative symbols for digits, corresponding to the %O conversion specification modifier. The value consists of semicolon-separated symbols. The first is the alternative symbol corresponding to zero, the second is the symbol corresponding to one, and so on. Up to 100 alternative symbols may be specified. The following table displays the correspondence between the items described above and the conversion specifiers used by date(1) and the strftime(3C), wcsftime(3C), and strptime(3C) functions.		

localedef	langinfo	Conversion
Keyword	Constant	Specifier
abday	ABDAY_X	%a
day	DAY_X	%A
abmon	ABMON_x	%b
mon	MON	%B
d_t_fmt	D_T_FMT	%C
date_fmt	DATE_FMT	%C
d_fmt	D_FMT	%x
t_fmt	T_FMT	%X
am_pm	AM_STR	%p
am_pm	PM_STR	%p
t_fmt_ampm	T_FMT_AMPM	%r

	localedef	langinfo	Conversion		
	Keyword	Constant	Specifier		
	era	ERA	%EC, %Eg,		
			%EG, %EY, %EY		
	era_d_fmt	ERA_D_FMT	%Ex		
	era_t_fmt	ERA_T_FMT	%EX		
	era_d_t_fmt	ERA_D_T_FMT	%EC		
	alt_digits	ALT_DIGITS	%O		
LC_TIME General Information	Although certain of the field descriptors in the POSIX locale (such as the name of the month) are shown with initial capital letters, this need not be the case in other locales. Programs using these fields may need to adjust the capitalization if the output is going to be used at the beginning of a sentence. The LC_TIME descriptions of abday, day, mon, and abmon imply a Gregorian style calendar (7-day weeks, 12-month years, leap years, and so forth). Formatting time strings for other types of calendars is outside the scope of this document set. As specified under date in Locale Definition and strftime(3C), the field descriptors corresponding to the optional keywords consist of a modifier followed by a traditional field descriptor (for instance %Ex). If the optional keywords are not supported by the implementation or are unspecified for the current locale, these field descriptors are treated as the traditional field descriptor. For instance, assume the following keywords:				
	alt_digits "Oth" ; "1st	" ; "2nd" ; "3rd" ; "4th" ;	"5th" ; \		
	"6th" ; "7th" ; "8th"	; "9th" ; "10th"			
	d_fmt "The %Od day of %B in %Y"				
	On 7/4/1776, the %x field descriptor would result in "The 4th day of July in 1776" while 7/14/1789 would come out as "The 14 day of July in 1789" It can be noted that the above example is for illustrative purposes only; the %0 modifier is primarily intended to provide for Kanji or Hindi digits in date formats.				
LC_MESSAGES	The LC_MESSAGES category defines the format and values for affirmative and negative responses.				
		e recognized as part of the loca accepts upper-case versions of			
	regex(rand consists of an extended 1 5)) that describes the acceptab n expecting an affirmative or n	le affirmative response to a		

noexpr	The operand consists of an extended regular expression that describes the acceptable negative response to a question expecting an affirmative or negative response.					
yesstr	that can l	The operand consists of a fixed string (not a regular expression) that can be used by an application for composition of a message that lists an acceptable affirmative response, such as in a prompt.				
nostr	The operand consists of a fixed string that can be used by an application for composition of a message that lists an acceptable negative response. The format and values for affirmative and negative responses of the POSIX locale follow; the code listing depicting the localedef input, the table representing the same information with the addition of nl langinfo() constants.					
	LC_MESSAGES # This is the POSIX locale definition for # the LC_MESSAGES category. # yesexpr " <circumflex><left-square-bracket><y><y><right-square-bracket></right-square-bracket></y></y></left-square-bracket></circumflex>					
	# noexpr " <circumflex><left-square-bracket><n><n><right-square-brac #</right-square-brac </n></n></left-square-bracket></circumflex>					
	#		" yesstr "yes"			
	yesstr	-				
	yesstr nostr	"no"				
	yesstr	"no"				
localedef Key	yesstr nostr END LC_MM	"no"	POSIX Locale Value			
localedef Key yesexpr	yesstr nostr END LC_MM	"no" ESSAGES	POSIX Locale Value "^[yY]"			
-	yesstr nostr END LC_MM	"no" ESSAGES langinfo Constant				
yesexpr	yesstr nostr END LC_MM	"no" ESSAGES langinfo Constant YESEXPR	"^[_Y Y] "			

SEE ALSO

date(1), locale(1), localedef(1), sort(1), tr(1), uniq(1), localeconv(3C), nl_langinfo(3C), setlocale(3C), strcoll(3C), strftime(3C), strptime(3C), strxfrm(3C), wcscoll(3C), wcsftime(3C), wcsxfrm(3C), wctype(3C), attributes(5), charmap(5), extensions(5), regex(5)

man(5)

NAME	man – macros to format Reference Manual pages					
SYNOPSIS	nroff -man filename					
	troff -man f	ilename				
DESCRIPTION	These macros are used to lay out the reference pages in this manual. Note: if <i>filename</i> contains format input for a preprocessor, the commands shown above must be piped through the appropriate preprocessor. This is handled automatically by the man(1) command. See the "Conventions" section.					
	Any text argument <i>t</i> may be zero to six words. Quotes may be used to include SPACE characters in a "word". If <i>text</i> is empty, the special treatment is applied to the next input line with text to be printed. In this way . I may be used to italicize a whole line, or . SB may be used to make small bold letters.					
	and is reset to c	A prevailing indent distance is remembered between successive indented paragraphs, and is reset to default value upon reaching a non-indented paragraph. Default units for indents <i>i</i> are ens.				
	Type font and size are reset to default values before each paragraph, and after processing font and size setting macros.					
	These strings are predefined by -man:					
	$\setminus R$ '®',	*R '®', '(Reg)' in nroff.				
	*S Change to default type size.					
Requests	* n.t.l. = next text line; p.i. = prevailing indent					
	Request	Cause	If no	Explanation		
		Break	Argument			
	.в <i>t</i>	no	<i>t</i> =n.t.l.*	Text is in bold font.		
	.BI t	no	<i>t</i> =n.t.l.	Join words, alternating bold and italic.		
	.BR t	no	<i>t</i> =n.t.l.	Join words, alternating bold and roman.		
	.DT	no	.5i 1i	Restore default tabs.		
	.HP i	yes	<i>i</i> =p.i.*	Begin paragraph with hanging indent. Set prevailing indent to <i>i</i> .		
	.I <i>t</i>	no	t=n.t.l.	Text is italic.		
	.IB t	no	t=n.t.l.	Join words, alternating italic and bold.		
	.IP x i	yes	<i>x</i> =""	Same as $. TP$ with tag x .		
	.IR t	no	t=n.t.l.	Join words, alternating italic and roman.		

Request Cause If no Explanation Break Argument .IX t no Index macro, for SunSoft internal use. Begin left-aligned paragraph. Set .LP yes prevailing indent to .5i. .P Same as .LP. yes .PD d d=.4vSet vertical distance between no paragraphs. Same as .LP. .PP yes End of relative indent. Restores .RE yes prevailing indent. .RB t no *t*=n.t.l. Join words, alternating roman and bold. .RI t t=n.t.l. Join words, alternating roman and italic. no .RS i *i*=p.i. Start relative indent, increase indent by *i*. yes Sets prevailing indent to .5i for nested indents. .SB t Reduce size of text by 1 point, make text no bold. .SH t Section Heading. yes t=n.t.l. Reduce size of text by 1 point. .SM t no .ss t yes t=n.t.l. Section Subheading. .TH n s d f m yes Begin reference page *n*, of of section *s*; *d* is the date of the most recent change. If present, *f* is the left page footer; *m* is the main page (center) header. Sets prevailing indent and tabs to .5i. .TP i *i*=p.i. Begin indented paragraph, with the tag yes given on the next text line. Set prevailing indent to *i*. .TX t p no Resolve the title abbreviation *t*; join to punctuation mark (or text) p. Conventions When formatting a manual page, man examines the first line to determine whether it requires special processing. For example a first line consisting of: ′\" t indicates that the manual page must be run through the tbl(1) preprocessor.

man(5)

A typical manual page for a co	ommand or fu	nction is laid out as follows:	
.TH title [1-9]	The name of the command or function, which serves as the title of the manual page. This is followed by the number of the section in which it appears.		
.SH NAME	The name, or list of names, by which the command is called, followed by a dash and then a one-line summary of the action performed. All in roman font, this section contains no troff(1) commands or escapes, and no macro requests. It is used to generate the windex database, which is used by the whatis(1) command.		
.SH SYNOPSIS			
	typed on word mu italics, a v that you s items are	ax of the command and its arguments, as the command line. When in boldface, a st be typed exactly as printed. When in word can be replaced with an argument supply. References to bold or italicized not capitalized in other sections, even y begin a sentence.	
	Syntactic	symbols appear in roman face:	
	[]	An argument, when surrounded by brackets is optional.	
	I	Arguments separated by a vertical bar are exclusive. You can supply only one item from such a list.	
		Arguments followed by an ellipsis can be repeated. When an ellipsis follows a bracketed set, the expression within the brackets can be repeated.	
	directive,	d, the data declaration, or #include is shown first, followed by the function on. Otherwise, the function declaration is	
.SH DESCRIPTION	external beh files or data, standard out implementat section atten	overview of the command or function's avior. This includes how it interacts with and how it handles the standard input, tput and standard error. Internals and tion details are normally omitted. This npts to provide a succinct overview in the question, "what does it do?"	

		man(0)
		Literal text from the synopsis appears in constant width, as do literal filenames and references to items that appear elsewhere in the reference manuals. Arguments are italicized.
		If a command interprets either subcommands or an input grammar, its command interface or input grammar is normally described in a USAGE section, which follows the OPTIONS section. The DESCRIPTION section only describes the behavior of the command itself, not that of subcommands.
	.SH OPTIONS	The list of options along with a description of how each affects the command's operation.
	.SH RETURN VALUES	A list of the values the library routine will return to the calling program and the conditions that cause these values to be returned.
	.SH EXIT STATUS	A list of the values the utility will return to the calling program or shell, and the conditions that cause these values to be returned.
	.SH FILES	A list of files associated with the command or function.
	.SH SEE ALSO	A comma-separated list of related manual pages, followed by references to other published materials.
	.SH DIAGNOSTICS	A list of diagnostic messages and an explanation of each.
	.SH BUGS	A description of limitations, known defects, and possible problems associated with the command or function.
FILES	/usr/share/lib/tmac/an	
	/usr/share/man/windex	
SEE ALSO	<pre>man(1), nroff(1), troff(1), w</pre>	whatis(1)
	Dale Dougherty and Tim O'Re	eilly, Unix Text Processing

man(5)

NAME	mansun – macros to format Reference Manual pages				
SYNOPSIS	nroff -mansun filename				
	troff -mansu	ın <i>filename</i>			
DESCRIPTION	These macros are used to lay out the reference pages in this manual. Note: if <i>filename</i> contains format input for a preprocessor, the commands shown above must be piped through the appropriate preprocessor. This is handled automatically by man(1). See the "Conventions" section.				
	Any text argument <i>t</i> may be zero to six words. Quotes may be used to include SPACE characters in a "word". If <i>text</i> is empty, the special treatment is applied to the next input line with text to be printed. In this way . I may be used to italicize a whole line, or .SB may be used to make small bold letters.				
	A prevailing indent distance is remembered between successive indented paragraphs, and is reset to default value upon reaching a non-indented paragraph. Default units for indents <i>i</i> are ens.				
	Type font and size are reset to default values before each paragraph, and after processing font and size setting macros.				
	These strings a	re predefined by	-mansun:		
	*R '®', '(Reg)' in nroff.				
	*S Cha	nge to default ty	pe size.		
Requests	* n.t.l. = next text line; p.i. = prevailing indent				
	Request	Cause	If no	Explanation	
		Break	Argument		
	.в <i>t</i>	no	<i>t</i> =n.t.l.*	Text is in bold font.	
	.BI t	no	t=n.t.l.	Join words, alternating bold and italic.	
	.BR t	no	t=n.t.l.	Join words, alternating bold and Roman.	
	.DT	no	.5i 1i	Restore default tabs.	
	.HP i	yes	<i>i</i> =p.i.*	Begin paragraph with hanging indent. Set prevailing indent to <i>i</i> .	
	.I t	no	<i>t</i> =n.t.l.	Text is italic.	
	.IB t	no	t=n.t.l.	Join words, alternating italic and bold.	
	.IP x i	yes	<i>x</i> =""	Same as $. TP$ with tag x .	
	.IR t	no	t=n.t.l.	Join words, alternating italic and Roman.	

	Request	Cause	If no	Explanation
		Break	Argument	
	.IX t	no	-	Index macro, for SunSoft internal use.
	.LP	yes	-	Begin left-aligned paragraph. Set prevailing indent to .5i.
	.P	yes	-	Same as .LP.
	.PD d	no	<i>d</i> =.4v	Set vertical distance between paragraphs.
	.PP	yes	-	Same as .LP.
	.RE	yes	-	End of relative indent. Restores prevailing indent.
	.RB t	no	<i>t</i> =n.t.l.	Join words, alternating Roman and bold.
	.RI <i>t</i>	no	t=n.t.l.	Join words, alternating Roman and italic.
	.RS i	yes	<i>i</i> =p.i.	Start relative indent, increase indent by <i>i</i> . Sets prevailing indent to .5i for nested indents.
	.SB t	no	-	Reduce size of text by 1 point, make text bold.
	.SH t	yes	-	Section Heading.
	.SM t	no	<i>t</i> =n.t.l.	Reduce size of text by 1 point.
	.ss t	yes	<i>t</i> =n.t.l.	Section Subheading.
	. TH <i>n s d f m</i>	yes	-	Begin reference page n , of of section s ; d is the date of the most recent change. If present, f is the left page footer; m is the main page (center) header. Sets prevailing indent and tabs to .5i.
	.TP i	yes	<i>i</i> =p.i.	Begin indented paragraph, with the tag given on the next text line. Set prevailing indent to <i>i</i> .
	.TX t p	no	-	Resolve the title abbreviation t ; join to punctuation mark (or text) p .
Conventions				tines the first line to determine whether st line consisting of:
	'\" t			
	indicates that t	he manual page	must be run thro	bugh the tbl(1) preprocessor.

A typical manual page for a co	ommand or fu	nction is laid out as follows:
.TH title [1-8]	the title of th	the command or function, which serves as ne manual page. This is followed by the ne section in which it appears.
.SH NAME	called, follow summary of this section c escapes, and	r list of names, by which the command is ved by a dash and then a one-line the action performed. All in Roman font, contains no troff(1) commands or no macro requests. It is used to generate database, which is used by the whatis(1)
.SH SYNOPSIS		
	typed on word mus italics, a v that you s items are	ax of the command and its arguments, as the command line. When in boldface, a st be typed exactly as printed. When in vord can be replaced with an argument supply. References to bold or italicized not capitalized in other sections, even y begin a sentence.
	Syntactic	symbols appear in Roman face:
	[]	An argument, when surrounded by brackets is optional.
	I	Arguments separated by a vertical bar are exclusive. You can supply only one item from such a list.
		Arguments followed by an ellipsis can be repeated. When an ellipsis follows a bracketed set, the expression within the brackets can be repeated.
	directive,	d, the data declaration, or #include is shown first, followed by the function on. Otherwise, the function declaration is
.SH DESCRIPTION	external beh files or data, standard out implementat section atten	overview of the command or function's avior. This includes how it interacts with and how it handles the standard input, tput and standard error. Internals and tion details are normally omitted. This npts to provide a succinct overview in e question, "what does it do?"

		mansun(5)
		Literal text from the synopsis appears in constant width, as do literal filenames and references to items that appear elsewhere in the reference manuals. Arguments are italicized.
		If a command interprets either subcommands or an input grammar, its command interface or input grammar is normally described in a USAGE section, which follows the OPTIONS section. The DESCRIPTION section only describes the behavior of the command itself, not that of subcommands.
	.SH OPTIONS	The list of options along with a description of how each affects the command's operation.
	.SH FILES	A list of files associated with the command or function.
	.SH SEE ALSO	A comma-separated list of related manual pages, followed by references to other published materials.
	.SH DIAGNOSTICS	A list of diagnostic messages and an explanation of each.
	.SH BUGS	A description of limitations, known defects, and possible problems associated with the command or function.
FILES	/usr/share/lib/tmac/an	Isun
	/usr/share/man/windex	
SEE ALSO	<pre>man(1), nroff(1), troff(1), v</pre>	whatis(1)
	Dale Dougherty and Tim O'R	eilly, Unix Text Processing
	1	

NAME	me – maci	ros for forma	atting paper	rs
SYNOPSIS	nroff -n	ne [<i>options</i>]	filename	
	troff -n	ne [<i>options</i>]	filename	
DESCRIPTION	facility for	ackage of nroff and troff macro definitions provides a canned formatting 7 for technical papers in various formats. When producing 2-column output on a 1 filter the output through col(1).		
		on with this p		elow. Many nroff and troff requests are unsafe in owever, these requests may be used with impunity
	.bp	begin new	page	
	.br	break outp	ut line here	
	.sp n	insert <i>n</i> spa	acing lines	
	.ls <i>n</i>	(line spacir	ng) <i>n</i> =1 sing	gle, $n=2$ double space
	.na	no alignme	ent of right	margin
	.ce <i>n</i>	center next	<i>n</i> lines	
	.ul <i>n</i>	underline r	next <i>n</i> lines	
	.sz +n	add <i>n</i> to po	oint size	
	-	the eqn(1), cceptable as		efer(1), and tbl(1) preprocessors for equations and
REQUESTS		owing list, "i o. This list is		n" refers to the first .pp, .lp, .ip, .np, .sh, or e.
	Request	Initial	Cause	Explanation
		Value	Break	
	. (c	-	yes	Begin centered block.
	. (d	-	no	Begin delayed text.
	.(f	-	no	Begin footnote.
	.(1	-	yes	Begin list.
	. (q	-	yes	Begin major quote.
	. (x <i>x</i>	-	no	Begin indexed item in index <i>x</i> .
	.(z	-	no	Begin floating keep.
	.)c	-	yes	End centered block.

Request	Initial	Cause	Explanation
	Value	Break	
.)d	-	yes	End delayed text.
.)f	-	yes	End footnote.
.)1	-	yes	End list.
.)q	-	yes	End major quote.
.)x	-	yes	End index item.
.)z	-	yes	End floating keep.
.++ m H	-	no	Define paper section.
			m defines the part of the paper,
			and can be C (chapter), A (appendix), P (preliminary, for instance,
			abstract, table of contents, etc.),
			B (bibliography), RC (chapters
			renumbered from page one each
			chapter), or RA (appendix renumbered
			from page one).
.+c T	-	yes	Begin chapter (or appendix, etc.,
			as set by .++). <i>T</i> is
			the chapter title.
.1c	1	yes	One column format on a new page.
.2c	1	yes	Two column format.
.EN	-	yes	Space after equation produced by eqn
			or neqn.
. EQ <i>x y</i>	-	yes	Precede equation; break out and
			add space. Equation number is y.
			The optional argument x may be I
			to indent equation (default),
			L to left-adjust the equation, or
			<i>C</i> to center the equation.

Standards, Environments, and Macros 247

Request	Initial	Cause	Explanation
	Value	Break	
.GE	-	yes	End gremlin picture.
.GS	-	yes	Begin gremlin picture.
.PE	-	yes	End pic picture.
.PS	-	yes	Begin pic picture.
.TE	-	yes	End table.
.TH	-	yes	End heading section of table.
.TS X	-	yes	Begin table; if x is H table
			has repeated heading.
.acAN	-	no	Set up for ACM style output.
			A is the Author's name(s), N is the
			total number of pages. Must be giver
			before the first initialization.
.b <i>x</i>	no	no	Print <i>x</i> in boldface; if no argument
			switch to boldface.
.ba + <i>n</i>	0	yes	Augments the base indent by n .
			This indent is used to set the indent
			on regular text (like paragraphs).
.bc	no	yes	Begin new column.
.bi x	no	no	Print <i>x</i> in bold italics
			(nofill only).
.bu	-	yes	Begin bulleted paragraph.
bx x	no	no	Print <i>x</i> in a box (nofill only).
.ef 'x'y'z	,,,,,	no	Set even footer to <i>x y z</i> .
.eh' $x'y'z$,,,,,	no	Set even header to $x y z$.
.fo′x′y′z	,,,,,	no	Set footer to <i>x y z</i> .
.hx	-	no	Suppress headers and footers on
			next page.
.he' $x'y'z$,,,,,	no	Set header to $x y z$.

Request	Initial	Cause	Explanation	
	Value	Break		
.hl	-	yes	Draw a horizontal line.	
.i <i>x</i>	no	no	Italicize <i>x</i> ; if <i>x</i> missing, italic	
			text follows.	
.ip x y	no	yes	Start indented paragraph, with	
			hanging tag x . Indentation is	
			y ens (default 5).	
.lp	yes	yes	Start left-blocked paragraph.	
.lo	-	no	Read in a file of local macros	
			of the form $. *x$. Must be	
			given before initialization.	
.np	1	yes	Start numbered paragraph.	
.of 'x'y'z	,,,,,	no	Set odd footer to x y z.	
.oh 'x'y'z	,,,,,	no	Set odd header to x y z.	
.pd	-	yes	Print delayed text.	
.pp	no	yes	Begin paragraph. First line indented.	
.r	yes	no	Roman text follows.	
.re	-	no	Reset tabs to default values.	
.sc	no	no	Read in a file of special characters	
			and diacritical marks. Must be	
			given before initialization.	
. sh n x	-	yes	Section head follows, font	
			automatically bold. <i>n</i> is level	
			of section, <i>x</i> is title of section.	
.sk	no	no	Leave the next page blank.	
			Only one page is remembered ahead.	
.sm <i>x</i>	-	no	Set <i>x</i> in a smaller pointsize.	
.sz+n	10p	no	Augment the point size by n points.	
.th	no	no	Produce the paper in thesis format.	

Standards, Environments, and Macros 249

me(5)	
-------	--

	Request	Initial	Cause	Explanation
		Value	Break	
				Must be given before initialization.
	.tp	no	yes	Begin title page.
	.u <i>x</i>	-	no	Underline argument (even in troff).
				(Nofill only).
	.uh	-	yes	Like . sh but unnumbered.
	.xp x	-	no	Print index <i>x</i> .
FILES	/usr/shar	e/lib/tma	ac/e	
	/usr/shar	e/lib/tma	ac/*.me	
SEE ALSO	col(1), eqn	(1), nroff(1	l), refer(1)),tbl(1),troff(1)

				1111(0)
NAME	mm – text formatting (1	memorandum) m	acros	
SYNOPSIS	<pre>nroff -mm [options]</pre>	filename		
	troff -mm [options]	filename		
DESCRIPTION	facility for various style	es of articles, thes lineprinter, or w	es, and book hen reverse	itions provides a formatting s. When producing 2-column line motions are needed, filter the defined below.
	superset of the standard Bell Labs macros have l	d -mm macro pack been removed; fo	kages as sup r instance, it	n written at Berkeley and is a plied by Bell Labs. Some of the is assumed that the user has little was generated at Whippany Labs.
		requests below n	nay be used	unction with this package. with impunity after initialization, on:
	.bp begin new p	page		
	.br break outpu	it line		
	.sp <i>n</i> insert n space	cing lines		
	.cen center next	n lines		
	.ls <i>n</i> line spacing	: <i>n</i> =1 single, <i>n</i> =2	double spac	e
	.na no alignmer	nt of right margin		
	preprocessors for equat	tions, tables, and	of the tbl(1), eqn(1) and refer(1)
REQUESTS	Here is a table of macro	DS.		
	Macro Name	Initial Value	Break? Reset?	Explanation
	.1C	on	у,у	one column format on a new page
				· · · · · · · · · · · · · · · · · · ·

у,у

y

y

_

t=1;*i*=.Li;*s*=0

.2C[*l*]

.AL[t][i][s]

.AE

Standards, Environments, and Macros 251

two column format *l*=line length

Start automatic list type t=[1,A,a,I,i] 1=arabic numbers; A=uppercase letters a=lowercase letters; I=uppercase Roman

numerals; i=lowercase Roman numerals indentation *i*; separation *s*

end abstract

mm(5)

mm(5)

Macro Name	Initial Value	Break? Reset?	Explanation
.AS m [n]	<i>n</i> =0	у	begin abstract
.AU	-	у	author's name
. AV x	-	у	signature and date line of verifier <i>x</i>
.B <i>x</i>	-	n	embolden <i>x;</i> if no <i>x,</i> switch to boldface
.BE	-	у	end block text
.BI x y	-	n	embolden <i>x</i> and underline <i>y</i>
.BL	_	у	bullet list
.BR x y	-	n	embolden <i>x</i> and use Roman font for <i>y</i>
.BS	_	n	start block text
. CN	-	у	same as .DE (nroff)
.CS	-	у	cover sheet
.CW	_	n	same as .DS I (nroff)
.DE	-	у	end display
.DF[p][f][rp]	p=L;f=N	У	start floating display; position p=[L,C,CB] L=left; I=indent; C=center; CB=center block fill f=[N,Y]; right position rp (fill only)
.DL[<i>i</i>][<i>s</i>]	_	у	start dash list
.DS[p][f][rp]	p=L;f=N	у	begin static display (see . DF for argument descriptions)
.EC x [n]	<i>n</i> =1	у	equation title; equation <i>x</i> ; number <i>n</i>
. EF <i>x</i>	-	n	even footer appears at the bottom of even-numbered pages; <i>x</i> =" <i>l</i> ' <i>c</i> ' <i>r</i> " <i>l</i> =left; <i>c</i> =center; <i>r</i> =right
. EH <i>x</i>	_	n	even header appears at the top of even-numbered pages; <i>x</i> =" <i>l</i> ' <i>c</i> ' <i>r</i> " <i>l</i> =left; <i>c</i> =center; <i>r</i> =right
.EN	-	у	end displayed equation produced by eqn
.EQ	-	у	break out equation produced by eqn
. EX x [n]	<i>n</i> =1	у	exhibit title; exhibit <i>x</i>

Macro Name	Initial Value	Break? Reset?	Explanation
			number n
.FD[f][r]	<i>f</i> =10; <i>r</i> =1	n	<pre>set footnote style format f=[0-11]; renumber r=[0,1]</pre>
.FE	-	у	end footnote
.FG x [n]	<i>n</i> =1	у	figure title; figure <i>x</i> ; number <i>n</i>
.FS	-	n	start footnote
.н <i>l</i> [<i>t</i>]	-	у	produce numbered heading level <i>l</i> =[1-7]; title <i>t</i>
. HU t	-	у	produce unnumbered heading; titl t
. I <i>x</i>	-	n	underline <i>x</i>
.IB x y	-	n	underline <i>x</i> and embolden <i>y</i>
.IR x y	-	n	underline x and use Roman font of y
.LE[<i>s</i>]	<i>s</i> =0	у	end list; separation s
.LI[<i>m</i>][<i>p</i>]	-	у	start new list item; mark <i>m</i>
			prefix <i>p</i> (mark only)
.ML <i>m</i> [<i>i</i>][<i>s</i>]	<i>s</i> =0	у	start marked list; mark <i>m</i> indentation <i>i</i> ; separation <i>s</i> =[0,1]
. MT <i>x</i>		у	memo title; title <i>x</i>
.ND X		n	no date in page footer; <i>x</i> is date on cover
.NE	-	у	end block text
.NS	-	у	start block text
. OF <i>x</i>	-	n	odd footer appears at the bottom o odd-numbered pages; x="l'c'r" l=left; c=center; r=right
.OF x	-	n	odd header appears at the top of odd-numbered pages; <i>x</i> =" <i>l</i> ' <i>c</i> ' <i>r</i> " <i>l</i> =left; <i>c</i> =center; <i>r</i> =right
.OP	-	у	skip to the top of an odd-number page
.P[t]	<i>t</i> =0	у,у	begin paragraph; <i>t</i> =[0,1] 0=justified; 1=indented

Macro Name	Initial Value	Break? Reset?	Explanation
. PF x	-	n	page footer appears at the bottom of every page; x="l'c'r" l=left; c=center; r=right
. PH <i>x</i>	-	n	page header appears at the top of every page; x="l' c' r" l=left; c=center; r=right
.R	on	n	return to Roman font
.RB <i>x y</i>	_	n	use Roman on <i>x</i> and embolden <i>y</i>
.RI x y	_	n	use Roman on <i>x</i> and underline <i>y</i>
.RP x	-	у,у	released paper format ? <i>x</i> =no stops title on first
.RS	5n	у,у	right shift: start level of relative indentation
. S m n	-	n	set character point size & vertical space character point size <i>m</i> ; vertical space <i>n</i>
. SA <i>x</i>	x=1	n	justification; <i>x</i> =[0,1]
. SK <i>x</i>	_	у	skip x pages
.SM	-	n	smaller; decrease point size by 2
.SP[x]	-	у	leave <i>x</i> blank lines
. TB x [n]	<i>n</i> =1	у	table title; table <i>x</i> ; number <i>n</i>
.TC	-	у	print table of contents (put at end of input file)
.TE	-	у	end of table processed by tbl
.TH	_	у	end multi-page header of table
.TL	-	n	title in boldface and two points larger
.TM	_	n	UC Berkeley thesis mode
. TP i	у	у	<i>i</i> =p.i. Begin indented paragraph, with the tag given on the next text line. Set prevailing indent to <i>i</i> .
. TS <i>x</i>	-	у,у	begin table; if <i>x</i> =H table has multi-page header
.TY	_	у	display centered title CONTENTS

254 man pages section 5: Standards, Environments, and Macros • Last Revised 1 Jan 1997

Macro Name	Initial Value	Break? Reset?	Explanation
.VL <i>i</i> [<i>m</i>][<i>s</i>]	<i>m</i> =0; <i>s</i> =0	у	start variable-item list; indentation <i>i</i> mark-indentation <i>m</i> ; separation <i>s</i>

REGISTERS

Formatting distances can be controlled in -mm by means of built-in number registers. For example, this sets the line length to 6.5 inches:

```
.nr LL 6.5i
```

Here is a table of number registers and their default values:

Name	Register Controls	Takes Effect	Default	
Cl	contents level	table of contents	2	
De	display eject	display	0	
Df	display floating	display	5	
Ds	display spacing	display	1v	
Hb	heading break	heading	2	
Нс	heading centering	heading	0	
Ні	heading indent	heading	1	
Ні	heading spacing	heading	1	
Hu	heading unnumbered	heading	2	
Li	list indentation	list	6(nroff) 5(troff)	
Ls	list spacing	list	6	
Pi	paragraph indent	paragraph	5	
Pt	paragraph type	paragraph	1	
Si	static indent	display	5(nroff) 3(troff)	

When resetting these values, make sure to specify the appropriate units. Setting the line length to 7, for example, will result in output with one character per line. Setting Pi to 0 suppresses paragraph indentation

Here is a list of string registers available in -mm; they may be used anywhere in the text:

Name	String's Function			
*Q	<pre>quote (" in nroff, `` in troff)</pre>			
/*U	<pre>unquote(" in nroff, '' in troff)</pre>			
*-	dash (in nroff, — in troff)			
* (MO	month (month of the year)			
* (DY	day (current date)			
\ * *	automatically numbered footnote			
\ * '	acute accent (before letter)			
\ * `	grave accent (before letter)			
* ^	circumflex (before letter)			
*,	cedilla (before letter)			
*:	umlaut (before letter)			
*~	tilde (before letter)			
\ (BU	bullet item			
\ (DT	date (month day, yr)			
\ (EM	em dash			
\(Lf	LIST OF FIGURES title			
\(Lt	LIST OF TABLES title			
\(Lx	LIST OF EXHIBITS title			
\(Le	LIST OF EQUATIONS title			
\(Rp	REFERENCES title			
\ (Tm	trademark character (TM)			

When using the extended accent mark definitions available with .AM, these strings should come after, rather than before, the letter to be accented.

FILES

/usr/share/lib/tmac/m

/usr/share/lib/tmac/mm.[nt] nroff and troff definitions of mm.

256 man pages section 5: Standards, Environments, and Macros • Last Revised 1 Jan 1997

ATTRIBUTES | See attributes(5) for descriptions of the following attributes:

ATTRIBUTE TYPE	ATTRIBUTE VALUE
Availability	SUNWdoc

SEE ALSO col(1), eqn(1), nroff(1), refer(1), tbl(1), troff(1), attributes(5)

BUGS Floating keeps and regular keeps are diverted to the same space, so they cannot be mixed together with predictable results.

ms(5)

(0)							
NAME	ms – text for	matting mac	ros				
SYNOPSIS	nroff -ms	[options] file	ename				
	troff -ms [options] filename						
DESCRIPTION	This package of nroff(1) and troff(1) macro definitions provides a formatting facility for various styles of articles, theses, and books. When producing 2-column output on a terminal or lineprinter, or when reverse line motions are needed, filter the output through col(1). All external -ms macros are defined below.						
	superset of t Bell Labs ma	Note: this -ms macro package is an extended version written at Berkeley and is a superset of the standard -ms macro packages as supplied by Bell Labs. Some of the Bell Labs macros have been removed; for instance, it is assumed that the user has little interest in producing headers stating that the memo was generated at Whippany Labs.					
	However, th	Many nroff and troff requests are unsafe in conjunction with this package. However, the first four requests below may be used with impunity after initialization, and the last two may be used even before initialization:					
	.bp b	. bp begin new page					
	.br break output line						
	sp n insert n spacing lines						
	. ce n center next n lines						
	.1s n line spacing: $n=1$ single, $n=2$ double space						
	.na n	o alignment o	of right marg	in			
	Font and point size changes with \f and \s are also allowed; for example, \fIword\fR will italicize <i>word</i> . Output of the tbl(1), eqn(1) and refer(1) preprocessors for equations, tables, and references is acceptable as input.						
REQUESTS	Macro Name Initial Value Break? Reset? Explanation						
	. AB x	.AB x - y begin abstract; if x=no do not label abstract					
	.AE	-	у	end abstract			
	.AI	-	у	author's institution			
		1	1	1			

258 man pages section 5: Standards, Environments, and Macros • Last Revised 25 Feb 1992

_

_

_

_

n

у

n

у

y

better accent mark definitions

embolden *x*; if no *x*, switch to boldface

begin text to be enclosed in a box

end boxed text and print it

author's name

.AM

.AU

.в х

.B1

.B2

Macro Name	Initial Value	Break? Reset?	Explanation	
.BT	date	n	bottom title, printed at foot of page	
.BX X	-	n	print word <i>x</i> in a box	
.CM	if t	n	cut mark between pages	
.CT	-	у,у	chapter title: page number moved to CF (TM only)	
.DA x	if n	n	force date <i>x</i> at bottom of page; today if no <i>x</i>	
.DE	-	у	end display (unfilled text) of any kind	
.DS x y	Ι	у	begin display with keep; x=I, L, C, B; y=indent	
.ID y	8n,.5i	у	indented display with no keep; <i>y</i> =indent	
.LD	-	у	left display with no keep	
.CD	-	у	centered display with no keep	
.BD	-	у	block display; center entire block	
.EF <i>x</i>	-	n	even page footer x (3 part as for .tl)	
. EH <i>x</i>	-	n	even page header x (3 part as for .tl)	
.EN	-	у	end displayed equation produced by eqn	
. EQ <i>x y</i>	-	у	break out equation; <i>x</i> =L,I,C; <i>y</i> =equation number	
.FE	-	n	end footnote to be placed at bottom of page	
.FP	-	n	numbered footnote paragraph; may be redefined	
.FS x	-	n	start footnote; <i>x</i> is optional footnote label	
.HD	undef	n	optional page header below header margin	
. I <i>x</i>	-	n	italicize <i>x</i> ; if no <i>x</i> , switch to italics	
.IP x y	-	у,у	indented paragraph, with hanging tag x; y=indent	
. IX <i>x y</i>	-	у	index words $x y$ and so on (up to 5 levels)	
.KE	-	n	end keep of any kind	
.KF	-	n	begin floating keep; text fills remainder of page	
.KS	-	у	begin keep; unit kept together on a single page	
.LG	-	n	larger; increase point size by 2	
.LP	-	у,у	left (block) paragraph.	
.MC x	-	у,у	multiple columns; <i>x</i> =column width	
.ND x	if t	n	no date in page footer; <i>x</i> is date on cover	

ms(5)

Macro Name	Initial Value	Break? Reset?	Explanation	
. NH <i>x y</i>	-	у,у	numbered header; <i>x</i> =level, <i>x</i> =0 resets, <i>x</i> =S sets to <i>y</i>	
.NL	10p	n	set point size back to normal	
. OF <i>x</i>	-	n	odd page footer x (3 part as for .tl)	
. OH <i>x</i>	-	n	odd page header x (3 part as for .tl)	
.P1	if TM	n	print header on first page	
.PP	-	у,у	paragraph with first line indented	
.PT	- % -	n	page title, printed at head of page	
. PX <i>x</i>	-	у	print index (table of contents); x=no suppresses title	
.QP	-	у,у	quote paragraph (indented and shorter)	
.R	on	n	return to Roman font	
.RE	5n	у,у	retreat: end level of relative indentation	
.RP x	-	n	released paper format; <i>x</i> =no stops title on first page	
.RS	5n	у,у	right shift: start level of relative indentation	
.SH	-	у,у	section header, in boldface	
.SM	-	n	smaller; decrease point size by 2	
.TA	8n,5n	n	set TAB characters to 8n 16n (nroff) or 5n 10	
. TC <i>x</i>	-	у	print table of contents at end; <i>x</i> =no suppresses title	
.TE	-	у	end of table processed by tbl	
.TH	-	у	end multi-page header of table	
.TL	-	у	title in boldface and two points larger	
.TM	off	n	UC Berkeley thesis mode	
. TS <i>x</i>	-	у,у	begin table; if <i>x</i> =H table has multi-page header	
.UL X	_	n	underline <i>x</i> , even in troff	
.UX X	-	n	UNIX; trademark message first time; <i>x</i> appended	
. XA <i>x y</i>	-	у	another index entry; <i>x</i> =page or no for none; y=indent	
.XE	-	у	end index entry (or series of . IX entries)	
.XP	_	у,у	paragraph with first line indented, others indented	

Macro Name	Initial Value	Break? Reset?	-	
.XS x y	-	у	begin index entry; <i>x</i> =page or no for none; <i>y</i> =indent	
.1C	on	у,у	one column format, on a new page	
.2C	-	у,у	begin two column format	
.] –	-	n	beginning of refer reference	
.[0	-	n	end of unclassifiable type of reference	
.[N	-	n	N= 1:journal-article, 2:book, 3:book-article, 4:report	

REGISTERS

Formatting distances can be controlled in -ms by means of built-in number registers. For example, this sets the line length to 6.5 inches:

.nr LL 6.5i

Here is a table of number registers and their default values:

Name	Register Controls	Takes Effect	Default
PS	point size	paragraph	10
VS	vertical spacing	paragraph	12
LL	line length	paragraph	6i
LT	title length	next page	same as LL
FL	footnote length	next .FS	5.5i
PD	paragraph distance	paragraph	1v (if n), .3v (if t)
DD	display distance	displays	1v (if n), .5v (if t)
PI	paragraph indent	paragraph	5n
QI	quote indent	next .QP	5n
FI	footnote indent	next .FS	2n
PO	page offset	next page	0 (if n), ≈1i (if t)
HM	header margin	next page	1i
FM	footer margin	next page	1i
FF	footnote format	next .FS	0 (1, 2, 3 available)

261

Standards, Environments, and Macros

ms(5)

When resetting these values, make sure to specify the appropriate units. Setting the line length to 7, for example, will result in output with one character per line. Setting FF to 1 suppresses footnote superscripting; setting it to 2 also suppresses indentation of the first line; and setting it to 3 produces an .IP-like footnote paragraph.

Here is a list of string registers available in -ms; they may be used anywhere in the text:

Name	String's Function
*Q	<pre>quote (" in nroff, " in troff)</pre>
/*U	unquote (" in nroff, " in troff)
*-	dash (in nroff, -in troff)
*(MO	month (month of the year)
*(DY	day (current date)
**	automatically numbered footnote
* <i>'</i>	acute accent (before letter)
/*'	grave accent (before letter)
*^	circumflex (before letter)
*,	cedilla (before letter)
*:	umlaut (before letter)
*~	tilde (before letter)

When using the extended accent mark definitions available with .AM, these strings should come after, rather than before, the letter to be accented.

FILES	/usr/share/lib/tmac/s
	/usr/share/lib/tmac/ms.???
SEE ALSO	col(1), eqn(1), nroff(1), refer(1), tbl(1), troff(1)
BUGS	Floating keeps and regular keeps are diverted to the same space, so they cannot be mixed together with predictable results.

ms(5)

NAME

DESCRIPTION

The mount_nfs(1M) and share_nfs(1M) commands each provide a way to specify the security mode to be used on an NFS file system through the sec=mode option. *mode* can be either sys, dh, krb5, krb5i, krb5p, or none. These security modes may also be added to the automount maps. Note that mount_nfs(1M) and automount(1M) do not support sec=none at this time.

nfssec - overview of NFS security modes

The sec=mode option on the share_nfs(1M) command line establishes the security mode of NFS servers. If the NFS connection uses the NFS Version 3 protocol, the NFS clients must query the server for the appropriate *mode* to use. If the NFS connection uses the NFS Version 2 protocol, then the NFS client will use the default security mode, which is currently sys. NFS clients may force the use of a specific security mode by specifying the sec=mode option on the command line. However, if the file system on the server is not shared with that security mode, the client may be denied access.

If the NFS client wants to authenticate the NFS server using a particular (stronger) security mode, the client will want to specify the security mode to be used, even if the connection uses the NFS Version 3 protocol. This guarantees that an attacker masquerading as the server does not compromise the client.

The NFS security modes are described below. Of these, the krb5, krb5i, krb5p modes use the Kerberos V5 protocol for authenticating and protecting the shared filesystems. Before these can be used, the system must be configured to be part of a Kerberos realm (see SEAM(5).

- sys Use AUTH_SYS authentication. The user's UNIX user-id and group-ids are passed in the clear on the network, unauthenticated by the NFS server. This is the simplest security method and requires no additional administration. It is the default used by Solaris NFS Version 2 clients and Solaris NFS servers.
- dh Use a Diffie-Hellman public key system (AUTH_DES, which is referred to as AUTH_DH in the forthcoming Internet RFC).
- krb5 Use Kerberos V5 protocol to authenticate users before granting access to the shared filesystem.
- krb5i Use Kerberos V5 authentication with integrity checking (checksums) to verify that the data has not been tampered with.
- krb5p User Kerberos V5 authentication, integrity checksums, and privacy protection (encryption) on the shared filesystem. This provides the most secure filesystem sharing, as all traffic is encrypted. It should be noted that performance might suffer on some systems when using krb5p, depending on the computational intensity of the encryption algorithm and the amount of data being transferred.
- none Use null authentication (AUTH_NONE). NFS clients using AUTH_NONE have no identity and are mapped to the anonymous user nobody by NFS

nfssec(5)

	servers. A client using a security mode other than the one with which a Solaris NFS server shares the file system will have its security mode mapped to AUTH_NONE. In this case, if the file system is shared with sec=none, users from the client will be mapped to the anonymous user. The NFS security mode none is supported by share_nfs(1M), but not by mount_nfs(1M) or automount(1M).	
FILES	/etc/nfssec.conf	NFS security service configuration file.
ATTRIBUTES	See attributes(5) for descriptions of the following attributes:	
	ATTRIBUTE TYPE	ATTRIBUTE VALUE
	Availability	SUNWnfscr
SEE ALSO	<pre>automount(1M), mount_nfs secure_rpc(3NSL), attrib</pre>	s(1M), share_nfs(1M), rpc_clnt_auth(3NSL), utes(5)
NOTES	/etc/nfssec.conf lists the intended to be user-configura	NFS security services. Do not edit this file. It is not ble.

			puni_dutitok_eneek(0)
NAME	pam_authtok_che	ck – authentication and p	assword management module
SYNOPSIS	pam_authtok_chec	k.so.l	
DESCRIPTION	pam_authtok_check provides functionality to the Password Management stack. The implementation of pam_sm_chauthtok(), performs a number of checks on the construction of the newly entered password. pam_sm_chauthtok() is invoked twice by the PAM framework, once with flags set to PAM_PRELIM_CHECK, and once with flags set to PAM_UPDATE_AUTHTOK. This module only performs its checks during the first invocation. This module expects the current authentication token in the PAM_OLDAUTHTOK item, the new (to be checked) password in the PAM_AUTHTOK item, and the login name in the PAM_USER item. The checks performed by this module are:		
	length	The password length sh specified in /etc/defa	nould not be less that the minimum ault/passwd.
	circular shift	The password should n	ot be a circular shift of the login name.
	complexity	The password should connumeric or special char	ontain at least two alpha characters and one acter.
	variation	The old and new passw	ords must differ by at least three positions.
	The following opt	ion may be passed to the	module:
	debug	syslog(3C) debugging	information at the LOG_DEBUG level
ERRORS	If the password in PAM_AUTHTOK passes all tests, PAM_SUCCESS is returned. If any of the tests fail, PAM_AUTHTOK_ERR is returned.		
FILES	/etc/default/passwd Contains the value for PASSLENGTH, the default minimal password length.		
ATTRIBUTES	See attributes(5) for descriptions of the following attributes:		
	ATT	RIBUTE TYPE	ATTRIBUTE VALUE
	Interface Stability		Evolving
	MT Level		MT-Safe with exceptions
SEE ALSO	passwd(1), pam(3PAM), pam_chauthtok(3PAM), syslog(3C), libpam(3LIB), pam.conf(4), attributes(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5), pam_unix_session(5)		
NOTES	The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.		

pam_authtok_check(5)

The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

pam_authtok_get(5)

		pani_autitok_get(5)		
NAME	pam_authtok_get – authentica	tion and password management module		
SYNOPSIS	pam_authtok_get.so.1			
DESCRIPTION	the PAM stack. It implements	The pam_authtok_get service module provides password prompting funtionality to the PAM stack. It implements pam_sm_authenticate() and pam_sm_chauthtok (), providing functionality to both the Authentication Stack and the Password Management Stack.		
Authentication Service	The implementation of pam_sm_authenticate(3PAM) prompts the user name if not set and then tries to get the authentication token from the pam handle. If the token is not set, it then prompts the user for a password and stores it in the PAM item PAM_AUTHTOK. This module is meant to be the first module on an authentication stack where users are to authenticate using a keyboard.			
Password Management Service	Due to the nature of the PAM Password Management stack traversal mechanism, the pam_sm_chauthtok(3PAM) function is called twice. Once with the PAM_PRELIM_CHECK flag, and one with the PAM_UPDATE_AUTHTOK flag.			
	In the first (PRELIM) invocation, the implementation of pam_sm_chauthtok(3PAM) moves the contents of the PAM_AUTHTOK (current authentication token) to PAM_OLDAUTHTOK, and subsequentially prompts the user for a new password. This new password is stored in PAM_AUTHTOK.			
	If a previous module has set PAM_AUTHTOK prior to the invocation of pam_authtok_get, this module turns into a NO-OP and immediately returns PAM_SUCCESS.			
	In the second (UPDATE) invocation, the user is prompted to Re-enter his password. The pam_sm_chauthtok implementation verifies this reentered password with the password stored in PAM_AUTHTOK. If the passwords match, the module returns PAM_SUCCESS.			
	The following option can be passed to the module:			
	debug syslog(30	C) debugging information at the LOG_DEBUG level		
ERRORS	The authentication service returns the following error codes:			
	PAM_SUCCESS	Successfully obtains authentication token		
	PAM_SYSTEM_ERR	Fails to retrieve username, username is NULL or empty		
	The password management se	rvice returns the following error codes:		
	PAM_SUCCESS	Successfully obtains authentication token		
	PAM_AUTHTOK_ERR	Authentication token manipulation error		

pam_authtok_get(5)

 $\label{eq:attributes} \textbf{ATTRIBUTES} ~|~ \textbf{See attributes}(5) ~ for descriptions of the following attributes:$

	ATTRIBUTE TYPE	ATTRIBUTE VALUE
	Interface Stability	Evolving
	MT Level	MT-Safe with exceptions
SEE ALSO	<pre>pam(3PAM), pam_authenticate(3PAM), attributes(5), pam_authtok_check(5), pam_authtok_store(5), pam_dhkeys(5), pam_unix_account(5), pam_unix_auth</pre>	,pam_authtok_get(5), ,pam_passwd_auth(5),pam_unix(5),
NOTES	The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.	
	The pam_unix(5) module might not be sup functionality is provided by pam_authtok pam_authtok_store(5), pam_dhkeys(5), pam_unix_account(5), pam_unix_auth	pported in a future release. Similar check(5), pam_authtok_get(5), , pam_passwd_auth(5),

pam_	_authtok_	$_store(5)$
------	-----------	-------------

NAME				
	pam_authtok_store – password management module			
SYNOPSIS	pam_authtok_stor	e.so.l		
DESCRIPTION	<pre>pam_authtok_store provides functionality to the PAM password management stack. It provides one function: pam_sm_chauthtok().</pre>			
		h flags set to PAM_UPDAT en for the user specified l	TE_AUTHTOK, this module updates the by PAM_USER.	
	The authentication token PAM_OLDAUTHTOK can be used to authenticate the user against repositories that need updating (NIS, LDAP). After successful updates, the new authentication token stored in PAM_AUTHTOK is the user's valid password.			
	repository is to be	This module honors the PAM_REPOSITORY item, which, if set, specifies which repository is to be updated. If PAM_REPOSITORY is unset, it follows the nsswitch.conf(4).		
	The following opti	on can be passed to the r	nodule:	
	debug	syslog(3C) debugging	information at the LOG_DEBUG level	
	server_policy	5	for the user, as specified by PAM_USER, is a he authentication token before updating.	
ERRORS	PAM_SUCCESS Successfully obtains authentication token			
	PAM_SYSTEM_ERF	0	username, service name, old password or ord, user name null or empty, or password	
ATTRIBUTES	See attributes(5) for descriptions of the following attributes:		following attributes:	
	ATTE	RIBUTE TYPE	ATTRIBUTE VALUE	
	Interface Stability			
	Interface Stability		Evolving	
	MT Level		Evolving MT-Safe with exceptions	
SEE ALSO	MT Level pam(3PAM), pam_ libpam(3LIB), pam_ pam_authtok_ge	<pre>m.conf(4), attributes et(5), pam_dhkeys(5), pa</pre>	~	
SEE ALSO NOTES	MT Level pam(3PAM), pam_ libpam(3LIB), pam_ pam_authtok_ge pam_unix_accou The interfaces in l	<pre>m.conf(4), attributes et(5), pam_dhkeys(5), pa unt(5), pam_unix_auth</pre>	MT-Safe with exceptions pam_chauthtok(3PAM), syslog(3C), s(5), pam_authtok_check(5), am_passwd_auth(5), pam_unix(5), (5), pam_unix_session(5) e only if each thread within the	
	MT Level pam(3PAM), pam_ libpam(3LIB), pam_ pam_authtok_gepam_unix_accou The interfaces in 1 multi-threaded app The pam_unix(5) functionality is propam_authtok_st	m.conf(4), attributes et(5), pam_dhkeys(5), p int(5), pam_unix_auth ibpam(3LIB) are MT-Safe plication uses its own PA module might not be sup ovided by pam_authtok core(5), pam_dhkeys(5)	MT-Safe with exceptions pam_chauthtok(3PAM), syslog(3C), s(5), pam_authtok_check(5), am_passwd_auth(5), pam_unix(5), (5), pam_unix_session(5) e only if each thread within the	
	MT Level pam(3PAM), pam_ libpam(3LIB), pam_ pam_authtok_gepam_unix_accou The interfaces in 1 multi-threaded app The pam_unix(5) functionality is propam_authtok_st	m.conf(4), attributes et(5), pam_dhkeys(5), p int(5), pam_unix_auth ibpam(3LIB) are MT-Safe plication uses its own PA module might not be sup ovided by pam_authtok core(5), pam_dhkeys(5)	MT-Safe with exceptions pam_chauthtok(3PAM), syslog(3C), s(5), pam_authtok_check(5), am_passwd_auth(5), pam_unix(5), (5), pam_unix_session(5) e only if each thread within the M handle. poported in a future release. Similar c_check(5), pam_authtok_get(5), , pam_passwd_auth(5),	

pam_dhkeys(5)

2				
NAME	pam_dhkeys – authe	pam_dhkeys – authentication Diffie-Hellman keys management module		
SYNOPSIS	pam_dhkeys.so.1			
DESCRIPTION	The pam_dhkeys.so.1 service module provides functionality to two PAM services: Secure RPC authentication and Secure RPC authentication token management.			
	Secure RPC authentication differs from regular unix authentication because NIS+ and other ONC RPCs use Secure RPC as the underlying security mechanism.			
	The following options may be passed to the module:			
	debug	slog(3C)debugging information at LOG_DEBUG level		
	nowarn	rn off warning messages		
Authentication Services	If the user has Diffie-Hellman keys, pam_sm_authenticate() establishes secret keys for the user specified by the PAM_USER (equivalent to running keylogin(1)), using the authentication token found in the PAM_AUTHTOK item. Not being able to establish the secret keys results in an authentication error if the NIS+ repository is used to authenticate the user and the NIS+ table permissions require secure RPC credentials to access the password field. If pam_sm_setcred() is called with PAM_ESTABLISH_CRED and the user's secure RPC credentials need to be established, these credentials are set. This is equivalent to running keylogin(1).			
	If the credentials could not be set and PAM_SILENT is not specified, a diagnostic message is displayed. If pam_setcred() is called with PAM_DELETE_CRED, the user's secure RPC credentials are unset. This is equivalent to running keylogout(1).			
	PAM_REINITIALIZ PAM_IGNORE.	_CRED and PAM_REFRESH_CRED are not supported and return		
Authentication Token Management	The pam_sm_chauthtok() implementation checks whether the old login password decrypts the users secret keys. If it doesn't this module prompts the user for an old Secure RPC password and stores it in a pam data item called SUNW_OLDRPCPASS. This data item can be used by the store module to effectively update the users secret keys.			
ERRORS	The authentication service returns the following error codes:			
	PAM_SUCCESS	Credentials set successfully.		
	PAM_IGNORE	Credentials not needed to access the password repository.		
	PAM_USER_UNKNOW	PAM_USER is not set, or the user is unknown.		
	PAM_AUTH_ERR	No secret keys were set. PAM_AUTHTOK is not set, no credentials are present or there is a wrong password.		
	PAM_BUF_ERR	Module ran out of memory.		
	PAM_SYSTEM_ERR	NIS+ subsystem failed .		

pam_dhkeys(5)

The authentication token management returns the following error codes:

PAM_SUCCESS	Old rpc password is set in SUNW_OLDRPCPASS
PAM_USER_UNKNOWN	User in PAM_USER is unknown.
PAM_AUTHTOK_ERR	User did not provide a password that decrypts the secret keys.
PAM_BUF_ERR	Module ran out of memory.

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

ATTRIBUTE TYPE	ATTRIBUTE VALUE
Interface Stability	Evolving
MT Level	MT-Safe with exceptions

SEE ALSO keylogin(1), keylogout(1), pam(3PAM), pam_authenticate(3PAM), pam_chauthtok(3PAM), pam_setcred(3PAM), pam_get_item(3PAM), pam_set_data(3PAM), pam_get_data(3PAM), syslog(3C), libpam(3LIB), pam.conf(4), attributes(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_passwd_auth(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5), pam_unix_session(5)

NOTES The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.

The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

pam_dial_auth(5)

NAME	pam_dial_auth – authentication management PAM module for dialups		
SYNOPSIS	/usr/lib/security/pam_dial_auth.so.1		
DESCRIPTION	The dialup PAM module, /usr/lib/security/pam_dial_auth.so.1, authenticates a user according to the /etc/dialups and /etc/d_passwd files. Only pam_sm_authenticate() is implemented within this module. pam_sm_setcred() is a null function. /usr/lib/security/pam_dial_auth.so.1 is designed to be stacked immediately below the /usr/lib/security/pam_unix.so.1 module for the login service.		
	<pre>pam_sm_authenticate() performs authentication only if both the /etc/dialups and /etc/d_passwd files exist. The user's terminal line is checked against entries in the /etc/dialups file. If there is a match, the user's shell is compared against entries in the /etc/d_passwd file. If there is a matching entry, the user is prompted for a password which is validated against the entry in the /etc/d_passwd file. If the passwords match, the user is authenticated. The following option may be passed in to this service module:</pre>		
	debug syslog(3C) debugging	g information at LOG_DEBUG level.	
ATTRIBUTES	See attributes(5) for description of the f	following attributes:	
	ATTRIBUTE TYPE	ATTRIBUTE VALUE	
	MT Level	MT-Safe with exceptions	
SEE ALSO	<pre>pam(3PAM), pam_authenticate(3PAM), d_passwd(4), dialups(4), libpam(3LIB), pam.conf(4), attributes(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5), pam_unix_session(5)</pre>		
NOTES	The interfaces in libpam(3LIB) are MT-Safe only if each thread within the		
	The interfaces in Tibpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle. The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).		

272 man pages section 5: Standards, Environments, and Macros • Last Revised 11 Dec 2001

NAME	pam_krb5 – authentication, account, session, and password management PAM modules for Kerberos V5		
SYNOPSIS	/usr/lib/security/pam_krb5.so.1		
DESCRIPTION	The Kerberos V5 service module for PAM, /usr/lib/security/pam_krb5.so.1, provides functionality for all four PAM modules: authentication, account management, session management, and password management. The pam_krb5.so.1 module is a shared object that can be dynamically loaded to provide the necessary functionality upon demand. Its path is specified in the PAM configuration file.		
Kerberos Authentication Module	a user, pam sm authenticate(), and to refresh the Kerberos credentials cache,		
		not needed.	
	debug	Provides syslog(3C) debugging information at LOG_DEBUG level.	
	nowarn	Turns off warning messages.	
	use_first_pass	Requests Kerberos V5 authentication with the user's initial password (entered when the user authenticated to the first authentication module in the stack). If Kerberos V5 authentication fails, or if no password has been entered, it quits and does not prompt the user for a password. This option should only be used if the authentication service is designated as optional in the pam.conf configuration file.	
	try_first_pass	Requests Kerberos V5 authentication with the user's initial password (entered when the user authenticated to the first authentication module in the stack). If	

pam_krb5(5)

		Kerberos V5 authentication fails, or if no password has been entered, the user is prompted for a password with the prompt "Kerberos Password:".
	use_xfn_pass	Requests Kerberos V5 authentication with a mapped password that has been stored under XFN. If Kerberos V5 authentication fails, or if no password has been entered, it quits and does not prompt the user for a password. This option should only be used if the authentication service is designated as optional in the pam.conf configuration file.
	try_xfn_Pass	Requests Kerberos V5 authentication with a mapped password that has been stored under XFN. If Kerberos V5 authentication fails, or if no password has been stored, the user is prompted for a password with the prompt "Kerberos Password:".
Kerberos V5 Account Management Module	The Kerberos account management component provides a function to perform account management, pam_sm_acct_mgmt(). This function checks to see if the pam_krb5 authentication module has noted that the user's password has not expired. The following options may be passed in to the Kerberos V5 service module:	
	debug	Provides syslog(3C) debugging information at LOG_DEBUG level
	nowarn	Turn off warning messages.
Kerberos V5 Session Management Module	The Kerberos V5 session management component provides functions to initiate pam_sm_open_session() and terminate pam_sm_close_session() Kerberos V5 sessions. For Kerberos V5, pam_sm_open_session is a null function. pam_close_session destroys a principal's credential cache as well as the kernel Kerberos credentials if the session being closed is the last open session on this server for the calling principal.	
Kerberos V5 Password Management Module	The Kerberos V5 password management component provides a function to change passwords pam_sm_chauthtok() in the Key Distribution Center (KDC) database. The following options can be passed in to the Kerberos V5 password module:	
moutie	debug	Provides syslog(3C) debugging information at LOG_DEBUG level.
	nowarn	Turns off warning messages.
	use_first_pass	Requests Kerberos V5 authentication with the user's initial password (entered when the user authenticated to the first authentication module in the stack). If Kerberos V5 authentication fails, or if no password has been entered, it quits and does not prompt the user for a password. If authentication succeeds, the user is prompted by "New KRB5 password:" for a new

		F
		password. The user is then prompted a second time for the new password for verification and the KDC database is updated with the new password if both responses match.
	try_first_pass	Requests Kerberos V5 authentication with the user's initial password (entered when the user authenticated to the first authentication module in the stack). If Kerberos V5 authentication fails, or if no password has been entered, the user is prompted for a password with the prompt "Old KRB5 Password:". If authentication succeeds, the user is prompted by "New KRB5 password:" for a new password. The user is then prompted a second time for the new password for verification and the KDC database is updated with the new password if both responses match.
	use_xfn_pass	Requests Kerberos V5 authentication with a mapped password that has been stored under XFN. If Kerberos V5 authentication fails, or if no password has been stored, it quits and does not prompt the user for a password. If authentication succeeds, the user is prompted by "New KRB5 password:" for a new password. The user is then prompted a second time for the new password for verification and the KDC database is updated with the new password if both responses match.
	try_xfn_pass	Requests Kerberos V5 authentication with a mapped password that has been stored under XFN. If Kerberos V5 authentication fails, or if no password has been stored, the user is prompted for a password with the prompt "Old KRB5 Password:". If authentication succeeds, the user is prompted by "New KRB5 password:" for a new password. The user is then prompted a second time for the new password for verification and the KDC database is updated with the new password if both responses match.
Sample pam.conf File	The following is a sample pam.conf configuration file with Kerberos V5 support. Please note that this is only intended to give the flavor of the pam.conf Kerberos V5 entries and is not complete.	
	<pre># # Authentication management # login auth required /usr/lib/security/\$ISA/pam_unix.so.1 login auth optional /usr/lib/security/\$ISA/pam_krb5.so.1 try_first_pass # # Account management</pre>	

pam_krb5(5)

#

dtlogin account required /usr/lib/security/\$ISA/pam_unix.so.1 dtlogin account optional /usr/lib/security/\$ISA/pam_krb5.so.1 # # Session management # other session required /usr/lib/security/\$ISA/pam_unix.so.1 other session optional /usr/lib/security/\$ISA/pam_krb5.so.1 # # Password management # other password required /usr/lib/security/\$ISA/pam_unix.so.1 other password optional /usr/lib/security/\$ISA/pam_unix.so.1 other password optional /usr/lib/security/\$ISA/pam_krb5.so.1 try_first_pass

The Kerberos V5 module entries typically follow the Unix module entries. Thus, the Kerberos V5 modules are "stacked" behind the Unix module. For the login service, the Kerberos V5 authentication module runs after the Unix module. Its entry is optional, so the user can still login if it fails, assuming that the previous Unix module succeeded. If the entry designates required instead of optional, the user cannot login if Kerberos V5 authentication fails. Because the try_first_pass option is designated, it tries the user's password entered for the Unix module. If Kerberos V5 authentication fails, or no password has been entered, the user is prompted for the Kerberos V5 password. For all session related services, the Kerberos V5 session module runs after the Unix module. For the dtlogin service, the Kerberos V5 account management module runs after the Unix module. For all password changing related services, the Kerberos V5 module runs after the Unix module. Because the try_first_pass option is designated, if the initial password entered for the Unix module authenticates Kerberos V5 successfully, the old Kerberos V5 password is not requested from the user; only the new Kerberos V5 password is requested.

ATTRIBUTES See attributes(5) for description of the following attributes:

ATTRIBUTE TYPE	ATTRIBUTE VALUE
MT Level	MT-Safe with exceptions
Interface Stability	Evolving

- SEE ALSO keylogin(1), ktutil(1), pam(3PAM), pam_authenticate(3PAM), syslog(3C), libpam(3LIB), pam.conf(4), attributes(5), SEAM(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5), pam_unix_session(5)
 - **NOTES** The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.

pam_krb5(5)

The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

NAME	pam_ldap – authentication, account, and password management PAM module for LDAP
SYNOPSIS	/usr/lib/security/pam_ldap.so.1
DESCRIPTION	The pam_ldap module implements pam_sm_authenticate, pam_sm_setcred, pam_sm_acct_mgmt, and pam_sm_chauthtok, the functions that provide functionality for the PAM authentication, account management and password management stacks. The pam_ldap module ties the authentication, account management and password change functionality to the functionality of the supporting LDAP server. For authentication, pam_ldap can authenticate the user directly to any LDAP directory server by using any supported authentication mechanism, such as DIGEST-MD5. However, the account management and password change components of pam_ldap will only work with the bundled Sun ONE Directory Server. The Sun ONE Directory Server user account management, that is, password and account lockout policy, must be properly configured on the server before it can be used by pam_ldap to provide the account management, password aging, and password syntax checking controls. Refer to the Sun ONE Directory Server Administrator's Guide that is cited in the NOTES section.
	<pre>pam_ldap must be used in conjunction with the modules that support the UNIX authentication, password, and account management., which are pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_passwd_auth(5), pam_unix_account(5), and pam_unix_auth(5). pam_ldap is designed to be stacked directly below these modules. If other modules are designed to be stacked in this manner, the modules can be stacked below the pam_ldap module. The EXAMPLES section shows how the UNIX modules are stacked with pam_ldap. When stacked together, the UNIX modules are used to control local accounts, such as root. pam_ldap is used for control with the network accounts, that is, LDAP users. For the stacks to work, pam_unix_auth, pam_unix_account, pam_passwd_auth, and pam_authtok_store must to configured with the binding control flag and the server_policy option. This configuration allows local account override of a network account.</pre>
LDAP Authentication Module	The LDAP authentication module verifies the identity of a user. The pam_sm_authenticate function uses the password entered by the user to attempt to authenticate to the LDAP server. If successful, the user is authenticated.
	The authentication method used is either defined in the client profile , or the authentication method is configured by using the ldapclient(1M) command. To determine the authentication method to use, this module first attempts to use the authentication method that is defined, for service pam_ldap, for example, serviceAuthenticationMethod:pam_ldap:sasl/DIGEST-MD5. If no authentication method is defined, pam_ldap uses the default authentication method. If neither are set, the authentication fails. This module skips the configured authentication method if the authentication method is set to none. The pam_sm_setcred(3PAM) function does nothing. This function always returns PAM_IGNORE.

	The following options may be passed to the LDAP service module:		
	debug	syslog(3C) debugging information at LOG_DEBUG level.	
	nowarn	Turn off warning messages.	
	These options are case sensitiv	ve, and the options must be used exactly as presented	
LDAP Account Management Module	The LDAP account management module validates the user's account. The pam_sm_acct_mgmt(3PAM) function authenticates to the LDAP server to verify that the user's password has not expired, or that the user's account has not been locked. The following options may be passed to the LDAP service module:		
	debug	syslog(3C) debugging information at LOG_DEBUG level.	
	nowarn	Turn off warning messages.	
	These options are case sensitiv	ve, and the options must be used exactly as presented	
LDAP Password Management Module	The preferred way to configure password management for LDAP is by using the pam_authtok_store(5) module and by specify ing the server_policy option. Use the pam_authtok_store function instead of pam_ldap for password change. When password management is configured this way, both the local and LDAP accounts are handled.pam_authtok_store(5) updates the passwords in all the repositories configured by nsswitch.conf(4).pam_ldap updates only the password in the LDAP password database.		
	The LDAP password management module provides the pam_sm_chauthtok() function to change passwords in the LDAP database.		
	The following options may be	passed to the LDAP service module:	
	debug	syslog(3C) debugging information at LOG_DEBUG level.	
	nowarn	Turn off warning messages.	
	These options are case sensitive , and the options must be used exactly as presented here.		
ERRORS	The authentication service ret	urns the following error codes:	
	PAM_SUCCESS	Authentication successful	
	PAM_MAXTRIES	Maximum number of authentication attempts exceeded	
	PAM_AUTH_ERR	Authentication failure	
	PAM_USER_UNKNOWN	No account present for user	

	PAM_BUF_ERR	Memory buffer error	
	PAM_SYSTEM_ERR	System error	
	The account management se	vice returns the following error codes:	
	PAM_SUCCESS	User allowed access to account	
	PAM_NEW_AUTHTOK_REQD	New authentication token required	
	PAM_ACCT_EXPIRED	User account has expired	
	PAM_PERM_DENIED	User denied access to account at this time	
	PAM_USER_UNKNOWN	No account present for user	
	PAM_BUF_ERROR	Memory buffer error	
	PAM_SYSTEM_ERR	System error	
	The password management	service returns the following values:	
	PAM_SUCCESS	Successfully updates authentication token	
	PAM_PERM_DENIED	No permission to update authentication token	
	PAM_AUTHTOK_ERR	Authentication token manipulation error	
	PAM_USER_UNKNOWN	No account present for user	
	PAM_BUF_ERR	Memory buffer error	
	PAM_SYSTEM_ERR	System error	
EXAMPLES	EXAMPLE 1 Using pam_ldap V	Vith Authentication	
	<pre>The following is a configuration for the login service when using pam_ldap. The service name login can be substituted for any other authentication service such dtlogin or su. Lines that begin with the # symbol are comments, and these line ignored. # Authentication management for login service is stacked. # If pam_unix_auth succeeds, pam_ldap is not invoked. # The control flag "binding" provides a local overriding # remote (LDAP) control. The "server_policy" option is used # to tell pam_unix_auth.so.1 to ignore the LDAP users.</pre>		
	EXAMPLE 2 Using pam_ldap V	Vith Account Management	
	The following is a configuration for account management when using pam_ld Lines that begin with the # symbol are ignored.		

```
EXAMPLE 2 Using pam ldap With Account Management
                                                       (Continued)
# Account management for all services is stacked
# If pam_unix_account succeeds, pam_ldap is not invoked.
# The control flag "binding" provides a local overriding
# remote (LDAP) control. The "server policy" option is used
# to tell pam_unix_account.so.1 to ignore the LDAP users.
otheraccountrequisitepam_roles.so.1otheraccountrequiredpam_projects.so.1otheraccountbindingpam_unix_account.so.1otheraccountrequiredpam_ldap.so.1
EXAMPLE 3 Using pam_authtok_store With Password Management For Both Local and
LDAP Accounts
The following is a configuration for password management when using
pam authtok store instead of pam ldap. This configuration works because
pam authtok store updates password in all the repositories configured by
nsswitch.conf(4). Lines that begin with the # symbol are comments , and the lines
are ignored.
# Password management (authentication)
passwd auth binding pam passwd auth.so.1 server policy
passwd auth required pam_ldap.so.1
# Password management (updates)
# This is the preferred stack, since it updates
# passwords stored both in the local /etc files and
# in the LDAP directory. The "server policy"
# option is used to tell pam_authtok_store to
# follow the LDAP server's policy when updating
# passwords stored in the LDAP directory
other password required pam_dhkeys.so.1
other password requisite pam_authtok_get.so.1
other password requisite pam_authtok_check.so.1
other password required pam_authtok_store.so.1 server_policy
EXAMPLE 4 Using pam 1dap With Password Management if There are no Local Accounts
Use the following configuration for password management when using pam ldap.
Lines that begin with the # symbol are comments, and athe comments are ignored.
# Password management (authentication)
# The control flag "binding" provides a local overriding
# remote (LDAP) control. The server_policy option is used
# to tell pam_passwd_auth.so.1 to ignore the LDAP users.
passwd auth binding pam_passwd_auth.so.1 server_policy
passwd auth required pam_ldap.so.1
```

Password management (updates)
This stack is limited to updating password stored in the

	EXAMPLE 4 Using pam_ldap With Password Management if There are no Local Accounts (<i>Continued</i>)		
	# LDAP directory. The preferred method is shown in Example 3.		
	other password required pam_ldap.so.1		
FILES	/var/ldap/ldap_client_file		
		The LDAP configuration files of the client. Do not manually modify these files, as these files may not be human readable. Use ldapclient(1M) to update these files.	
	/etc/pam.conf	PAM configuration file.	
ATTRIBUTES	See attributes(5) for descriptions of the following attributes:		
	ATTRIBUTE TYPE	ATTRIBUTE VALUE	
	MT-Level	MT-Safe with exceptions	
	Stability Level	Evolving	

- SEE ALSO ldap(1), idsconfig(1M), ldap_cachemgr(1M), ldapclient(1M), libpam(3LIB), pam(3PAM), pam_sm_authenticate(3PAM), pam_sm_chauthtok(3PAM), pam_sm_setcred(3PAM), syslog(3C), pam.conf(4), attributes(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5)
 - **NOTES** The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.

For information on how to configure the user account management, including password and account lockout policy for the bundled Sun ONE Directory Server, please browse the html file

/usr/iplanet/ds5/manual/en/slapd/ag/password.htm.

282 man pages section 5: Standards, Environments, and Macros • Last Revised 21 Jan 2003

pam_	_passwd_	$_auth(5)$
------	----------	------------

			puni_pusswu_uuu(0)
NAME	pam_passwd_auth – authentication module for password		
SYNOPSIS	pam_passwd_auth.so.1		
DESCRIPTION	pam_passwd_auth provides authentication functionality to the password service as implemented by passwd(1). It differs from the standard PAM authentication modules in its prompting behavior.		
	The name of the user whose password attributes are to be updated must be present in the PAM_USER item. This can be accomplished due to a previous call to pam_start(3PAM), or explicitly set by pam_set_item(3PAM). Based on the current user-id and the repository that is to by updated, the module determines whether a password is necessary for a successful update of the password repository, and if so, which password is required.		
	The following option can be p	passed to the r	nodule:
	debug	syslog(3C) level) debugging information at the LOG_DEBUG
	nowarn	Turn off war	rning messages
	server_policy	If the account authority for the user, as specified by PAM_USER, is a server, do not apply the Unix policy from the passwd entry in the name service switch.	
ERRORS	The following error codes are	returned:	
	PAM_BUF_ERR	Memory bu	ffer error
	PAM_IGNORE	Ignore mod	ule, not participating in result
	PAM_SUCCESS	Successfully	obtains authentication token
	PAM_SYSTEM_ERR	System erro	r
ATTRIBUTES	See attributes(5) for descriptions of the following attributes:		
	ATTRIBUTE TYPE		ATTRIBUTE VALUE
	Interface Stability		Evolving
	MT Level		MT-Safe with exceptions
SEE ALSO	<pre>passwd(1), pam(3PAM), pam_authenticate(3PAM), pam_start(3PAM), pam_set_item(3PAM), syslog(3C), libpam(3LIB), pam.conf(4), attributes(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5), pam_unix_session(5)</pre>		
NOTES	The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.		
			Standarda Environmenta and Maaraa 202

pam_passwd_auth(5)

This module relies on the value of the current real UID, this module is only safe for MT-applications that don't change UIDs during the call to pam_authenticate(3PAM).

The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

pam_projects(5)

NAME	pam_projects – account management PAM module for projects		
SYNOPSIS	/usr/lib/security/pam_projects.s	0.1	
DESCRIPTION	The projects service module for PAM, /usr/lib/security/pam_projects.so.1, provides functionality for the account management PAM module. The pam_projects.so.1 module is a shared object that can be dynamically loaded to provide the necessary functionality upon demand. Its path is specified in the PAM configuration file.		
	<pre>pam_projects.so.1 is designed to be stacked on top of the pam_unix_account.so.1 module for all services. This module is normally configured as "required", implying that any user lacking a default project will be denied login.</pre>		
Projects Account Management Module	The project account management component provides a function to perform account management, pam_sm_acct_mgmt(). This function uses the getdefaultproj() function (see getprojent(3PROJECT)) to retrieve the user's default project entry from the project(4) database. It then sets the project ID attribute of the calling process, using the settaskid(2) system call.		
	If the user does not belong to any project defined in the project(4) database, or if the settaskid() system call failed to set the project ID attribute of the calling process, the module will display an error message and will return error code PAM_PERM_DENIED.		
ATTRIBUTES	See attributes(5) for description of the following attributes:		
	ATTRIBUTE TYPE	ATTRIBUTE VALUE	
	ATTRIBUTE TYPE MT-Level	-	
SEE ALSO	MT-Level settaskid(2),getprojent(3PROJECT), pam_acct_mgmt(3PAM),pam.conf(4),pr ,pam_authtok_check(5),pam_authtok	ATTRIBUTE VALUE MT-Safe with exceptions Libpam(3LIB), pam(3PAM), roject(4), attributes(5) _get(5), pam_authtok_store(5),	
SEE ALSO	MT-Level settaskid(2),getprojent(3PROJECT), pam_acct_mgmt(3PAM),pam.conf(4),pr	ATTRIBUTE VALUE MT-Safe with exceptions Libpam(3LIB), pam(3PAM), roject(4), attributes(5) _get(5), pam_authtok_store(5), am_unix(5), pam_unix_account(5),	
SEE ALSO NOTES	MT-Level settaskid(2), getprojent(3PROJECT), 7 pam_acct_mgmt(3PAM), pam.conf(4), pr , pam_authtok_check(5), pam_authtok_ pam_dhkeys(5), pam_passwd_auth(5), pa	ATTRIBUTE VALUE MT-Safe with exceptions Libpam(3LIB), pam(3PAM), roject(4), attributes(5) _get(5), pam_authtok_store(5), am_unix(5), pam_unix_account(5), (5) e only if each thread within the	
	MT-Level settaskid(2), getprojent(3PROJECT), 7 pam_acct_mgmt(3PAM), pam.conf(4), pr , pam_authtok_check(5), pam_authtok_ pam_dhkeys(5), pam_passwd_auth(5), pa pam_unix_auth(5), pam_unix_session(The interfaces in libpam(3LIB) are MT-Safe	ATTRIBUTE VALUE MT-Safe with exceptions Libpam(3LIB), pam(3PAM), roject(4), attributes(5) _get(5), pam_authtok_store(5), am_unix(5), pam_unix_account(5), (5) e only if each thread within the M handle. ported in a future release. Similar _check(5), pam_authtok_get(5), pam_passwd_auth(5),	

pam_rhosts_auth(5)

NAME	pam_rhosts_auth - authentication management PAM module using ruserok()		
SYNOPSIS	/usr/lib/security/pam_rhosts_auth.so.1		
DESCRIPTION	The rhosts PAM module, /usr/lib/security/pam_rhosts_auth.so.1, authenticates a user via the rlogin authentication protocol. Only pam_sm_authenticate() is implemented within this module. pam_sm_authenticate() uses the ruserok(3SOCKET) library function to authenticate the rlogin or rsh user.pam_sm_setcred() is a null function.		
	<pre>/usr/lib/security/pam_rhosts_auth.so.1 is designed to be stacked on top of the /usr/lib/security/pam_unix.so.1 module for both the rlogin and rsh services. This module is normally configured as <i>sufficient</i> so that subsequent authentication is performed only on failure of pam_sm_authenticate(). The following option may be passed in to this service module:</pre>		
	debug syslog(3C) debugging	; information at LOG_DEBUG level.	
ATTRIBUTES	See attributes(5) for descriptions of the	following attributes:	
	ATTRIBUTE TYPE	ATTRIBUTE VALUE	
	MT Level	MT-Safe with exceptions	
SEE ALSO	<pre>pam(3PAM), pam_authenticate(3PAM), libpam(3LIB), pam.conf(4), attributes</pre>		
NOTES	The interfaces in libpam() are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.		

pam_roles(5)

NAME	pam_roles – Role Account Management PAM module for Solaris
SYNOPSIS	/usr/lib/security/pam_roles.so.1
DESCRIPTION	The Role Account Management module for PAM, /usr/lib/security/pam_roles.so.1, provides functionality for one PAM module: Account management. The pam_roles.so.1 is a shared object that can be dynamically loaded to provide the necessary functionality upon demand. Its path is specified in the PAM configuration file.
Role Account Management Module	The Role account management component provides a function to check for authorization to assume a role. It prevents direct logins to a role. It uses the user_attr(4) database to specify which users can assume which roles.
	The following options may be passed to the Role Authentication service module:
	debug syslog(3C) debugging information at LOG_DEBUG level.
	If PAM_USER (see pam_set_item(3PAM)) is specified as type normal in the user_attr(4) database, the module returns PAM_IGNORE.
	If PAM_RUSER (see pam_set_item(3PAM)) is not set, the uid of the process loading the module is used to determine PAM_RUSER.
	The module returns success if the user_attr(4) entry for PAM_RUSER has an entry in the roles field for PAM_USER; otherwise it returns PAM_PERM_DENIED.
	This module is generally stacked above the account management module pam_unix.so.1. The error messages indicating that roles cannot be logged into correctly are only issued if the user has entered the correct password.
	Here are some sample entries from pam.conf(4) demonstrating the use of the pam_roles.so.1 module:
	dtlogin account requisite /usr/lib/security/\$ISA/pam_roles.so.1 dtlogin account required /usr/lib/security/\$ISA/pam_unix.so.1 #
	su account requisite /usr/lib/security/\$ISA/pam_roles.so.1 su account requisite /usr/lib/security/\$ISA/pam_roles.so.1 #
	<pre>rlogin account requisite /usr/lib/security/\$ISA/pam_roles.so.1 rlogin account required /usr/lib/security/\$ISA/pam_unix.so.1 #</pre>
	The dtlogin program invokes pam_roles.so.1. PAM_RUSER is the username corresponding to the uid of the dtlogin process, which is 0. The user_attr entry for root user (uid 0) is empty, so all role logins are prevented through dtlogin. The same rule applies to login.
	The su program invokes pam_roles.so.1. PAM_RUSER is the username of the userid of the shell that invokes su. A user needs the appropriate entry in the roles list in user_attr(4) to be able to su to another user.

pam_roles(5)

	In the example above, the rlogin program invokes the pam_roles.so.1 module. The module checks for PAM_RUSER and determines whether the role being assumed, PAM_RUSER, is in the roles list of the userattr entry for PAM_RUSER. If it is in the roles list, the module returns PAM_SUCCESS; otherwise it returns PAM_PERM_DENIED.
SEE ALSO	<pre>keylogin(1), libpam(3LIB), pam(3PAM), pam_acct_mgmt(3PAM), pam_setcred(3PAM), pam_set_item(3PAM), syslog(3C), pam.conf(4), user_attr(4), attributes(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5), pam_unix_session(5)</pre>
NOTES	The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.
	The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

		Puni-punipic(c)	
NAME	pam_sample – a sample PAM	module	
SYNOPSIS	/usr/lib/security/pam_sample.so.1		
DESCRIPTION	The SAMPLE service module for PAM is divided into four components: authentication, account management, password management, and session management. The sample module is a shared object that is dynamically loaded to provide the necessary functionality.		
SAMPLE Authentication Component	The SAMPLE authentication module, typically /usr/lib/security/pam_sample.so.1, provides functions to test the PAM framework functionality using the pam_sm_authenticate(3PAM) call. The SAMPLE module implementation of the pam_sm_authenticate(3PAM) function compares the user entered password with the password set in the pam.conf(4) file, or the string "test" if a default test password has not been set. The following options may be passed in to the SAMPLE Authentication module:		
	debug	Syslog debugging information at the LOG_DEBUG level.	
	passwd=newone	Sets the password to be "newone."	
	first_pass_good	The first password is always good when used with the use_first_pass or try_first_pass option.	
	first_pass_bad	The first password is always bad when used with the use_first_pass or try_first_pass option.	
	always_fail	Always returns PAM_AUTH_ERR.	
	always_succeed	Always returns PAM_SUCCESS.	
	always_ignore	Always returns PAM_IGNORE.	
	use_first_pass	Use the user's initial password (entered when the user is authenticated to the first authentication module in the stack) to authenticate with the SAMPLE module. If the passwords do not match, or if this is the first authentication module in the stack, quit and do not prompt the user for a password. It is recommended that this option only be used if the SAMPLE authentication module is designated as <i>optional</i> in the pam. conf configuration file.	
	try_first_pass	Use the user's initial password (entered when the user is authenticated to the first authentication module in the stack) to authenticate with the SAMPLE module. If the passwords do not match, or if this is the first authentication module in the stack, prompt the user for a password. The SAMPLE module pam_sm_setcred(3PAM) function always returns PAM_SUCCESS.	

pam_	_samp	le(5)
------	-------	-------

SAMPLE Account Management Component	The SAMPLE Account Management Component, typically pam_sample.so.1, implements a simple access control scheme that limits machine access to a list of authorized users. The list of authorized users is supplied as option arguments to the entry for the SAMPLE account management PAM module in the pam.conf file. Note that the module always permits access to the root super user. The option field syntax to limit access is shown below: allow= <i>name[,name]</i> allow= <i>name [allow=name]</i> The example pam.conf show below permits only larry to login directly.rlogin is allowed only for don and larry. Once a user is logged in, the user can use su if the user are sam or eric.				
	login	account	require	pam_sample.so.1	allow=larry
	dtlogin	account	require	pam_sample.so.1	allow=larry
	rlogin	account	require	pam_sample.so.1	allow=don allow=larry
	su	account	require	pam_sample.so.1	allow=sam,eric
SAMPLE Password Management SAMPIPPPESSFBI Management Component ATTRIBUTES	The debug and nowarn options are also supported. The SAMPLE Password Management Component function (pam_sm_chauthtok(3PAM)), always returns PAM_SUCCESS. The SAMPLE Session Management Component functions (pam_sm_open_session(3PAM), pam_sm_close_session(3PAM)) always return PAM_SUCCESS. See attributes(5) for description of the following attributes:				
	ATTRIBUTE T	יסע		ATTRIBUTE VALUE	
	MT Level			MT-Safe with exceptions	
SEE ALSO NOTES	pam(3PAM), p pam_sm_clos pam_sm_setc The interfaces	se_session(3P cred(3PAM),li	AM), pam_sr .bpam(3LIB), re MT-Safe or	M), pam_sm_chauthtol m_open_session(3PAM pam.conf(4), attribut), es(5)

pam_smartcard(5)

			F	
NAME	pam_smartcard	– PAM authentication m	odule for Smart Card	
SYNOPSIS	/usr/lib/security/pam_smartcard.so			
DESCRIPTION	/usr/lib/sec user's informati pam_smartcan provide the nec	ion (such as user name as cd.so module is a share	ed.so, provides functionality to obtain a nd password) for a smart card. The d object that can be dynamically loaded to n demand. Its path is specified in the PAM	
Smart Card Authentication Module		authentication componenticate(3PAM) function	ent provides the on to verify the identity of a smart card user.	
	passes this data verification is s	back to its underlying la uccessful, the module ret	on collects as user input the PIN number. It yer, OCF, to perform PIN verification. If urns PAM_SUCCESS, and passes the username M modules stacked below.pam_smartcard.	
	The following o	ptions can be passed to	he Smart Card service module:	
	debug	sysolg(3C) debugg	ing information at LOG_DEBUG level.	
	nowarn	Turn off warning me	ssages.	
	verbose	Turn on verbose aut	nentication failure reporting to the user.	
Smart Card Module Configuration	configuration fi	le(/etc/pam.conf).Fo	rtcard) can be configured in the PAM or example, the following configuration on on ment) forces a user to use a smart card for	
	The following are typical values set by 'smartcard -c enable', if the comman applied to the default configuration.			
	dtlogin dtlogin dtlogin	auth requisite auth required auth required	pam_smartcard.so.1 pam_authtok_get.so.1 pam_dhkeys.so.1	
	dtlogin dtsession dtsession dtsession dtsession	auth required auth requisite auth required auth required auth required	<pre>pam_unix_auth.so.1 pam_smartcard.so.1 pam_authtok_get.so.1 pam_dhkeys.so.1 pam_unix_auth.so.1</pre>	
SEE ALSO	pam_start(3P	AM), pam.conf(4), pam_store(5), pam_dhkeys	PAM), pam_authenticate(3PAM), _authtok_check(5), pam_authtok_get(5), (5), pam_passwd_auth(5), pam_unix(5), th(5), pam_unix_session(5)	

pam_smartcard(5)

NOTES The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

pam_unix(5)

NAME	pam_unix – authentication, account, session, and password management PAM modules for UNIX			
SYNOPSIS	/usr/lib/security/pam_unix.so.1			
DESCRIPTION	The UNIX service module for PAM, /usr/lib/security/pam_unix.so.1, provides functionality for all four PAM modules: authentication, account management, session management and password management. The pam_unix.so.1 module is a shared object that can be dynamically loaded to provide the necessary functionality upon demand. Its path is specified in the PAM configuration file.			
UNIX Authentication Module	The UNIX authentication component provides functions to verify the identity of a user, (pam_sm_authenticate()) and to set user specific credentials (pam_sm_setcred()). pam_sm_authenticate() compares the user entered password with the password from the UNIX password database. If the passwords match, the user is authenticated. If the user also has secure RPC credentials and the secure RPC password is the same as the UNIX password, then the secure RPC credentials are also obtained.			
	The following options may be	passed to the UNIX service module:		
	debug	syslog(3C) debugging information at LOG_DEBUG level.		
	nowarn	Turn off warning messages.		
	use_first_pass	It compares the password in the password database with the user's initial password (entered when the user authenticated to the first authentication module in the stack). If the passwords do not match, or if no password has been entered, it quits and does not prompt the user for a password. This option should only be used if the authentication service is designated as <i>optional</i> in the pam.conf configuration file.		
	try_first_pass	It compares the password in the password database with the user's initial password (entered when the user authenticated to the first authentication module in the stack). If the passwords do not match, or if no password has been entered, prompt the user for a password. When prompting for the current password, the UNIX authentication module will use the prompt, "password:" unless one of the following scenarios occur: 1. The option try_first_pass is specified and the password entered for the first module in the stack fails for the UNIX module.		

pam_unix(5)

		 The option try_first_pass is not specified, and the earlier authentication modules listed in the pam.conf file have prompted the user for the password. 	
		In these two cases, the UNIX authentication module will use the prompt "SYSTEM password:". The pam_sm_setcred() function sets user specific credentials. If the user had secure RPC credentials, but the secure RPC password was not the same as the UNIX password, then a warning message is printed. If the user wants to get secure RPC credentials, then keylogin(1) needs to be run.	
UNIX Account Management Module	management, pam entry from the UN	management component provides a function to perform account _sm_acct_mgmt(). The function retrieves the user's password IX password database and verifies that the user's account and expired. The following options may be passed in to the UNIX	
	debug	syslog(3C) debugging information at LOG_DEBUG level.	
	nowarn	Turn off warning messages.	
UNIX Session Management Module	The UNIX session management component provides functions to initiate pam_sm_open_session() and terminate pam_sm_close_session() UNIX sessions. For UNIX, pam_open_session updates the /var/adm/lastlog file. The account management module reads this file to determine the previous time the user logged in. The following options may be passed in to the UNIX service module:		
	debug	$syslog(3C)$ debugging information at LOG_DEBUG level.	
	nowarn	Turn off warning messages. pam_close_session is a null function.	
UNIX Password Management Module	passwords pam_sr must be <i>required</i> in	rd management component provides a function to change chauthtok() in the UNIX password database. This module pam.conf. It cannot be <i>optional</i> or <i>sufficient</i> . The following options to the UNIX service module:	
	debug	syslog(3C) Debugging information at LOG_DEBUG level.	
	nowarn	Turn off warning messages.	
	use_first_pass	It compares the password in the password database with the user's old password (entered to the first password module in the stack). If the passwords do not match, or if no password has been entered, it quits and does not prompt the user for the old password. It also attempts to use the new password (entered to the first	

pam_unix(5)

		pani_anix(5)
	for this mod	odule in the stack) as the new password ule. If the new password fails, it quits and ompt the user for a new password.
try_first_pass	with the use password m match, or if the user for the new pas module in th module. If the for a new pas the UNIX ac authentication unique nam password m authentication	the password in the password database r's old password (entered to the first odule in the stack). If the passwords do not no password has been entered, it prompts the old password. It also attempts to use sword (entered to the first password ne stack) as the new password for this ne new password fails, it prompts the user assword. If the user's password has expired, count module saves this information in the on handle using pam_set_data(), with a e, SUNW_UNIX_AUTHOK_DATA. The UNIX odule retrieves this information from the on handle using pam_get_data() to whether or not to force the user to update assword.
See attributes(5) for description of the following attributes:		
ATTRIBUTE TYPE		ATTRIBUTE VALUE
MT Level		MT-Safe with exceptions
<pre>syslog(3C), libpam(3LIB), p pam_authtok_check(5), pa pam_dhkeys(5), pam_passw</pre>	pam.conf(4), m_authtok_ d_auth(5),pa	attributes(5), get(5),pam_authtok_store(5), am_unix_account(5),
The pam_unix(5) module mi functionality is provided by p pam_authtok_store(5), pa	ght not be sup bam_authtok m_dhkeys(5),	pported in a future release. Similar check(5), pam_authtok_get(5), , pam_passwd_auth(5),
	See attributes(5) for descr ATTRIBUTE TYPE MT Level keylogin(1), pam(3PAM), pa syslog(3C), libpam(3LIB), pam_authtok_check(5), pa pam_authtok_check(5), pam_unix_auth(5), pam_unix_auth(5), pam_unix_auth(5), pam_unix_auth(5), pam_unix_auth(5), pam_unix_auth(5), pam_unix_because the pam_unix(5) module minister functionality is provided by pam_authtok_store(5), pam_unix_because the pam_authtok_store(5), pam_authtok_store(5), pam_unix_because the pam_authtok_store(5), pam_authtok_store	for this mod does not pro- try_first_pass It compares with the use password m match, or if the user for the new pas module in th module. If th for a new pa the UNIX ac authentication unique nam password m authentication determine w the user's pa See attributes(5) for description of the f

pam_unix_account(5)

NAME	pam_unix_account	t – PAM ace	count manage	ment module for UNIX
SYNOPSIS	pam_unix_account.so.1			
DESCRIPTION	pam_unix_account provides functionality to the PAM account management stack. The function pam(3PAM) function retrieves password aging information from the repositories specified in nsswitch.conf(4), and verifies that the user's account and password have not expired.			
	The following opti	ions can be	passed to the	module:
	debug	syslog(3	C) debugging	information at the LOG_DEBUG level
	nowarn	Turn off w	varning messa	ges
	server_policy	a server, c		for the user, as specified by PAM_USER, is ne Unix policy from the passwd entry in n.
ERRORS	The following valu	ies are retu	rned:	
	PAM_AUTHTOK_EX	VPIRED	Password ex	pired and no longer usable
	PAM_BUF_ERR		Memory but	ffer error
	PAM_IGNORE		Ignore modu	ule, not participating in result
	PAM_NEW_AUTHTC	OK_REQD	Obtain new	authentication token from the user
	PAM_SERVICE_ER	RR	Error in und	erlying service module
	PAM_SUCCESS		Successfully	obtains authentication token
ATTRIBUTES	See attributes(5) for descriptions of the following attributes:			
	ATTF	RIBUTE TYPE		ATTRIBUTE VALUE
	Interface Stability			Evolving
	MT Level			MT-Safe with exceptions
SEE ALSO	nsswitch.conf(pam_authtok_ge	4),attrib et(5),pam_	utes(5), pam_ authtok_st	<pre>syslog(3C), libpam(3LIB), pam.conf(4), _authtok_check(5), ore(5), pam_dhkeys(5), _unix_auth(5), pam_unix_session(5)</pre>
NOTES	The interfaces in 1 multi-threaded ap			e only if each thread within the M handle.
	functionality is propam_authtok_st	ovided by p core(5), pa	am_authtok m_dhkeys(5),	pported in a future release. Similar check(5), pam_authtok_get(5), , pam_passwd_auth(5), (5), and pam_unix_session(5).

NAME	pam_unix_auth – PAM a	uthentication mod	ule for UNIX
SYNOPSIS	pam_unix_auth.so.1		
DESCRIPTION	The pam_unix_auth module implements pam_sm_authenticate(), which provides functionality to the PAM authentication stack. It provides functions to verify that the password contained in the PAM item PAM_AUTHTOK is the correct password for the user specified in the item PAM_USER. If PAM_REPOSITORY is specified, then user's passwd is fetched from that repository. Otherwise the default nsswitch.conf(4) repository is searched for that user.		
	The following options ca	n be passed to the	module:
	a ser		for the user, as specified by PAM_USER, is ne Unix policy from the passwd entry in n.
ERRORS	The following values are	returned:	
	PAM_AUTH_ERR	Authenticat	ion failure
	PAM_BUF_ERR	Memory but	ffer error
	PAM_IGNORE	Ignore modu	ule, not participating in result
	PAM_PERM_DENIED	Permission of	denied
	PAM_SUCCESS	Successfully	obtains authentication token
	PAM_SYSTEM_ERR	System error	r
	PAM_USER_UNKNOWN	No account	present for user
ATTRIBUTES	See attributes(5) for a	descriptions of the	following attributes:
	ATTRIBUTE	ТҮРЕ	ATTRIBUTE VALUE
	Interface Stability		Evolving
	MT Level		MT-Safe with exceptions
SEE ALSO	<pre>pam(3PAM), pam_authe nsswitch.conf(4), att pam_authtok_get(5), p pam_passwd_auth(5), p pam_unix_session(5)</pre>	cributes(5),pam_ pam_authtok_st	ore(5), pam_dhkeys(5),
NOTES	The interfaces in libpan multi-threaded application		e only if each thread within the M handle.

pam_unix_auth(5)

The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

NAME	pam_unix_session – session management PAM module for UNIX		
SYNOPSIS	pam_unix_session.so.1		
DESCRIPTION	<pre>pam_unix_session provides functions to initiate pam_sm_open_session(3PAM) and to terminate pam_sm_close_session(3PAM).</pre>		
	pam_open_session updates the /var/adm/lastlog file. The account management module reads this file to determine the previous time the user logged in. pam_sm_close_session is a null function.		
	The following options can be passed to the module:		
	debug syslog(3C) debugging information at the LOG_DEBUG level		
	nowarn Turn off warning messages		
ERRORS	The following values are returned:		
	PAM_SUCCESS Successful completiton		
	PAM_SESSION_ERR Can not make or remove the entry for the specified session		
	PAM_USER_UNKNOWN No account is present for <i>user</i>		
ATTRIBUTES	See attributes(5) for de	scriptions of the following attributes:	

 ATTRIBUTE TYPE
 ATTRIBUTE VALUE

 Interface Stability
 Evolving

 MT Level
 MT-Safe with exceptions

SEE ALSO	<pre>pam(3PAM), pam_authenticate(3PAM), syslog(3C), libpam(3LIB), pam.conf(4), nsswitch.conf(4), attributes(5), pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix(5), pam_unix_account(5), pam_unix_auth(5),</pre>
NOTES	The interfaces in libpam(3LIB) are MT-Safe only if each thread within the multi-threaded application uses its own PAM handle.
	The pam_unix(5) module might not be supported in a future release. Similar functionality is provided by pam_authtok_check(5), pam_authtok_get(5), pam_authtok_store(5), pam_dhkeys(5), pam_passwd_auth(5), pam_unix_account(5), pam_unix_auth(5), and pam_unix_session(5).

prof(5)

NAME	prof – profile within a function
SYNOPSIS	#define MARK #include <prof.h></prof.h>
	<pre>void MARK(name);</pre>
DESCRIPTION	MARK introduces a mark called <i>name</i> that is treated the same as a function entry point. Execution of the mark adds to a counter for that mark, and program-counter time spent is accounted to the immediately preceding mark or to the function if there are no preceding marks within the active function.
	<i>name</i> may be any combination of letters, numbers, or underscores. Each <i>name</i> in a single compilation must be unique, but may be the same as any ordinary program symbol.
	For marks to be effective, the symbol MARK must be defined before the header prof.h is included, either by a preprocessor directive as in the synopsis, or by a command line argument:
	cc -p -DMARK work.c
	If MARK is not defined, the MARK (<i>name</i>) statements may be left in the source files containing them and are ignored. prof -g must be used to get information on all labels.
EXAMPLES	In this example, marks can be used to determine how much time is spent in each loop. Unless this example is compiled with MARK defined on the command line, the marks are ignored.
	<pre>#include <prof.h> work() { int i, j; MARK(loop1); ((i 0 0 i 0 0000 i 0)) { </prof.h></pre>
	<pre>for (i = 0; i < 2000; i++) {</pre>
SEE ALSO	<pre>profil(2), monitor(3C)</pre>

	rbac(5)	
NAME	rbac – role-based access control	
DESCRIPTION	The addition of role-based access control (RBAC) to the Solaris operating environment gives developers the opportunity to deliver fine-grained security in new and modified applications. RBAC is an alternative to the all-or-nothing security model of traditional superuser-based systems. With RBAC, an administrator can assign privileged functions to specific user accounts (or special accounts called roles).	
	There are two ways to give applications privileges:	
	1. Administrators can assign special attributes such as setUID to application binaries (executable files).	
	2. Administrators can assign special attributes such as setUID to applications using execution profiles.	
	Special attribute assignment along with the theory behind RBAC is discussed in detail in "Role Based Access Control" chapter of the <i>System Administration Guide: Security</i> <i>Services</i> . This chapter describes what authorizations are and how to code for them.	
Authorizations	An authorization is a unique string that represents a user's right to perform some operation or class of operations. Authorization definitions are stored in a database called auth_attr(4). For programming authorization checks, only the authorization name is significant.	
	Some typical values in an auth_attr database are shown below.	
	<pre>solaris.jobs.:::Cron and At Jobs::help=JobHeader.html solaris.jobs.grant:::Delegate Cron & At Administration::help=JobsGrant.html solaris.jobs.admin:::Manage All Jobs::help=AuthJobsAdmin.html solaris.jobs.user:::Cron & At User::help=JobsUser.html</pre>	
	Authorization name strings ending with the grant suffix are special authorizations that give a user the ability to delegate authorizations with the same prefix and functional area to other users.	
Creating Authorization Checks	To check authorizations, use the chkauthattr(3SECDB) library function, which verifies whether or not a user has a given authorization. The synopsis is:	
	<pre>int chkauthattr(const char *authname, const char *username);</pre>	
	The chkauthattr() function checks the policy.conf(4), user_attr(4), and prof_attr(4) databases in order for a match to the given authorization.	
	If you are modifying existing code that tests for root UID, you should find the test in the code and replace it with the chkauthattr() function. A typical root UID check is shown in the first code segment below. An authorization check replacing it is shown in the second code segment; it uses the solaris.jobs.admin authorization and a variable called real_login representing the user.	
	EXAMPLE 1 Standard root check	
	<pre>ruid = getuid();</pre>	

rbac(5)

rbac(5)

```
EXAMPLE 1 Standard root check
                               (Continued)
if ((eflag || lflag || rflag) && argc == 1) {
        if ((pwp = getpwnam(*argv)) == NULL)
                crabort(INVALIDUSER);
        if (ruid != 0) {
               if (pwp->pw_uid != ruid)
                        crabort (NOTROOT);
                else
                        pp = getuser(ruid);
        } else
                pp = *argv++;
} else {
EXAMPLE 2 Authorization check
ruid = getuid();
if ((pwp = getpwuid(ruid)) == NULL)
       crabort(INVALIDUSER);
strcpy(real_login, pwp->pw_name);
if ((eflag || lflag || rflag) && argc == 1) {
        if ((pwp = getpwnam(*argv)) == NULL)
                crabort(INVALIDUSER);
        if (!chkauthattr("solaris.jobs.admin", real login)) {
                if (pwp->pw_uid != ruid)
                        crabort (NOTROOT);
                else
                        pp = getuser(ruid);
        } else
                pp = *argv++;
} else {
```

For new applications, find an appropriate location for the test and use chkauthattr() as shown above. Typically the authorization check makes an access decision based on the identity of the calling user to determine if a privileged action (for example, a system call) should be taken on behalf of that user.

Applications that perform a test to restrict who can perform their security-relevant functionality are generally setuid to root. Programs that were written prior to RBAC and that are only available to the root user may not have such checks. In most cases, the kernel requires an effective user ID of root to override policy enforcement. Therefore, authorization checking is most useful in programs that are setuid to root.

For instance, if you want to write a program that allows authorized users to set the system date, the command must be run with an effective user ID of root. Typically, this means that the file modes for the file would be -rwsr-xr-x with root ownership.

Use caution, though, when making programs setuid to root. For example, the effective UID should be set to the real UID as early as possible in the program's initialization function. The effective UID can then be set back to root after the authorization check is performed and before the system call is made. On return from the system call, the effective UID should be set back to the real UID again to adhere to the principle of least privilege.

Another consideration is that LD_LIBRARY path is ignored for setuid programs (see SECURITY section in ld.so.l(1)) and that shell scripts must be modified to work properly when the effective and real UIDs are different. For example, the -p flag in Bourne shell is required to avoid resetting the effective UID back to the real UID.

Using an effective UID of root instead of the real UID requires extra care when writing shell scripts. For example, many shell scripts check to see if the user is root before executing their functionality. With RBAC, these shell scripts may be running with the effective UID of root and with a real UID of a user or role. Thus, the shell script should check euid instead of uid. For example,

```
WHO=`id | cut -f1 -d" "`
if [ ! "$WHO" = "uid=0(root)" ]
then
            echo "$PROG: ERROR: you must be super-user to run this script."
            exit 1
fi
should be changed to
```

Authorizations can be explicitly checked in shell scripts by checking the output of the auths(1) utility. For example,

for auth in `auths | tr , " "` NOTFOUND
do
 ["\$auth" = "solaris.date"] && break # authorization found
done
 if ["\$auth" != "solaris.date"]
 then
 echo >&2 "\$PROG: ERROR: you are not authorized to set the date"
 exit 1
 fi
SEE ALSO
 ld.so.1(1), chkauthattr(3SECDB), auth_attr(4), policy.conf(4),
 prof_attr(4), user_attr(4)

System Administration Guide: Security Services

NAME	noon internationalized basis and outer ded needed oversession matching			
NAME	regex – internationalized basic and extended regular expression matching			
DESCRIPTION	Regular Expressions (REs) provide a mechanism to select specific strings from a set of character strings. The Internationalized Regular Expressions described below differ from the Simple Regular Expressions described on the regexp(5) manual page in the following ways:			
	 both Basic and Extended Regular Expressions are supported 			
	 the Internationalization features—character class, equivalence class, and multi-character collation—are supported. 			
	The Basic Regular Expression (BRE) notation and construction rules described in the BASIC REGULAR EXPRESSIONS section apply to most utilities supporting regular expressions. Some utilities, instead, support the Extended Regular Expressions (ERE) described in the EXTENDED REGULAR EXPRESSIONS section; any exceptions for both cases are noted in the descriptions of the specific utilities using regular expressions. Both BREs and EREs are supported by the Regular Expression Matching interfaces regcomp(3C) and regexec(3C).			
BASIC REGULAR				
EXPRESSIONS BREs Matching a Single Character	A BRE ordinary character, a special character preceded by a backslash, or a period matches a single character. A bracket expression matches a single character or a single collating element. See RE Bracket Expression, below.			
BRE Ordinary Characters				
	The interpretation of an ordinary character preceded by a backslash ($\)$ is undefined, except for:			
	 the characters), (, {, and } the digits 1 to 9 inclusive (see BREs Matching Multiple Characters, below) a character inside a bracket expression. 			
BRE Special Characters	A BRE <i>special character</i> has special properties in certain contexts. Outside those contexts, or when preceded by a backslash, such a character will be a BRE that matches the special character itself. The BRE special characters and the contexts in which they have their special meaning are:			
	. [\ The period, left-bracket, and backslash are special except when used in a bracket expression (see RE Bracket Expression, below). An expression containing a [that is not preceded by a backslash and is not part of a bracket expression produces undefined results.			
	* The asterisk is special except when used:			
	in a bracket expressionas the first character of an entire BRE (after an initial ^, if any)			

	 as the first character of a subexpression (after an initial ^, if any); see BREs Matching Multiple Characters, below.
	^ The circumflex is special when used:
	 as an anchor (see BRE Expression Anchoring, below). as the first character of a bracket expression (see RE Bracket Expression, below).
	\$ The dollar sign is special when used as an anchor.
Periods in BREs	A period (.), when used outside a bracket expression, is a BRE that matches any character in the supported character set except NUL.
RE Bracket Expression	A bracket expression (an expression enclosed in square brackets, []) is an RE that matches a single collating element contained in the non-empty set of collating elements represented by the bracket expression.
	The following rules and definitions apply to bracket expressions:
	1. A <i>bracket expression</i> is either a matching list expression or a non-matching list expression. It consists of one or more expressions: collating elements, collating symbols, equivalence classes, character classes, or range expressions (see rule 7 below). Portable applications must not use range expressions, even though all implementations support them. The right-bracket (]) loses its special meaning and represents itself in a bracket expression if it occurs first in the list (after an initial circumflex (^), if any). Otherwise, it terminates the bracket expression, unless it appears in a collating symbol (such as [.].]) or is the ending right-bracket for a collating symbol, equivalence class, or character class. The special characters:
	* [\ (period, asterisk, left-bracket and backslash, respectively) lose their special meaning within a bracket expression. The character sequences:
	[. [= [: (left-bracket followed by a period, equals-sign, or colon) are special inside a bracket expression and are used to delimit collating symbols, equivalence class expressions, and character class expressions. These symbols must be followed by a valid expression and the matching terminating sequence .], =] or :], as described in the following items.
	2. A <i>matching list</i> expression specifies a list that matches any one of the expressions represented in the list. The first character in the list must not be the circumflex. For example, [abc] is an RE that matches any of the characters a, b or c.
	3. A <i>non-matching list</i> expression begins with a circumflex (^), and specifies a list that matches any character or collating element except for the expressions represented in the list after the leading circumflex. For example, [^abc] is an RE that matches any character or collating element except the characters a, b, or c. The circumflex will have this special meaning only when it occurs first in the list, immediately following the left-bracket.

- 4. A *collating symbol* is a collating element enclosed within bracket-period ([..]) delimiters. Multi-character collating elements must be represented as collating symbols when it is necessary to distinguish them from a list of the individual characters that make up the multi-character collating element. For example, if the string ch is a collating element in the current collation sequence with the associated collating symbol <ch>, the expression [[.ch.]] will be treated as an RE matching the character sequence ch, while [ch] will be treated as an RE matching symbols will be recognized only inside bracket expressions. This implies that the RE [[.ch.]] *c matches the first to fifth character in the string chchch. If the string is not a collating element in the current collating sequence definition, or if the collating element has no characters associated with it, the symbol will be treated as an invalid expression.
- 5. An *equivalence class expression* represents the set of collating elements belonging to an equivalence class. Only primary equivalence classes will be recognised. The class is expressed by enclosing any one of the collating elements in the equivalence class within bracket-equal ([==]) delimiters. For example, if a, and belong to the same equivalence class, then [[=a=]b], [[==]b] and [[==]b] will each be equivalent to [ab]. If the collating element does not belong to an equivalence class, the equivalence class expression will be treated as a *collating symbol*.
- 6. A *character class expression* represents the set of characters belonging to a character class, as defined in the LC_CTYPE category in the current locale. All character classes specified in the current locale will be recognized. A character class expression is expressed as a character class name enclosed within bracket-colon ([::]) delimiters.

The following character class expressions are supported in all locales:

[:alnum:]	[:cntrl:]	[:lower:]	[:space:]
[:alpha:]	[:digit:]	[:print:]	[:upper:]
[:blank:]	[:graph:]	[:punct:]	[:xdigit:]
In addition, character class expressions of the form:			

[:name:]

are recognized in those locales where the *name* keyword has been given a charclass definition in the LC_CTYPE category.

7. A *range expression* represents the set of collating elements that fall between two elements in the current collation sequence, inclusively. It is expressed as the starting point and the ending point separated by a hyphen (–).

Range expressions must not be used in portable applications because their behavior is dependent on the collating sequence. Ranges will be treated according to the current collating sequence, and include such characters that fall within the range based on that collating sequence, regardless of character values. This, however, means that the interpretation will differ depending on collating sequence. If, for instance, one collating sequence defines as a variant of a, while another

	defines it as a letter following z , then the expression $[-z]$ is valid in the first language and invalid in the second.
	In the following, all examples assume the collation sequence specified for the POSIX locale, unless another collation sequence is specifically defined.
	The starting range point and the ending range point must be a collating element or collating symbol. An equivalence class expression used as a starting or ending point of a range expression produces unspecified results. An equivalence class can be used portably within a bracket expression, but only outside the range. For example, the unspecified expression $[[=e=]-f]$ should be given as $[[=e=]e-f]$. The ending range point must collate equal to or higher than the starting range point; otherwise, the expression will be treated as invalid. The order used is the order in which the collating elements are specified in the current collation definition. One-to-many mappings (see locale(5)) will not be performed. For example, assuming that the character eszet is placed in the collation sequence after r and s, but before t, and that it maps to the sequence ss for collation purposes, then the expression $[r-s]$ matches only r and s, but the expression $[s-t]$ matches s, beta, or t.
	The interpretation of range expressions where the ending range point is also the starting range point of a subsequent range expression (for instance $[a-m-o]$) is undefined.
	The hyphen character will be treated as itself if it occurs first (after an initial ^, if any) or last in the list, or as an ending range point in a range expression. As examples, the expressions [-ac] and [ac–] are equivalent and match any of the characters a, c, or –; [^–ac] and [^ac–] are equivalent and match any characters except a, c, or –; the expression [%– –] matches any of the characters between % and – inclusive; the expression [–—@] matches any of the characters between – and @ inclusive; and the expression [a–—@] is invalid, because the letter a follows the symbol – in the POSIX locale. To use a hyphen as the starting range point, it must either come first in the bracket expression or be specified as a collating symbol, for example: [][.–.]–0], which matches either a right bracket or any character or collating element that collates between hyphen and 0, inclusive.
	If a bracket expression must specify both – and], the] must be placed first (after the $^$, if any) and the – last within the bracket expression.
	Note: Latin-1 characters such as $$ or $$ are not printable in some locales, for example, the ja locale.
BREs Matching Multiple Characters	The following rules can be used to construct BREs matching multiple characters from BREs matching a single character:
Characters	1. The concatenation of BREs matches the concatenation of the strings matched by each component of the BRE.
	2. A <i>subexpression</i> can be defined within a BRE by enclosing it between the character pairs \(and \). Such a subexpression matches whatever it would have matched without the \(and \), except that anchoring within subexpressions is optional behavior; see BRE Expression Anchoring, below. Subexpressions can be arbitrarily nested.
	Standarda Environmenta and Maaraa 207

BRE

)		
	characters as was matched by a suppreceding the \n . The character n subexpression (the one that begins corresponding paired $\)$). The expression the \n . For example, the expression adjacent appearances of the match a. The limit of nine back-results the use of a single digit identifier.	atches the same (possibly empty) string of bexpression enclosed between $\langle \text{ and } \rangle$) must be a digit from 1 to 9 inclusive, <i>n</i> th is with the <i>n</i> th $\langle and ends with thepression is invalid if less than n subexpressionsxpression \langle (.* \rangle) \rangle matches a line consistinge same string, and the expression \langle (a \rangle)^* \rangle fails toferences to subexpressions in the RE is based onThis does not imply that only ninea. The following is a valid BRE with ten$
	((((ab))*c))*d)(ef)*(gh))	$2\$ (ij) * (kl) * (mn) * (op) * (qr) *
	followed by the special character a	aracter, a subexpression or a back-reference is asterisk (*), together with that asterisk it matches currences of the BRE would match. For example, lent when matching the string ab.
	followed by an <i>interval expression</i> of with that interval expression it may the BRE would match. The values $0 \le m \le n \le \{\text{RE}_\text{DUP}_\text{MAX}\}$, where occurrences and <i>n</i> specifies the matches exactly <i>m</i> occurrences of the matches exactly <i>m</i> occurrences exactly <i>m</i> occurrences of the matches exactly <i>m</i> occurrences of the matches exactly <i>m</i> occurre	aracter, a subexpression, or a back-reference is of the format $\{m, \}, \{m, \}$ or $\{m,n\}$, together tches what repeated consecutive occurrences of of <i>m</i> and <i>n</i> will be decimal integers in the range <i>m</i> specifies the exact or minimum number of eximum number of occurrences. The expression ces of the preceding BRE, $\{m, \}$ matches at least es any number of occurrences between <i>m</i> and <i>n</i> ,
		ccccccd, the BRE c\{3\} is matched by \(ab\)\{4,\} is not matched at all and the BRE s ten to thirteen.
	The behavior of multiple adjacent dup undefined results.	plication symbols (* and intervals) produces
Precedence	The order of precedence is as shown i	n the following table:
	BRE Precedence (from high to lo	N)
	collation-related bracket symbols	[= =] [: :] []
	escaped characters	\ <special character=""></special>
	bracket expression	
	subexpressions/back-references	() n
	single-character-BRE duplication	* \{ <i>m,n</i> \}
	concatenation	

308 man pages section 5: Standards, Environments, and Macros • Last Revised 12 Jul 1999

	anchoring	^\$
		I
BRE Expression Anchoring	ssion A BRE can be limited to matching strings that begin or end a line; this is called <i>anchoring</i> . The circumflex and dollar sign special characters will be considered B anchors in the following contexts:	
 A circumflex ([^]) is an anchor when used as the first of The implementation may treat circumflex as an anchor character of a subexpression. The circumflex will anchor beginning of a string; only sequences starting at the firs be matched by the BRE. For example, the BRE ^ab matcher abcdef, but fails to match in the string cdefab. A por leading circumflex in a subexpression to match a literal 		umflex as an anchor when used as the first rcumflex will anchor the expression to the es starting at the first character of a string will le, the BRE ^ab matches ab in the string tring cdefab. A portable BRE must escape a
	The implementation may treat a do character of a subexpression. The de	en used as the last character of an entire BRE. llar sign as an anchor when used as the last ollar sign will anchor the expression to the end llar sign can be said to match the end-of-string
	3. A BRE anchored by both ^ and \$ m BRE ^abcdef\$ matches strings cor	atches only an entire string. For example, the nsisting only of abcdef.
	4. ^ and \$ are not special in subexpres	ssions.
	Note: The Solaris implementation does	not support anchoring in BRE subexpressions.
EXTENDED REGULAR EXPRESSIONS	The rules specififed for BREs apply to Extended Regular Expressions (EREs) with the following exceptions:	
	The characters , +, and ? have special meaning, as defined below.	
	 The { and } characters, when used as the duplication operator, are not preceded by backslashes. The constructs \ { and \ } simply match the characters { and }, respectively. 	
	• The back reference operator is not s	upported.
	 Anchoring (^\$) is supported in sub 	expressions.
EREs Matching a Single Character	matches a single character. A bracket ex	naracter preceded by a backslash, or a period xpression matches a single character or a single <i>ingle character</i> enclosed in parentheses matches sees would have matched.
ERE Ordinary Characters	character in the supported character se	tches itself. An ordinary character is any t, except for the ERE special characters listed in e interpretation of an ordinary character ed.

regex(5)

ERE Special Characters	An <i>ERE special character</i> has special properties in certain contexts. Outside those contexts, or when preceded by a backslash, such a character is an ERE that mat the special character itself. The extended regular expression special characters a contexts in which they have their special meaning are:	
	.[\(The period, left-bracket, backslash, and left-parenthesis are special except when used in a bracket expression (see RE Bracket Expression, above). Outside a bracket expression, a left-parenthesis immediately followed by a right-parenthesis produces undefined results.
)	The right-parenthesis is special when matched with a preceding left-parenthesis, both outside a bracket expression.
	* + ? {	The asterisk, plus-sign, question-mark, and left-brace are special except when used in a bracket expression (see RE Bracket Expression, above). Any of the following uses produce undefined results:
		 if these characters appear first in an ERE, or immediately following a vertical-line, circumflex or left-parenthesis if a left-brace is not part of a valid interval expression.
	1	The vertical-line is special except when used in a bracket expression (see RE Bracket Expression, above). A vertical-line appearing first or last in an ERE, or immediately following a vertical-line or a left-parenthesis, or immediately preceding a right-parenthesis, produces undefined results.
	^	The circumflex is special when used:
		 as an anchor (see ERE Expression Anchoring, below). as the first character of a bracket expression (see RE Bracket Expression, above).
	\$	The dollar sign is special when used as an anchor.
Periods in EREs	- · · ·	n used outside a bracket expression, is an ERE that matches any pported character set except NUL.
ERE Bracket Expression	The rules for ERE Bracket Expressions are the same as for Basic Regular Expressions; see RE Bracket Expression, above).	
EREs Matching Multiple Characters	EREs matching a single character:	
Churacters	matched by eac parentheses ma For example, bo	of EREs matches the concatenation of the character sequences th component of the ERE. A concatenation of EREs enclosed in atches whatever the concatenation without the parentheses matches. both the ERE cd and the ERE (cd) are matched by the third and r of the string abcdefabcdef.
	I	

		0 ()	
	2. When an ERE matching a single character or an ERE enclosed in parentheses is followed by the special character plus-sign (+), together with that plus-sign it matches what one or more consecutive occurrences of the ERE would match. For example, the ERE b+ (bc) matches the fourth to seventh characters in the string acabbbcde; [ab] + and [ab] [ab] * are equivalent.		
	what zero or more consecutive occurrent the ERE b*c matches the first character	k (*), together with that asterisk it matches ces of the ERE would match. For example, in the string cabbbcde, and the ERE acters in the string cabbbcdebbbbbbcdbc.	
	4. When an ERE matching a single charact followed by the special character questi- question-mark it matches what zero or o would match. For example, the ERE b?o acabbbcde.	on-mark (?), together with that	
	5. When an ERE matching a single character or an ERE enclosed in parentheses is followed by an <i>interval expression</i> of the format $\{m\}$, $\{m,\}$ or $\{m,n\}$, together with that interval expression it matches what repeated consecutive occurrences of the ERE would match. The values of <i>m</i> and <i>n</i> will be decimal integers in the range $0 \le m \le n \le \{\text{RE}_\text{DUP}_\text{MAX}\}$, where <i>m</i> specifies the exact or minimum number of occurrences and <i>n</i> specifies the maximum number of occurrences. The expression $\{m\}$ matches exactly <i>m</i> occurrences of the preceding ERE, $\{m_i\}$ matches at least <i>m</i> occurrences and $\{m,n\}$ matches any number of occurrences between <i>m</i> and <i>n</i> , inclusive.		
	For example, in the string abababccccccd the ERE c{3} is matched by characters seven to nine and the ERE (ab){2,} is matched by characters one to six.		
	The behavior of multiple adjacent duplicati produces undefined results.	on symbols (+, *, ? and intervals)	
ERE Alternation	Two EREs separated by the special character vertical-line () match a string that is matched by either. For example, the ERE a((bc) d) matches the string abc and the string ad. Single characters, or expressions matching single characters, separated by the vertical bar and enclosed in parentheses, will be treated as an ERE matching a single character.		
ERE Precedence	The order of precedence will be as shown in the following table:		
	ERE Precedence (from high to low)		
	collation-related bracket symbols	[= =] [: :] []	
	escaped characters	<pre>\<special character=""></special></pre>	
	bracket expression	[]	

	grouping	()
	single-character-ERE duplication	$* + ? \{m, n\}$
	concatenation	
	anchoring	^ \$
	alternation	
	For example, the ERE abba cde matche (rather than the string abbade or abbcde, of precedence than alternation).	because concatenation has a higher order
ERE Expression Anchoring	An ERE can be limited to matching strings that begin or end a line; this is called <i>anchoring</i> . The circumflex and dollar sign special characters are considered ERE anchors when used anywhere outside a bracket expression. This has the following effects:	
 A circumflex (^) outside a bracket expression anchors the expression or subexpression it begins to the beginning of a string; such an expression or subexpression can match only a sequence starting at the first character of a s For example, the EREs ^ab and (^ab) match ab in the string abcdef, but fail match in the string cdefab, and the ERE a^b is valid, but can never match be the a prevents the expression ^b from matching starting at the first character 		of a string; such an expression or ce starting at the first character of a string. atch ab in the string abcdef, but fail to a^b is valid, but can never match because
	can match only a sequence ending at the EREs ef\$ and (ef\$) match ef in the string	ing; such an expression or subexpression e last character of a string. For example, the abcdef, but fail to match in the string never match because the f prevents the
SEE ALSO	<pre>localedef(1), regcomp(3C), attribute</pre>	s(5), environ(5), locale(5), regexp(5)

	Tegerp(3)
NAME	regexp, compile, step, advance – simple regular expression compile and match routines
SYNOPSIS	<pre>#define INIT declarations #define GETC(void) getc code #define PEEKC(void) peekc code #define UNGETC(void) ungetc code #define RETURN(ptr) return code #define ERROR(val) error code extern char *loc1, *loc2, *locs;</pre>
	<pre>#include <regexp.h></regexp.h></pre>
	<pre>char *compile(char *instring, char *expbuf, const char *endfug, int</pre>
	<pre>int step(const char *string, const char *expbuf);</pre>
	<pre>int advance(const char *string, const char *expbuf);</pre>
DESCRIPTION	Regular Expressions (REs) provide a mechanism to select specific strings from a set of character strings. The Simple Regular Expressions described below differ from the Internationalized Regular Expressions described on the regex(5) manual page in the following ways:
	 only Basic Regular Expressions are supported
	 the Internationalization features—character class, equivalence class, and multi-character collation—are not supported.
	The functions step(), advance(), and compile() are general purpose regular expression matching routines to be used in programs that perform regular expression matching. These functions are defined by the <regexp.h> header.</regexp.h>
	The functions $step()$ and $advance()$ do pattern matching given a character string and a compiled regular expression as input.
	The function compile() takes as input a regular expression as defined below and produces a compiled expression that can be used with step() or advance().
Basic Regular Expressions	A regular expression specifies a set of character strings. A member of this set of strings is said to be matched by the regular expression. Some characters have special meaning when used in a regular expression; other characters stand for themselves.
	The following one-character REs match a single character:
	1.1 An ordinary character (<i>not</i> one of those discussed in 1.2 below) is a one-character RE that matches itself.
	1.2 A backslash ($\$) followed by any special character is a one-character RE that matches the special character itself. The special characters are:

Standards, Environments, and Macros 313

regexp(5)

	a.	., *, [, and \ (period, asterisk, left square bracket, and backslash, respectively), which are always special, <i>except</i> when they appear within square brackets ([]; see 1.4 below).
	b.	^ (caret or circumflex), which is special at the <i>beginning</i> of an <i>entire</i> RE (see 4.1 and 4.3 below), or when it immediately follows the left of a pair of square brackets ([]]) (see 1.4 below).
	с.	\$ (dollar sign), which is special at the end of an <i>entire</i> RE (see 4.2 below).
	d.	The character used to bound (that is, delimit) an entire RE, which is special for that RE (for example, see how slash (/) is used in the g command, below.)
1.3	A period (new-line.	.) is a one-character RE that matches any character except
1.4	one-charac the first ch matches a string. The minus (-) example, meaning i right squa character of matches ei through f	pty string of characters enclosed in square brackets ([]) is a cter RE that matches <i>any one</i> character in that string. If, however, haracter of the string is a circumflex (^), the one-character RE ny character <i>except</i> new-line and the remaining characters in the e ^ has this special meaning <i>only</i> if it occurs first in the string. The may be used to indicate a range of consecutive characters; for $[0-9]$ is equivalent to $[0123456789]$. The – loses this special f it occurs first (after an initial ^, if any) or last in the string. The re bracket (]) does not terminate such a string when it is the first within it (after an initial ^, if any); for example, []a-f] ither a right square bracket (]) or one of the ASCII letters a inclusive. The four characters listed in 1.2.a above stand for s within such a string of characters.
The follow	ving rules n	nay be used to construct REs from one-character REs:
2.1		A one-character RE is a RE that matches whatever the one-character RE matches.
2.2		A one-character RE followed by an asterisk (*) is a RE that matches 0 or more occurrences of the one-character RE. If there is any choice, the longest leftmost string that permits a match is chosen.
2.3		A one-character RE followed by $\{m\}, \{m, \}$, or $\{m,n\}$ is a RE that matches a <i>range</i> of occurrences of the one-character RE. The values of <i>m</i> and <i>n</i> must be non-negative integers less than 256; $\{m\}$ matches <i>exactly m</i> occurrences; $\{m,n\}$ matches <i>at least m</i> occurrences; $\{m,n\}$ matches <i>any number</i> of occurrences <i>between m</i> and <i>n</i> inclusive. Whenever a choice exists, the RE matches as many occurrences as possible.

	2.4	The concatenation of REs is a RE that matches the concatenation of the strings matched by each component of the RE.
	2.5	A RE enclosed between the character sequences \setminus (and \setminus) is a RE that matches whatever the unadorned RE matches.
	2.6	The expression $\ n$ matches the same string of characters as was matched by an expression enclosed between $\ (and \)$ <i>earlier</i> in the same RE. Here <i>n</i> is a digit; the sub-expression specified is that beginning with the <i>n</i> -th occurrence of $\ (counting from the left. For example, the expression ^{\land} (c. * \) \ 1 \ $ matches a line consisting of two repeated appearances of the same string.$
	An RE may be constrained to	match words.
	3.1	< constrains a RE to match the beginning of a string or to follow a character that is not a digit, underscore, or letter. The first character matching the RE must be a digit, underscore, or letter.
	3.2	> constrains a RE to match the end of a string or to precede a character that is not a digit, underscore, or letter.
	An <i>entire RE</i> may be constrain line (or both).	ed to match only an initial segment or final segment of a
	4.1	A circumflex (^) at the beginning of an entire RE constrains that RE to match an <i>initial</i> segment of a line.
	4.2	A dollar sign (\$) at the end of an entire RE constrains that RE to match a <i>final</i> segment of a line.
	4.3	The construction $^{entire} RE \ddagger$ constrains the entire RE to match the entire line.
	The null RE (for example, //) is equivalent to the last RE encountered.
Addressing with	Addresses are constructed as f	follows:
REs	1. The character "." addresses	s the current line.
	2. The character "\$" addresses	s the last line of the buffer.
	3. A decimal number <i>n</i> addre	sses the <i>n</i> -th line of the buffer.
		ed with the mark name character <i>x,</i> which must be an z). Lines are marked with the k command described

cgcxp(0)	
	5. A RE enclosed by slashes (/) addresses the first line found by searching <i>forward</i> from the line <i>following</i> the current line toward the end of the buffer and stopping at the first line containing a string matching the RE. If necessary, the search wraps around to the beginning of the buffer and continues up to and including the current line, so that the entire buffer is searched.
	6. A RE enclosed in question marks (?) addresses the first line found by searching <i>backward</i> from the line <i>preceding</i> the current line toward the beginning of the buffer and stopping at the first line containing a string matching the RE. If necessary, the search wraps around to the end of the buffer and continues up to and including the current line.
	 An address followed by a plus sign (+) or a minus sign (-) followed by a decimal number specifies that address plus (respectively minus) the indicated number of lines. A shorthand for .+5 is .5.
	8. If an address begins with + or –, the addition or subtraction is taken with respect to the current line; for example, –5 is understood to mean .–5.
	9. If an address ends with + or -, then 1 is added to or subtracted from the address, respectively. As a consequence of this rule and of Rule 8, immediately above, the address - refers to the line preceding the current line. (To maintain compatibility with earlier versions of the editor, the character ^ in addresses is entirely equivalent to) Moreover, trailing + and - characters have a cumulative effect, so refers to the current line less 2.
	10. For convenience, a comma (,) stands for the address pair 1, \$, while a semicolon (;) stands for the pair ., \$.
Characters With Special Meaning	Characters that have special meaning except when they appear within square brackets ([]) or are preceded by $\ are: ., *, [, \ .\ Other special characters, such as $ have special meaning in more restricted contexts.$
	The character ^ at the beginning of an expression permits a successful match only immediately after a newline, and the character \$ at the end of an expression requires a trailing newline.
	Two characters have special meaning only when used within square brackets. The character – denotes a range, [$c-c$], unless it is just after the open bracket or before the closing bracket, [$-c$] or [$c-$] in which case it has no special meaning. When used within brackets, the character ^ has the meaning <i>complement of</i> if it immediately follows the open bracket (example: [cc]); elsewhere between brackets (example: [c^2]) it stands for the ordinary character ^.
	The special meaning of the \backslash operator can be escaped only by preceding it with another \backslash , for example $\backslash \backslash$.
Macros	Programs must have the following five macros declared before the #include <regexp.h> statement. These macros are used by the compile() routine. The macros GETC, PEEKC, and UNGETC operate on the regular expression given as input to compile().</regexp.h>

	GETC	This macro returns the value of the next character (byte) in the regular expression pattern. Successive calls to GETC should return successive characters of the regular expression.		
	PEEKC	This macro returns the next character (byte) in the regular expression. Immediately successive calls to PEEKC should return the same character, which should also be the next character returned by GETC.		
	UNGETC	This macro causes the argument c to be returned by the next call to GETC and PEEKC. No more than one character of pushback is ever needed and this character is guaranteed to be the last character read by GETC. The return value of the macro UNGETC(c) is always ignored.		
	RETURN (<i>ptr</i>)	This macro is used on normal exit of the compile() routine. The value of the argument <i>ptr</i> is a pointer to the character after the last character of the compiled regular expression. This is useful to programs which have memory allocation to manage.		
	ERROR (val)	This macro is the abnormal return from the compile() routine. The argument <i>val</i> is an error number (see ERRORS below for meanings). This call should never return.		
compile()	compile() routine is as follows:			
	compile(<i>instring</i> , <i>expbuf</i> , <i>endbuf</i> , <i>eof</i>)			
	The first parameter, <i>instring</i> , is never used explicitly by the compile() routine but is useful for programs that pass down different pointers to input characters. It is sometimes used in the INIT declaration (see below). Programs which call functions to input characters or have characters in an external array can pass down a value of (char *) 0 for this parameter.			
	The next parameter, <i>expbuf</i> , is a character pointer. It points to the place where the compiled regular expression will be placed.			
	The parameter <i>endbuf</i> is one more than the highest address where the compiled regular expression may be placed. If the compiled expression cannot fit in (endbuf-expbuf) bytes, a call to ERROR (50) is made.			
	The parameter <i>eof</i> is the character which marks the end of the regular expression. This character is usually a $/$.			
	Each program that includes the <regexp.h> header file must have a #define statement for INIT. It is used for dependent declarations and initializations. Most often it is used to set a register variable to point to the beginning of the regular expression so that this register variable can be used in the declarations for GETC, PEEKC, and UNGETC. Otherwise it can be used to declare external variables that might be used by GETC, PEEKC and UNGETC. (See EXAMPLES below.)</regexp.h>			

regexp(5)			
<pre>step(), advance()</pre>	The first parameter to the step() and advance() functions is a pointer to a string of characters to be checked for a match. This string should be null terminated.		
	The second parameter, <i>expbuf</i> , is the compiled regular expression which was obtained by a call to the function compile().		
	The function step() returns non-zero if some substring of <i>string</i> matches the regular expression in <i>expbuf</i> and 0 if there is no match. If there is a match, two external character pointers are set as a side effect to the call to step(). The variable loc1 points to the first character that matched the regular expression; the variable loc2 points to the character after the last character that matches the regular expression. Thus if the regular expression matches the entire input string, loc1 will point to the first character of <i>string</i> and loc2 will point to the null at the end of <i>string</i> .		
	The function advance() returns non-zero if the initial substring of <i>string</i> matches the regular expression in <i>expbuf</i> . If there is a match, an external character pointer, loc2, is set as a side effect. The variable loc2 points to the next character in <i>string</i> after the last character that matched.		
	When $advance()$ encounters $a * or \setminus \{ \ \}$ sequence in the regular expression, it will advance its pointer to the string to be matched as far as possible and will recursively call itself trying to match the rest of the string to the rest of the regular expression. As long as there is no match, $advance()$ will back up along the string until it finds a match or reaches the point in the string that initially matched the $* \text{ or } \setminus \{ \ \}$. It is sometimes desirable to stop this backing up before the initial point in the string is reached. If the external character pointer locs is equal to the point in the string at sometime during the backing up process, $advance()$ will break out of the loop that backs up and will return zero.		
	The external variables circf, sed, and nbra are reserved.		
EXAMPLES	EXAMPLE 1 The following is an example of how the regular expression macros and calls might be defined by an application program:		
	<pre>#define INIT register char *sp = instring; #define GETC (*sp++) #define PEEKC (*sp) #define UNGETC(c) (sp) #define RETURN(*c) return; #define ERROR(c) regerr #include <regexp.h> (void) compile(*argv, expbuf, &expbuf[ESIZE],'\0'); if (step(linebuf, expbuf))</regexp.h></pre>		
	succeed;		
DIAGNOSTICS	The function compile() uses the macro RETURN on success and the macro ERROR on failure (see above). The functions step() and advance() return non-zero on a successful match and zero if there is no match. Errors are:		

	11	range endpoint too large.
	16	bad number.
	25	\ <i>digit</i> out of range.
	36	illegal or missing delimiter.
	41	no remembered search string.
	42	() imbalance.
	43	too many \setminus (.
	44	more than 2 numbers given in $\setminus \{ \setminus \}$.
	45	$\}$ expected after \backslash .
	46	first number exceeds second in $\setminus \{ \ \}$.
	49	[] imbalance.
	50	regular expression overflow.
1150	(E)	

SEE ALSO | regex(5)

Standards, Environments, and Macros 319

SEAM(5)

NAME | SEAM - overview of Sun Enterprise Authentication Mechanism

DESCRIPTION

SEAM (Sun Enterprise Authentication Mechanism) authenticates clients in a network environment, allowing for secure transactions. (A client may be a user or a network service) SEAM validates the identity of a client and the authenticity of transferred data. SEAM is a *single-sign-on* system, meaning that a user needs to provice a password only at the beginning of a session. SEAM is based on the KerberosTM system developed at MIT, and is compatible with Kerberos V5 systems over heterogeneous networks.

SEAM works by granting clients *tickets*, which uniquely identify a client, and which have a finite lifetime. A client possessing a ticket is automatically validated for network services for which it is entitled; for example, a user with a valid SEAM ticket may rlogin into another machine running SEAM without having to identify itself. Because each client has a unique ticket, its identity is guaranteed.

To obtain tickets, a client must first initialize the SEAM session, either by using the kinit(1) command or a PAM module. (See pam_krb5(5)). kinit prompts for a password, and then communicates with a *Key Distribution Center* (KDC). The KDC returns a *Ticket-Granting Ticket* (TGT) and prompts for a confirmation password. If the client confirms the password, it can use the Ticket-Granting Ticket to obtain tickets for specific network services. Because tickets are granted transparently, the user need not worry about their management. Current tickets may be viewed by using the klist(1) command.

Tickets are valid according to the system *policy* set up at installation time. For example, tickets have a default lifetime for which they are valid. A policy may further dictate that privileged tickets, such as those belonging to root, have very short lifetimes. Policies may allow some defaults to be overruled; for example, a client may request a ticket with a lifetime greater or less than the default.

Tickets can be renewed using kinit. Tickets are also *forwardable*, allowing you to use a ticket granted on one machine on a different host. Tickets can be destroyed by using kdestroy(1). It is a good idea to include a call to kdestroy in your .logout file.

Under SEAM, a client is referred to as a *principal*. A principal takes the following form:

primary/instance@REALM

primary

instance

A user, a host, or a service.

A qualification of the primary. If the primary is a host — indicated by the keyword host— then the instance is the fully-qualified domain name of that host. If the primary is a user or service, then the instance is optional. Some instances, such as admin or root, are privileged.

By taking advantage of the General Security Services API (GSS-API), SEAM offers, besides user authentication, two other types of security service: integrity, which authenticates the validity of transmitted data, and privacy, which encrypts transmitted data. Developers can take advantage of the GSS-API through the use of the RPCSEC_GSS API interface (see rpcsec_gss(3NSL)). EXAMPLESEXAMPLE 1 Examples of valid principal names The following are examples of valid principal names: joe joe/admin joe@ENG.ACME.COM rlogin/bigmachine.eng.acme.com@ENG.ACME.COM host/bigmachine.eng.acme.com@ENG.ACME.COMThe first four cases are user principals. In the first two cases, it is assumed that the user joe is in the same realm as the client, so no realm is specified. Note that joeand joe/admin has
The following are examples of valid principal names: joe joe/admin joe@ENG.ACME.COM joe/admin@ENG.ACME.COM rlogin/bigmachine.eng.acme.com@ENG.ACME.COM host/bigmachine.eng.acme.com@ENG.ACME.COM The first four cases are <i>user principals</i> . In the first two cases, it is assumed that the user joe is in the same realm as the client, so no realm is specified. Note that joeand
joe joe/admin joe@ENG.ACME.COM joe/admin@ENG.ACME.COM rlogin/bigmachine.eng.acme.com@ENG.ACME.COM host/bigmachine.eng.acme.com@ENG.ACME.COM The first four cases are <i>user principals</i> . In the first two cases, it is assumed that the user joe is in the same realm as the client, so no realm is specified. Note that joeand
<pre>joe/admin joe@ENG.ACME.COM joe/admin@ENG.ACME.COM rlogin/bigmachine.eng.acme.com@ENG.ACME.COM host/bigmachine.eng.acme.com@ENG.ACME.COM The first four cases are user principals. In the first two cases, it is assumed that the user joe is in the same realm as the client, so no realm is specified. Note that joeand</pre>
joe is in the same realm as the client, so no realm is specified. Note that joeand
different privileges from joe. The fifth case is a <i>service principal</i> , while the final case is a <i>host principal</i> . The word host is required for host principals. With host principals, the instance is the fully qualified hostname. Note that the words admin and host are reserved keywords.
SEE ALSO kdestroy(1), kinit(1), klist(1), kpasswd(1), krb5.conf(5)
Sun Enterprise Authentication Mechanism Guide
NOTES If you enter your username and kinit responds with this message:
Principal unknown (kerberos) you haven't been registered as a SEAM user. See your system administrator or the <i>Sun Enterprise Authentication Mechanism Guide</i> .

sgml	(5)
------	-----

NAMEsgml, solbook – Standard Generalized Markup LanguageDESCRIPTIONStandard Generalized Markup Language (SGML) is the ISO standard 8879:1986 that describes a syntax for marking up documents with tags that describe the purpose of the text rather than the appearance on the page. This form of morkup facilitates document interchange between different platforms and applications. SCML allows the management of information as data objects rather than text on a page.In an SGML document the main structural components are called elements. The organization and structure of a document and the meaning of elements are described in the Document Type Definition (DTD). Elements are the <i>tage</i> that identify the content. Element names may be descriptive of the content for case of use. For example -paras for paragraphs. Elements can have attributes which are used to modify or refine the properties or characteristics of the element. Within the DTD a valid context for each element is defined and a framework is provided for the types of elements that constitute a compliant document.Another component of the DTD is entities. Entities are a collection of characters that can be referenced as a unit. Entities are similar to constants in a programming language such as C. They can be defined and referenced. An entity can represent one character or symbol which does not appear on a standard keyboard, a word or group of words, or an entire separate sgml marked-up file. Entities allow reuse of standard text.SolBook ElementsElements are defined with a hierarchical structure that gives a structure to the document. The following is a description of Some of the elements from the SolBook DTD, which is a proper subset of DocBook DTD, is used when writing reference marula pages. The SolBook DTD contains a number of tags that are designed for the unique needs o	sgmi(5)				
describes a syntax for marking up documents with tags that describe the purpose of the text rather than the appearance on the page. This form of markup facilitates document interchange between different platforms and applications. SGML allows the management of information as data objects rather than text on a page.In an SGML document the main structural components are called elements. The organization and structure of a document and the meaning of elements are described in the Document Type Definition (DTD). Elements are the <i>tags</i> that identify the content. Element names may be descriptive of the content for case of use. For example e-paira> for paragraphs. Elements can have attributes which are used to modify or refine the properties or characteristics of the element. Within the DTD a valid context for each element is defined and a framework is provided for the types of elements that constitute a compliant document.Another component of the DTD is entities. Entities are a collection of characters that can be referenced as a unit. Entities are similar to constants in a programming language such as C. They can be defined and referenced. An entity can represent one character or symbol which does not appear on a standard keyboard, a word or group of words, or an entire separate sgml marked-up file. Entities allow reuse of standard text.SolBook ElementsElements are defined with a hierarchical structure that gives a structure to the document. The following is a description of some of the elements from the SolBook DTD which are used for reference pages.DOCTYPEThe first line in an SGML file that identifies the location of the DTD that is used to define the document. The <1D0CTYPE string is what the SGML-aware man(1) command uses to identify that a file is formatted in SCML rather than nrcff(1).RefIntryThe top lay	NAME	sgml, solbook – Standard Generalized Markup Language			
organization and structure of a document and the meaning of elements are described in the Document Type Definition (DTD). Elements are the <i>tags</i> that identify the content. Element names may be descriptive of the content for ease of use. For example <paras are="" attributes="" can="" elements="" for="" have="" modify="" or<br="" paragraphs.="" to="" used="" which=""></paras> refine the properties or characteristics of the element. Within the DTD a valid context for each element is defined and a framework is provided for the types of elements that constitute a compliant document.Another component of the DTD is entities. Entities are a collection of characters that can be referenced as a unit. Entities are similar to constants in a programming language such as C. They can be defined and referenced. An entity can represent one character or symbol which does not appear on a standard keyboard, a word or group of words, or an entire separate sgml marked-up file. Entities allow reuse of standard text.SolBook ElementsElements are defined with a hierarchical structure that gives a structure to the document. The following is a description of some of the elements from the SolBook DTD which is a used for reference pages.SolBook ElementsElements are defined with a hierarchical structure that gives a structure to the document. The following is a description of some of the elements from the SolBook DTD which are used for reference pages.SolBook ElementsRefementyRefEntryThe top layer element that contains a reference page is <refentrys. all="" of="" text<br="" the=""></refentrys.> and other tags must be contained within this tag.RefMetaThe next tag in a reference page is <refentrys. all="" of="" text<br="" the=""></refentrys.> and other tags must be contained within this tag.referenceThe inst is the title of the reference page. It is equivalent to the name of the refe	DESCRIPTION	describes a syntax the text rather than document intercha	for marking up documents with tags that describe the purpose of a the appearance on the page. This form of markup facilitates ange between different platforms and applications. SGML allows the		
that can be referenced as a unit. Entities are similar to constants in a programming language such as C. They can be defined and referenced. An entity can represent one character or symbol which does not appear on a standard keyboard, a word or group of words, or an entire separate sgml marked-up file. Entities allow reuse of standard text.There is no single standard DTD , but the de facto standard for the computer industry is the DocBook DTD , developed and maintained by the Davenport Group. Within Sun, the SolBook DTD , which is a proper subset of DocBook DTD , is used when writing reference manual pages. The SolBook DTD contains a number of tags that are designed for the unique needs of the reference pages.SolBook ElementsElements are defined with a hierarchical structure that gives a structure to the document. The following is a description of some of the elements from the SolBook DTD which are used for reference pages.DOCTYPEThe first line in an SGML file that identifies the location of the DTD that is used to define the document. The DCTYPE string is what the SGML -aware man(1)<br/ command uses to identify that a file is formatted in SGML rather than nroff(1).RefEntryThe top layer element that contains a reference page is <refentry>. All of the text and other tags must be contained within this tag.RefMetaThe next tag in a reference page is <refmeta>, which is a container for several other tags. They are:<rd>This is the title of the reference page. It is equivalent to the name of the reference page's file name, without the section number extension.</br></br></rd></br></refmeta></refentry>		organization and structure of a document and the meaning of elements are described in the Document Type Definition (DTD). Elements are the <i>tags</i> that identify the content. Element names may be descriptive of the content for ease of use. For example <para> for paragraphs. Elements can have attributes which are used to modify or refine the properties or characteristics of the element. Within the DTD a valid context for each element is defined and a framework is provided for the types of elements that constitute a compliant document.</para>			
is the DocBook DTD , developed and maintained by the Davenport Group. Within Sun, the SolBook DTD , which is a proper subset of DocBook DTD , is used when writing reference manual pages. The SolBook DTD contains a number of tags that are designed for the unique needs of the reference pages.SolBook ElementsElements are defined with a hierarchical structure that gives a structure to the document. The following is a description of some of the elements from the SolBook DTD which are used for reference pages.DOCTYPEThe first line in an SGML file that identifies the location of the DTD that is used to define the document. The <1DOCTYPE string is what the SGML -aware man(1) command uses to identify that a file is formatted in SGML rather than nroff(1).RefEntryThe top layer element that contains a reference page is <refentry>. All of the text and other tags must be contained within this tag.RefMetaThe next tag in a reference page is <refmeta>, which is a container for several other tags. They are: <re>Crefentrytitle>This is the title of the reference page. It is equivalent to the name of the reference page's file name, without the section number extension.</re></refmeta></refentry>		that can be referenced as a unit. Entities are similar to constants in a programming language such as C. They can be defined and referenced. An entity can represent one character or symbol which does not appear on a standard keyboard, a word or group of words, or an entire separate sgml marked-up file. Entities allow reuse of standard text.			
document. The following is a description of some of the elements from the SolBook DTD which are used for reference pages.DOCTYPEThe first line in an SGML file that identifies the location of the DTD that is used to define the document. The string is what the SGML -aware man(1)<br/ command uses to identify that a file is formatted in SGML rather than nroff(1).RefEntryThe top layer element that contains a reference page is <refentry>. All of the text and other tags must be contained within this tag.RefMetaThe next tag in a reference page is <refmeta>, which is a container for several other tags. They are:This is the title of the reference page. It is equivalent to the name of the reference page's file name, without the section number extension.This is the section number that the reference page resides in. The</refmeta></refentry>		is the DocBook DTD , developed and maintained by the Davenport Group. Within Sun, the SolBook DTD , which is a proper subset of DocBook DTD , is used when writing reference manual pages. The SolBook DTD contains a number of tags that are			
define the document. The string is what the SGML -aware man(1)<br/ command uses to identify that a file is formatted in SGML rather than nroff(1).RefEntryThe top layer element that contains a reference page is <refentry>. All of the text and other tags must be contained within this tag.RefMetaThe next tag in a reference page is <refmeta>, which is a container for several other tags. They are:This is the title of the reference page. It is equivalent to the name of the reference page's file name, without the section number extension.<</refmeta></refentry>	SolBook Elements	document. The foll	lowing is a description of some of the elements from the SolBook		
and other tags must be contained within this tag. RefMeta The next tag in a reference page is <refmeta>, which is a container for several other tags. They are: <refentrytitle> This is the title of the reference page. It is equivalent to the name of the reference page's file name, without the section number extension. <manvolnum> This is the section number that the reference page resides in. The</manvolnum></refentrytitle></refmeta>	DOCTYPE	define the docume	nt. The $< !DOCTYPE$ string is what the SGML -aware man(1)		
tags. They are: If the title of the reference page. It is equivalent to the name of the reference page's file name, without the section number extension. <manvolnum> This is the section number that the reference page resides in. The</manvolnum>	RefEntry				
 the reference page's file name, without the section number extension. <manvolnum> This is the section number that the reference page resides in. The</manvolnum> 	RefMeta		eference page is <refmeta>, which is a container for several other</refmeta>		
1 0		<refentrytitle></refentrytitle>	the reference page's file name, without the section number		
· ·		<manvolnum></manvolnum>			

			sgml(5)
	<refmiscinfo></refmiscinfo>	information. Meta	nore <refmiscinfo> tags which contain <i>meta</i> information is information about the reference scinfo> tag has the class attribute. There are re routinely used.</refmiscinfo>
		date	This is the date that the file was last modified. By consensus this date is changed only when the technical information on the page changes and not simply for an editorial change.
		sectdesc	This is the section title of the reference page; for example User Commands. The value of this attribute may be a text entity reference.
		software	This is the name of the software product that the topic discussed on the reference page belongs to. For example UNIX commands are part of the SunOS x.x release. The value of this attribute may be a text entity reference.
		arch	This is the architectural platform limitation of the subject discussed on the reference page. If there are no limitations the value used is generic. Other values are sparc and x86.
		copyright	This attribute contains the Sun Microsystems copyright. Any other copyrights that may pertain to the individual reference page file should be entered as separate <refmiscinfo> entries. The value of this attribute may be a text entity reference.</refmiscinfo>
RefNameDiv	reference page. <r< th=""><th colspan="2">ins the equivalent information to the .TH macro line in an nroff(1) 2. <refnamediv> contains three tags. These tags contain the text that is er the '-' (dash) on the NAME line.</refnamediv></th></r<>	ins the equivalent information to the .TH macro line in an nroff(1) 2. <refnamediv> contains three tags. These tags contain the text that is er the '-' (dash) on the NAME line.</refnamediv>	
	<refname></refname>	There may be mor <refname> must <refentrytitle each is separated b</refentrytitle </refname>	es of the topics that are discussed in the file. e than one <refname> for a page. The first match the name of the file and the e>. If there are more than one <refname> tags, by a ',' (comma). The comma is generated by the files, so it should not be typed. This is referred to ext.</refname></refname>
	<refpurpose></refpurpose>	This is a short sum the reference page	dash on the NAME line is contained in this tag. nmary of what the object or objects described on do or are used for. The dash is also d should not be typed in.
	<refdiscriptor></refdiscriptor>		<refentrytitle> is a general topic descriptor ed objects that are discussed on the same page.</refentrytitle>

sgml(5)

		In this case the first tag after the <refnamediv> is a <refdiscriptor>. The <refname> tags follow. Only one <refdiscriptor> is allowed, and it should match the <refentrytitle>.</refentrytitle></refdiscriptor></refname></refdiscriptor></refnamediv>	
RefSynopsisDiv	The SYNOPSIS line of the reference page is contained by this tag. There is a <title> that usually contains an entity reference. The text is the word SYNOPSIS. There are several tags within <refsynopsisdiv> that are designed specifically for the type of synopsis that is used in the different reference page sections. The three types are:</th></tr><tr><th></th><th><cmdsynopsis></th><th>Used for commands and utilities pages.</th></tr><tr><th></th><th><funcsynopsis></th><th>Used for programming interface pages.</th></tr><tr><th></th><th><synopsis></th><th>Used for pages that do not fall into the other two categories.</th></tr><tr><th>RefSect1</th><th colspan=3>This tag is equivalent to the .SH nroff macro. It contains a <title> element that is the title of the reference page section. Section names are the standard names such as DESCRIPTION, OPTIONS, PARAMETERS, SEE ALSO, and others. The contents of the <title> may be a text entity reference.</th></tr><tr><th>RefSect2</th><th colspan=3>This tag is equivalent to the .SS nroff macro. It contains a <title> element that contains the text of the sub-section heading. <refsect2> tags may also be used within a <refsynopsisdiv> as a sub-section heading for the SYNOPSIS section.</th></tr><tr><th>Block Elements</th><th colspan=3>There are a number of block elements that are used for grouping text. This is a list of some of these elements.</th></tr><tr><th></th><th><para></th><th>This tag is used to contain a paragraph of text.</th></tr><tr><th></th><th><variablelist></th><th>This tag is used to create two column lists. For example
descriptions for command options, where the first
column lists the option and the second column
describes the option.</th></tr><tr><th></th><th><orderedlist></th><th>An list of items in a specific order.</th></tr><tr><th></th><th><itemizedlist></th><th>A list of items that are marked with a character such as a bullet or a dash.</th></tr><tr><th></th><th><literallayout></th><th>Formatted program output as produced by a program
or command. This tag is a container for lines set off
from the main text in which line breaks, tabs, and
leading white space are significant.</th></tr><tr><th></th><th><programlisting></th><th>A segment of program code. Line breaks and leading white space are significant.</th></tr><tr><th></th><th></th><th>This tag contains the layout and content for tabular formatting of information. has a required <title>.</th></tr><tr><th></th><th></th><th></th></tr></tbody></table></title>		

		sgnu(5)
	<informaltable></informaltable>	This tag is the same as the tag except the <title> is not required.</th></tr><tr><th></th><th><example></th><th>This tag contains examples of source code or usage of commands. It contains a required <title>.</th></tr><tr><th></th><th><informalexample></th><th>This tag is the same as the <example> tag except the <title> is not required.</th></tr><tr><th>Inline Elements</th><th>The inline elements are used t</th><th>for tagging text.</th></tr><tr><th></th><th><command></th><th>An executable program or the entry a user makes to execute a command.</th></tr><tr><th></th><th><function></th><th>A subroutine in a program or external library.</th></tr><tr><th></th><th><literal></th><th>Contains any literal string.</th></tr><tr><th></th><th><parameter></th><th>An argument passed to a computer program by a function or routine.</th></tr><tr><th></th><th><inlineequation></th><th>An untitled mathematical equation occurring in-line.</th></tr><tr><th></th><th><link></th><th>A hypertext link to text within a book, in the case of the reference manual it is used to cross reference to another reference page.</th></tr><tr><th></th><th><olink></th><th>A hypertext link used to create cross references to books other than the reference manual.</th></tr><tr><th></th><th><xref></th><th>A cross reference to another part of the same reference page.</th></tr><tr><th>SEE ALSO</th><th><pre>man(1), nroff(1), man(5)</pre></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr><tr><th></th><th></th><th></th></tr></tbody></table></title>

sgml(5)

smartcard(5)

NAME | smartcard – overview of smart card features on Solaris

DESCRIPTION

N The smart card framework provides a mechanism to abstract the details of interacting with smart cards and smart cardreaders (called card terminals). The framework is based on the OpenCard Framework V1.1 (OCF) with Sun extensions to allow OCF to operate in a multi-user environment. The core OCF software protocol stack is implemented as a system service daemon. This implementation allows smart cards and card terminals to be shared cooperatively among many different clients on the system while providing access control to the smart card and card terminal resources on a per-UID basis.

An event dispatcher is provided to inform clients of events occuring on the card and at the card terminal, such as card insertion and card removal.

A high-level authentication mechanism is provided to allow clients to perform smart card-based authentications without requiring knowledge of specific card or reader authentication features.

A set of applet administration tools is provided for JavaCards that support downloading Java applets (although applet build tools are not provided).

Administration of the smart card framework is provided with the smartcard(1M) command line administration utility and the smartcardguiadmin(1) GUI administration tool.

Support for several card terminals is provided:

- Sun External Smart Card Reader I (see ocf escr1(7D))
- Sun Internal Smart Card Reader I (see ocf iscr1(7D))
- Dallas iButton Serial Reader (see ocf ibutton(7D))

Additional card terminals can be supported by implementing smart card terminal interfaces in a shared library.

Support for several smart cards is provided:

- Schlumberger Cyberflex Access JavaCard
- Schlumberger MicroPayflex
- Dallas Semiconductor Java iButton JavaCard

Each of the supported cards has a complete set of OCF card services that implement the necessary functionality for authentication and secure storage of data. For the two supported JavaCards, an authentication and secure data storage applet is provided that can be loaded into these cards with the supplied applet administration tools. See smartcard(1M).

A PAM smart card module is provided to allow PAM clients to use smart card-based authentication. See pam smartcard(5)

smartcard(5)

CDE is able to use the PAM smart card module for dtlogin and dtsession authentication. CDE also uses the smart card framework event dispatcher to listen for events on the card terminal and provide corresponding visual feedback to the user.

SEE ALSO ocfserv(1M), smartcard(1M), libsmartcard(3LIB), pam_start(3PAM), pam_smartcard(5), ocf_escr1(7D), ocf_ibutton(7D), ocf_iscr1(7D), scmi2c(7D)

NAME standards, ANSI, C, C++, ISO, POSIX, POSIX.1, POSIX.2, SUS, SUSv2, SVID, SVID3, XNS, XNS4, XNS5, XPG, XPG3, XPG4, XPG4v2 – standards and specifications supported by Solaris

DESCRIPTION Solaris 9supports IEEE Std 1003.1 and IEEE Std 1003.2, commonly known as POSIX.1 and POSIX.2, respectively. The following table lists each version of these standards with a brief description and the SunOS or Solaris release that first conformed to it.

POSIX Standard	Description	Release
POSIX.1-1988	system interfaces and headers	SunOS 4.1
POSIX.1-1990	POSIX.1-1988 update	Solaris 2.0
POSIX.1b-1993	realtime extensions	Solaris 2.4
POSIX.1c-1996	threads extensions	Solaris 2.6
POSIX.2-1992	shell and utilities	Solaris 2.5
POSIX.2a-1992	interactive shell and utilities	Solaris 2.5

Solaris 9also supports the X/Open Common Applications Environment (CAE) Portability Guide Issue 3 (XPG3) and Issue 4 (XPG4), Single UNIX Specification (SUS, also known as XPG4v2), and Single UNIX Specification, Version 2 (SUSv2). Both XPG4 and SUS include Networking Services Issue 4 (XNS4). SUSv2 includes Networking Services Issue 5 (XNS5).

The following table lists each X/Open specification with a brief description and the SunOS or Solaris release that first conformed to it.

Description	Release
superset of POSIX.1-1988 containing utilities from SVID3	SunOS 4.1
superset of POSIX.1-1990, POSIX.2-1992, and POSIX.2a-1992 containing extensions to POSIX standards from XPG3	Solaris 2.4
superset of XPG4 containing historical BSD interfaces widely used by common application packages	Solaris 2.6
sockets and XTI interfaces	Solaris 2.6
superset of SUS extended to support POSIX.1b-1993, POSIX.1c-1996, and ISO/IEC 9899 (C Standard) Amendment 1	Solaris 7
	superset of POSIX.1-1988 containing utilities from SVID3 superset of POSIX.1-1990, POSIX.2-1992, and POSIX.2a-1992 containing extensions to POSIX standards from XPG3 superset of XPG4 containing historical BSD interfaces widely used by common application packages sockets and XTI interfaces superset of SUS extended to support POSIX.1b-1993, POSIX.1c-1996, and

	X/Open CAE Specification	Description	Release
	XNS5	superset and LP64-clean derivative of XNS4.	Solaris 7
	not be used for LP LP64-clean interfac	ation is safe for use only in ILP32 (32-b 64 (64-bit) application environments. U ces that are portable across ILP32 and 1 9 support both the ILP32 and ILP64 er	Jse XNS5, which has LP64 environments. Solaris
	Solaris releases 7 tl 98 Product Standar	hrough 9 have been branded to conformed.	m to The Open Group's UNIX
	Interface Definition since the develope business and since	through 9 support the interfaces speci n, Third Edition, Volumes 1 through 4 rs of this specification (UNIX Systems this specification defers to POSIX and greement about what is currently requi	(SVID3). Note, however, that Laboratories) are no longer in X/Open CAE specifications,
	support the ANSI	hop Compiler™ C 4.2 is installed, So X3.159-1989 Programming Language - guage - C (C) interfaces.	
		hop Compiler™ C 5.0 is installed, So IEC 9899 Amendment 1: C Integrity.	olaris releases 7 through 9
	9 support ISO/IEC of that standard ar standard adopted	hop Compiler C++ 5.0 is installed, C 14882:1998 Programming Languages e described in the compiler README from ISO/IEC 9899 Amendement 1 are partially supported on Solaris 2.6.	- C++. Unsupported features file. The features of the C++
Utilities	historical Solaris u unchanged; a new /usr/xpg4/bin. XPG4, SUS, or SUS variables should b	uired by POSIX.2, POSIX.2a, XPG4, SU tility behavior, the original Solaris vers version that is standard-conforming h For applications wishing to take advan Sv2 features, the PATH (sh or ksh) or p e set with /usr/xpg4/bin preceding cified by those specifications are found usr/ccs/bin.	sion of the utility is as been provided in ntage of POSIX.2, POSIX.2a, path (csh) environment any other directories in
Feature Test Macros	are desired beyond interfaces and head CAE), then it need standard. If the ap	s are used by applications to indicate a d those specified by the C standard. If a ders defined by a particular standard (only define the appropriate feature te plication is using interfaces and heade addition to defining the appropriate sta	an application uses only those such as POSIX or X/Open st macro specified by that rs not defined by that

Standards, Environments, and Macros 329

must also define __EXTENSIONS__. Defining __EXTENSIONS__ provides the application with access to all interfaces and headers not in conflict with the specified standard. The application must define __EXTENSIONS__ either at compile time or within the applicatio'n source files.

ANSI/ISO C

No feature test macros need to be defined to indicate that an application is a conforming C application.

ANSI/ISO C++

ANSI/ISO C++ does not define any feature test macros. If the standard C++ announcement macro __cplusplus is predefined to value 199711 or greater, the compiler operates in a standard-conforming mode, indicating C++ standards conformance. The value 199711 indicates conformance to ISO/IEC 14882:1998, as required by that standard. (As noted above, conformance to the standard is incomplete.) A standard-conforming mode is not available with compilers prior to Sun WorkShop C++ 5.0.

C++ bindings are not defined for POSIX or X/Open CAE, so specifying feature test macros such as _POSIX_SOURCE and _XOPEN_SOURCE can result in compilation errors due to conflicting requirements of standard C++ and those specifications.

POSIX

Applications that are intended to be conforming POSIX.1 applications must define the feature test macros specified by the standard before including any headers. For the standards listed below, applications must define the feature test macros listed. Application writers must check the corresponding standards for other macros that can be queried to determine if desired options are supported by the implementation.

POSIX Standard	Feature Test Macros
POSIX.1-1990	_POSIX_SOURCE
POSIX.1-1990 and POSIX.2-1992 C-Language Bindings Option	_POSIX_SOURCE and _POSIX_C_SOURCE=2
POSIX.1b-1993	_POSIX_C_SOURCE=199309L
POSIX.1c-1996	_POSIX_C_SOURCE=199506L

SVID3

The SVID3 specification does not specify any feature test macros to indicate that an application is written to meet SVID3 requirements. The SVID3 specification was written before the C standard was completed.

X/Open CAE

	specifications, use	e an application that conforms to one of the X/Open CAE the following guidelines. Applications need not set the POSIX is if they require both CAE and POSIX functionality.
	XPG3	The application must define _XOPEN_SOURCE with a value other than 500 (preferably 1).
	XPG4	The application must define _XOPEN_SOURCE with a value other than 500 (preferably 1) and set _XOPEN_VERSION=4.
	SUS (XPG4v2)	The application must define _XOPEN_SOURCE with a value other than 500 (preferably 1) and set _XOPEN_SOURCE_EXTENDED=1.
	SUSv2	The application must define _XOPEN_SOURCE=500.
Compilation	ANSI X3.159-1989 the cc and c89 ut and the compilation SPARC and x86 er /usr/ccs/lib/v	, SUS-, or SUSv2-conforming implementation must include an (ANSI C Language) standard-conforming compilation system and ilities. Solaris 7 through 9 were tested with the cc and c89 utilities in system provided by Sun WorkShop Compiler™ C 5.0 in the vironments. When cc is used to link applications, ralues-xpg4.0 must be specified on any link/load command line, way to build applications is described below.
		5-conforming application must include -1 XNS on any link/load ddition to defining the feature test macros specified for SUS or y.
	WorkShop Compidefine the macro feature), then the sproperly map function full support for IS	ppports the redefine_extname pragma feature (the Sun ler™ C 4.2 and Sun WorkShop Compiler™ C 5.0 compilers _PRAGMA_REDEFINE_EXTNAME to indicate that it supports this tandard headers use #pragma redefine_extname directives to tion names onto library entry point names. This mapping provides O C, POSIX, and X/Open namespace reservations. The Sun ler™ C 5.0 compiler was used for all branding and certification eases 7 through 9.
	directive to map ir In this instance, ap	ure is not supported by the compiler, the headers use the #define aternal function names onto appropriate library entry point names. plications should avoid using the explicit 64-bit file offset symbols (5) manual page, since these names are used by the implementation ative entry points.
		NorkShop Compiler [™] C 5.0, applications conforming to the d above should be compiled using the utilities and flags indicated in :
	Specification	Compiler/FlagBeature Test Macros
	ANSI/ISO C	c89 none

Standards, Environments, and Macros 331

Specification	Compiler/FlagReature Test Macros	
SVID3	cc -Xt	none
POSIX.1-1990	c89	_POSIX_SOURCE
POSIX.1-1990 and POSIX.2-1992 C-Language Bindings Option	c89	_POSIX_SOURCE and POSIX_C_SOURCE=2
POSIX.1b-1993	c89	_POSIX_C_SOURCE=199309L
POSIX.1c-1996	c89	_POSIX_C_SOURCE=199506L
CAE XPG3	cc -Xa	_XOPEN_SOURCE
CAE XPG4	c89	_XOPEN_SOURCE and _XOPEN_VERSION=4
SUS (CAE XPG4v2) (includes XNS4)	c89	_XOPEN_SOURCE and _XOPEN_SOURCE_EXTENDED=1
SUSv2 (includes XNS5)	c89	_XOPEN_SOURCE=500

For platforms supporting the LP64 (64-bit) programming environment where the SC5.0 Compilers have been installed, SUSv2–conforming LP64 applications using XNS5 library calls should be built with command lines of the form:

c89 \$(getconf XBS5_LP64_OFF64_CFLAGS) -D_XOPEN_SOURCE=500 \
 \$(getconf XBS5_LP64_OFF64_LDFLAGS) foo.c -o foo \
 \$(getconf XBS5_LP64_OFF64_LIBS) -lxnet

SEE ALSO | sysconf(3C), environ(5), lf64(5)

332 man pages section 5: Standards, Environments, and Macros • Last Revised 29 Aug 2001

NAME | sticky – mark files for special treatment

DESCRIPTION

The *sticky bit* (file mode bit 01000, see chmod(2)) is used to indicate special treatment of certain files and directories. A directory for which the sticky bit is set restricts deletion of files it contains. A file in a sticky directory may only be removed or renamed by a user who has write permission on the directory, and either owns the file, owns the directory, or is the super-user. This is useful for directories such as /tmp, which must be publicly writable, but should deny users permission to arbitrarily delete or rename the files of others.

If the sticky bit is set on a regular file and no execute bits are set, the system's page cache will not be used to hold the file's data. This bit is normally set on swap files of diskless clients so that accesses to these files do not flush more valuable data from the system's cache. Moreover, by default such files are treated as swap files, whose inode modification times may not necessarily be correctly recorded on permanent storage.

Any user may create a sticky directory. See chmod for details about modifying file modes.

FILES /tmp

- **SEE ALSO** chmod(1), chmod(2), chown(2), mkdir(2)
 - BUGS
- mkdir(2) will not create a directory with the sticky bit set.

term(5)

NAME | term – conventional names for terminals

DESCRIPTION

Terminal names are maintained as part of the shell environment in the environment variable TERM. See sh(1), profile(4), and environ(5). These names are used by certain commands (for example, tabs, tput, and vi) and certain functions (for example, see curses(3CURSES)).

Files under /usr/share/lib/terminfo are used to name terminals and describe their capabilities. These files are in the format described in terminfo(4). Entries in terminfo source files consist of a number of comma-separated fields. To print a description of a terminal *term*, use the command infocmp -I *term*. See infocmp(1M). White space after each comma is ignored. The first line of each terminal description in the terminfo database gives the names by which terminfo knows the terminal, separated by bar (|) characters. The first name given is the most common abbreviation for the terminal (this is the one to use to set the environment variable TERMINFO in \$HOME/.profile; see profile(4)), the last name given should be a long name fully identifying the terminal, and all others are understood as synonyms for the terminal name. All names but the last should contain no blanks and must be unique in the first 14 characters; the last name may contain blanks for readability.

Terminal names (except for the last, verbose entry) should be chosen using the following conventions. The particular piece of hardware making up the terminal should have a root name chosen, for example, for the AT&T 4425 terminal, att4425. This name should not contain hyphens, except that synonyms may be chosen that do not conflict with other names. Up to 8 characters, chosen from the set a through z and 0 through 9, make up a basic terminal name. Names should generally be based on original vendors rather than local distributors. A terminal acquired from one vendor should not have more than one distinct basic name. Terminal sub-models, operational modes that the hardware can be in, or user preferences should be indicated by appending a hyphen and an indicator of the mode. Thus, an AT&T 4425 terminal in 132 column mode is att4425–w. The following suffixes should be used where possible:

Suffix	Meaning	Example
-w	Wide mode (more than 80 columns)	att4425-w
-am	With auto. margins (usually default)	vt100–am
–nam	Without automatic margins	vt100–nam
- <i>n</i>	Number of lines on the screen	aaa–60
-na	No arrow keys (leave them in local)	c100–na
-np	Number of pages of memory	c100–4p

-rv	Reverse video	att4415–rv

To avoid conflicts with the naming conventions used in describing the different modes of a terminal (for example, -w), it is recommended that a terminal's root name not contain hyphens. Further, it is good practice to make all terminal names used in the terminfo(4) database unique. Terminal entries that are present only for inclusion in other entries via the use= facilities should have a '+' in their name, as in 4415+n1.

Here are some of the known terminal names: (For a complete list, enter the command ls -C /usr/share/lib/terminfo/?).

2621,hp2621	Hewlett-Packard 2621 series
2631	Hewlett-Packard 2631 line printer
2631-с	Hewlett-Packard 2631 line printer, compressed mode
2631-е	Hewlett-Packard 2631 line printer, expanded mode
2640,hp2640	Hewlett-Packard 2640 series
2645,hp2645	Hewlett-Packard 2645 series
3270	IBM Model 3270
33,tty33	AT&T Teletype Model 33 KSR
35,tty35	AT&T Teletype Model 35 KSR
37,tty37	AT&T Teletype Model 37 KSR
4000a	Trendata 4000a
4014,tek4014	TEKTRONIX 4014
40,tty40	AT&T Teletype Dataspeed 40/2
43,tty43	AT&T Teletype Model 43 KSR
4410,5410	AT&T 4410/5410 in 80-column mode, version 2
4410-nfk,5410-nfk	AT&T 4410/5410 without function keys, version 1
4410-nsl,5410-nsl	AT&T 4410/5410 without pln defined
4410-w,5410-w	AT&T 4410/5410 in 132-column mode
4410v1,5410v1	AT&T 4410/5410 in 80-column mode, version 1
4410v1-w,5410v1-w	AT&T 4410/5410 in 132-column mode, version 1
4415,5420	AT&T 4415/5420 in 80-column mode
4415–nl,5420–nl	AT&T 4415/5420 without changing labels

term(5)

4415-rv,5420-rv	AT&T 4415/5420 80 columns in reverse video
4415-rv-nl,5420-rv-nl	AT&T 4415/5420 reverse video without changing labels
4415-w,5420-w	AT&T 4415/5420 in 132-column mode
4415-w-nl,5420-w-nl	AT&T 4415/5420 in 132-column mode without changing labels
4415-w-rv,5420-w-rv	AT&T 4415/5420 132 columns in reverse video
4418,5418	AT&T 5418 in 80-column mode
4418-w,5418-w	AT&T 5418 in 132-column mode
4420	AT&T Teletype Model 4420
4424	AT&T Teletype Model 4424
4424-2	AT&T Teletype Model 4424 in display function group ii
4425,5425	AT&T 4425/5425
4425-fk,5425-fk	AT&T 4425/5425 without function keys
4425–nl,5425–nl	AT&T 4425/5425 without changing labels in 80-column mode
4425-w,5425-w	AT&T 4425/5425 in 132-column mode
4425-w-fk,5425-w-fk	AT&T 4425/5425 without function keys in 132-column mode
4425–nl–w,5425–nl–w	AT&T 4425/5425 without changing labels in 132-column mode
4426	AT&T Teletype Model 4426S
450	DASI 450 (same as Diablo 1620)
450–12	DASI 450 in 12-pitch mode
500,att500	AT&T-IS 500 terminal
510,510a	AT&T 510/510a in 80-column mode
513bct,att513	AT&T 513 bct terminal
5320	AT&T 5320 hardcopy terminal
5420_2	AT&T 5420 model 2 in 80-column mode
5420_2-w	AT&T 5420 model 2 in 132-column mode
5620,dmd	AT&T 5620 terminal 88 columns
5620–24,dmd–24	AT&T Teletype Model DMD 5620 in a 24x80 layer
5620–34,dmd–34	AT&T Teletype Model DMD 5620 in a 34x80 layer

336 man pages section 5: Standards, Environments, and Macros • Last Revised 3 Jul 1990

term(5)

610,610bct	AT&T 610 bct terminal in 80-column mode	
610-w,610bct-w	AT&T 610 bct terminal in 132-column mode	
630,630MTG	AT&T 630 Multi-Tasking Graphics terminal	
7300,pc7300,unix_pc	AT&T UNIX PC Model 7300	
735,ti	Texas Instruments TI735 and TI725	
745	Texas Instruments TI745	
dumb	generic name for terminals that lack reverse line-feed and other special escape sequences	
hp	Hewlett-Packard (same as 2645)	
lp	generic name for a line printer	
pt505	AT&T Personal Terminal 505 (22 lines)	
pt505–24	AT&T Personal Terminal 505 (24-line mode)	
sync	generic name for synchronous Teletype Model 4540-compatible terminals	

Commands whose behavior depends on the type of terminal should accept arguments of the form -T*term* where *term* is one of the names given above; if no such argument is present, such commands should obtain the terminal type from the environment variable TERM, which, in turn, should contain *term*.

FILES

SEE ALSO

sh(1), stty(1), tabs(1), tput(1), vi(1), infocmp(1M), curses(3CURSES), profile(4), terminfo(4), environ(5)

/usr/share/lib/terminfo/?/*

compiled terminal description database

vgrindefs(5)

NAME	vgrindefs – vgrind's language definition data base	
SYNOPSIS	SIS /usr/lib/vgrindefs	
DESCRIPTION	vgrindefs contains all language definitions for vgrind(1). Capabilities in vgrindefs are of two types: Boolean capabilities which indicate that the language has some particular feature and string capabilities which give a regular expression or keyword list. Entries may continue onto multiple lines by giving a \ as the last character of a line. Lines starting with # are comments.	

Capabilities The following table names and describes each capability.

Name	Туре	Description
ab	str	Regular expression for the start of an alternate form comment
ae	str	Regular expression for the end of an alternate form comment
bb	str	Regular expression for the start of a block
be	str	Regular expression for the end of a lexical block
cb	str	Regular expression for the start of a comment
се	str	Regular expression for the end of a comment
id	str	String giving characters other than letters and digits that may legally occur in identifiers (default '_')
kw	str	A list of keywords separated by spaces
lb	str	Regular expression for the start of a character constant
le	str	Regular expression for the end of a character constant
oc	bool	Present means upper and lower case are equivalent
pb	str	Regular expression for start of a procedure
pl	bool	Procedure definitions are constrained to the lexical level matched by the 'px' capability
px	str	A match for this regular expression indicates that procedure definitions may occur at the next lexical level. Useful for lisp-like languages in which procedure definitions occur as subexpressions of defuns.
sb	str	Regular expression for the start of a string
se	str	Regular expression for the end of a string
tc	str	Use the named entry as a continuation of this one
tl	bool	Present means procedures are only defined at the top lexical level

Regular Expressions	vgrindefs uses regular expressions similar to those of $ex(1)$ and $lex(1)$. The characters ' $^{\prime}$, ' $^{\prime}$, ' $^{\prime}$, and ' $^{\prime}$ are reserved characters and must be 'quoted' with a preceding \setminus if they are to be included as normal characters. The metasymbols and their meanings are:	
	\$	The end of a line
	*	The beginning of a line
	∖d	A delimiter (space, tab, newline, start of line)
	∖a	Matches any string of symbols (like '.*' in lex)
	/p	Matches any identifier. In a procedure definition (the 'pb' capability) the string that matches this symbol is used as the procedure name.
	()	Grouping
	1	Alternation
	?	Last item is optional
	\e Preceding any string means that the string will not match an inpustring if the input string is preceded by an escape character (\). This is typically used for languages (like C) that can include the string delimiter in a string by escaping it.	
	Unlike other regular expressions in the system, these match words and not characters. Hence something like '(tramp steamer)flies?' would match 'tramp', 'steamer', 'trampflies', or 'steamerflies'. Contrary to some forms of regular expressions, vgrindef alternation binds very tightly. Grouping parentheses are likely to be necessary in expressions involving alternation.	
Keyword List	The keyword list is just a list of keywords in the language separated by spaces. If the 'oc' boolean is specified, indicating that upper and lower case are equivalent, then all the keywords should be specified in lower case.	
EXAMPLES	EXAMPLE 1 A sample program.	
	The following entry, which describes the C language, is typical of a language entry.	
	<pre>C c the C programming language:\ :pb=^\d?*?\d?\p\d?(\a?\)(\d {):bb={:be=}:cb=/*:ce=*/:sb=":se=\e":\ :le=\e':tl:\ :kw=asm auto break case char continue default do double else enum\ extern float for fortran goto if int long register return short\ sizeof static struct switch typedef union unsigned void while #define\ #else #endif #if #ifdef #ifndef #include #undef # define endif\ ifdef ifndef include undef defined: Note that the first field is just the language name (and any variants of it). Thus the C</pre>	
	language could be specified to vgrind(1) as 'c' or 'C'.	

vgrindefs(5)

FILES	/usr/lib/vgrindefs	file containing vgrind descriptions
SEE ALSO	ex(1), lex(1), troff(1), vgr	ind(1)

NAME v

wbem - Web-Based Enterprise Management

DESCRIPTION Web-Based Enterprise Management (WBEM) is a set of management and Internet-related technologies intended to unify the management of enterprise computing environments. Developed by the Distributed Management Task Force (DMTF), WBEM enables organizations to deliver an integrated set of standards-based management tools that support and promote World Wide Web technology. The DMTF has developed a set of standards that make up WBEM. This set of standards includes:

> Common Information Model (CIM) - an object-oriented data model that describes the overall management of information in an enterprise network environment. CIM consists of a CIM specification and a CIM schema:

	CIM Specification	Consists of the language and methodology that describes management data.
	CIM Schema	Provides actual model descriptions of systems, applications, large area networks, and devices. The CIM Schema enables applications from different developers on different platforms to describe management data in a standard format. As a result, a variety of management applications can share this information.
•	1 91	ext Transport Protocol (HTTP) 1.1 is a transport erations to HTTP to allow implementations of CIM ndardized manner.
	1	.1 uses eXtensible Markup Language (XML), which is sents management information in textual form.

In addition to the XML representation, CIM information is also represented textually by the managed object format (MOF). These MOF representations are typically stored as text files that developers compile into a CIM Object Manager.

EXTENDED DESCRIPTION WBEM Tools and Services	Tools and services that enable developers to create and Services management applications and instrumentation that manage heterogeneous computer environments include:
	Solaris WBEM Services 2.5Solaris WBEM Software Development Kit 2.5
Solaris WBEM Services 2.5	These services consist of a set of value-added Services 2.5 components. These services make it easier for developers to create management applications that run in the Solaris operating environment. They also make the Solaris operating environment easier to manage. Solaris WBEM Services 2.5 consists of:
	 CIM Object Manager, CIM Repository, and MOF Compiler

wbem(5)

	 CIM and Solaris Schema, which is an extension schema of CIM. CIM and Solaris Schema is a collection of CIM classes that describe managed elements in the Solaris operating environment. These classes are available from the CIM Object Manager at start up.
	 Solaris Providers, which are programs that communicate information between the Solaris operating environment and the CIM Object Manager (providers get and set "dynamic" information about managed elements, acting as an intermediary between the CIM Object Manager and the managed elements).
	Solaris software providers have been developed for a variety of areas: users, roles, file systems, and network configuration, for example. A remote provider is also available to distribute agents away from the CIM Object Manager when required. Because of the incremental development capabilities of the WBEM instrumentation framework, developers can progressively and consistently add more providers for additional Solaris software services.
	 SNMP Adapter for WBEM, which enables Simple Network Management Protocol (SNMP) management applications to access system management information that is provided by Solaris WBEM Services. Used with the Solstice Enterprise Agent (SEA) Master Agent snmpdx(1M), the SNMP Adapter for WBEM maps SNMP requests into equivalent WBEM Common Information Model (CIM) properties or instances.
	The SNMP Adapter for WBEM also remaps the response from the CIM Object Manager into an SNMP response, which is returned to the management application.
	A mapping file contains the corresponding Object Identifier (OID), class name, property name, and Abstract Syntax Notation One (ASN.1) type for each object. Developers can create their own mapping files.
	• SNMP Provider, which enables WBEM services to deliver SNMP information.
Solaris WBEM Software Development Kit 2.5	This kit consists of a set of key application Software development tools that make it easier for developers to write management applications that can communicate with any WBEM-enabled management device. The Solaris WBEM Software Development Kit includes examples, documentation, and CIM Workshop, a graphical user interface through which developers can view and create classes and instances, through the remote method invocation (RMI) or the XML/HTTP protocol.
	Developers can also use this kit to write providers, which are programs that communicate with managed elements to access data.
	All management applications that developers create with the Solaris WBEM Software Development Kit run on the Java platform. The Solaris 9 WBEM Software Development Kit installs and runs in version 1.4 of the Java environment. Developers can use the kit to write standalone applications or applications that run in conjunction with Solaris WBEM Services.

wbem(5)

The Solaris WBEM Software Development Kit is described in the *Solaris WBEM SDK Development Guide*. Javadoc for the WBEM application programming interface is located at /usr/sadm/lib/wbem/doc/index.html.

Compatibility of Solaris WBEM Services with Existing Protocols

Adapters and converters enable Solaris WBEM Services of Solaris to work compatibly with existing protocols by mapping WBEM information to these protocols. One such protocol is Simple Network Management Protocol (SNMP).

Legacy management applications can administer WBEM-enabled software in the Solaris operating environment. Developers can write agents or providers that convert information from these protocols to WBEM, and they can write adapters that convert WBEM information into these protocols.

ATTRIBUTES

See attributes(5) for descriptions of the following attributes:

ATTRIBUTE TYPE	ATTRIBUTE VALUE
Availability	SPARC and x86
Architecture	SUNWwbapi, SUNWwbcor, SUNWwbcou, SUNWwbdev, SUNWwbdoc, SUNWwbpro
CSI	Enabled

SEE ALSO

appletviewer(1), cimworkshop(1M), init.wbem(1M), mofcomp(1M), mofreg(1M), snmpdx(1M), wbemadmin(1M), wbemconfig(1M), wbemlogviewer(1M), attributes(5) wbem(5)

Index

Α

account management PAM module for projects – pam_projects, 285 ANSI — standards and specifications supported by Solaris, 328 architecture — attributes of interfaces, 18 ascii — ASCII character set, 16 attributes - attributes of interfaces, 18 attributes of interfaces - architecture, 18 attributes of interfaces - attributes, 18 attributes of interfaces - availability, 18 attributes of interfaces — CSI, 18 attributes of interfaces - MT-Level, 18 attributes of interfaces - stability, 18 attributes — characteristics of commands, utilities, and device drivers Architecture, 18 Availability, 18 Interface Stability, 19 MT-Level, 22 authentication, account, session, and password management PAM modules for Kerberos V5 — pam_krb5, 273 authentication and password management module — pam_authtok_check, 265 authentication and password management module — pam_authtok_get, 267 authentication module for password pam_passwd_auth, 283 availability - attributes of interfaces, 18

С

C — standards and specifications supported by Solaris, 328 C++ — standards and specifications supported by Solaris, 328 character set description file — charmap, 26 charmap — character set description file, 26 Decimal Constants, 28 Declarations, 26 Format, 27 Ranges of Symbolic Names, 28 Symbolic Names, 26 code set conversion tables — iconv, 122 code set conversion tables — iconv_1250, 107 code set conversion tables — iconv_1251, 113 code set conversion tables — iconv_646, 126 code set conversion tables — iconv_852, 129 code set conversion tables — iconv_8859-1, 135 code set conversion tables — iconv_8859-2, 141 code set conversion tables — iconv_8859-5, 147 code set conversion tables — iconv dhn, 155 code set conversion tables — iconv_koi8-r, 159 iconv_mac_cyr, 167 code set conversion tables — iconv_maz, 175 code set conversion tables — iconv_pc_cyr, 179 iconv_unicode, 185 code set conversion tables — iconv_1250, 107 — iconv_1251, 113 — iconv_646, 126 — iconv_852, 129

code set conversion tables (Continued) — iconv_8859-1, 135 — iconv_8859-2, 141 — iconv_8859-5, 147 — iconv_dhn, 155 — iconv_koi8-r, 159 — iconv_mac_cyr, 167 — iconv_maz, 175 — iconv_pc_cyr, 179 compilation environment, transitional — Ifcompile64, 206 crypt_unix — traditional UNIX crypt algorithm, 32 CSI — attributes of interfaces, 18

D

data storage modules for the DHCP service dhcp_modules, 35 dhcp — Dynamic Host Configuration Protocol, 33 dhcp_modules --- data storage modules for the DHCP service, 35 document production man — macros to format manual pages, 238 mansun - macros to format manual pages, 242 me - macros to format technical papers, 246 mm — macros to format articles, theses and books, 251 ms - macros to format articles, theses and books, 258 Dynamic Host Configuration Protocol dhcp, 33

Ε

environ — user environment, 37 environment variables HOME, 37 LANG, 37 LC_COLLATE, 37 LC_CTYPE, 37 LC_MESSAGES, 37 LC_MONETARY, 37 environment variables (Continued) LC_NUMERIC, 37 LC_TIME, 37 MSGVERB, 37 NETPATH, 37 PATH, 37 SEV_LEVEL, 37 TERM, 37 TZ, 37 eqnchar — special character definitions for eqn, 42 extended file attributes — fsattr, 96 extensions — localedef extensions description file, 43

F

file format notation — formats, formats, 91 file name pattern matching — fnmatch, 63 filesystem — file system layout, 44 filesystem — file system organization, 44 Root File System, 44 /usr File System, 53 fnmatch — file name pattern matching, 63 fns — overview of FNS, 67 Composite Names, 67 FNS and Naming Systems, 68 FNS overview - fns, 67 overview of FNS References fns_references, 85 overview over DNS implementation fns_dns, 69 overview over files implementation fns_files, 71 overview over NIS+ implementation ---fns_nis+, 77 overview over NIS (YP) implementation fns_nis, 79 overview over X.500 implementation fns_x500, 88 fns - overview of FNS Why FNS?, 67 XFN, 67 fns_dns - overview of FNS over DNS implementation, 69

346 man pages section 5: Standards, Environments, and Macros • December 2003

fns_files — overview of FNS over files implementation, 71 fns_files - overview of FNS over /etc files implementation, FNS Policies and /etc Files, 71 fns_initial_context - overview of the FNS Initial Context, 73 fns_nis - overview of FNS over NIS (YP) implementation, 79 Federating NIS with DNS or X.500, 79 FNS Policies and NIS, 79 NIS Security, 79 fns_nis+ - overview of FNS over NIS+ implementation, 77 FNS Policies and NIS+, 77 fns_policies - overview of the FNS Policies, 81 fns_references - overview of FNS References, 85 Address Types, 85 Reference Types, 85 fns_x500 — overview of FNS over X.500 implementation, 88 formats - file format notation, 91 fsattr — extended file attributes, 96

L

iconv — code set conversion tables, 122 iconv_1250 — code set conversion tables for MS 1250 (Windows Latin 2), 107 iconv_1251 — code set conversion tables for MS 1251 (Windows Cyrillic), 113 iconv_646 — code set conversion tables for ISO 646, 126 iconv_852 - code set conversion tables for MS 852 (MS-DOS Latin 2), 129 iconv_8859-1 — code set conversion tables for ISO 8859-1 (Latin 1), 135

iconv_8859-2 - code set conversion tables for ISO 8859-2 (Latin 2), 141

iconv_8859-5 -- code set conversion tables for ISO 8859-5 (Cyrillic), 147

iconv_dhn — code set conversion tables for DHN (Dom Handlowy Nauki), 155

iconv_koi8-r - code set conversion tables for KOI8-R, 159

iconv_mac_cyr — code set conversion tables for Macintosh Cyrillic, 167

- iconv_maz code set conversion tables for Mazovia, 175
- iconv_pc_cyr code set conversion tables for Alternative PC Cyrillic, 179
- iconv_unicode -- code set conversion tables for Unicode, 185
- internationalized basic and extended regular expression matching — regex, 304
- isalist the native instruction sets known to Solaris, 190
- ISO standards and specifications supported by Solaris, 328

L large file status of utilities — largefile, 192 largefile — large file status of utilities, 192 Large file aware utilities, 192 Large file safe utilities, 194 lf64 — transitional interfaces for 64-bit file offsets, 196 Data Types, 196 System Interfaces, 198 lfcompile — large file compilation environment, Access to Additional Large File Interfaces, 203 lfcompile64 — transitional compilation environment, 206 Access to Additional Large File Interfaces, 206 live_upgrade, 208 locale — subset of a user's environment that depends on language and cultural conventions, 212 collating-element keyword, 220 collating-symbol keyword, 221 Collation Order, 222 LC_COLLATE, 219 LC_CTYPE, 215 LC_MESSAGES, 236 LC_MONETARY, 225 LC_NUMERIC, 229 LC_TIME, 231 LC_TIME C-language Access, 233

LC_TIME General Information, 236

Index 347

locale — subset of a user's environment that depends on language and cultural conventions (Continued) Locale Definition, 212

order_end keyword, 225 order_start keyword, 221 localedef extensions description file extensions, 43

Μ

macros mm, 251 ms, 258 to format Manual pages — man, 238 to format Manual pages - mansun, 242 to format technical papers - me, 246 man — macros to format manual pages, 238 mansun - macros to format manual pages, 242 manual pages macros to format manual pages - man, 238 Sun macros to format manual pages mansun, 242 mark files for special treatment - sticky, 333 me — macros to format technical papers, 246 mm — macros to format articles, theses and books, 251 ms - macros to format articles, theses and books, 258 MT-Level — attributes of interfaces, 18

Ν

native instruction sets known to Solaris isalist, 190 NFS and sticky bits — sticky, 333 nfssec — overview of NFS security modes, 263

0

overview of FNS - fns, 67

overview of FNS over DNS implementation fns dns, 69 overview of FNS over files implementation fns files, 71 overview of FNS over NIS+ implementation ---fns_nis+, 77 overview of FNS over NIS (YP) implementation — fns_nis, 79 overview of FNS over X.500 implementation fns_x500, 88 overview of FNS References fns_references, 85 overview of NFS security modes - nfssec, 263 overview of the FNS Initial Context fns_initial_context, 73 overview of the FNS Policies - fns_policies, 81

Ρ

PAM account management module for UNIX pam_unix_account, 296 pam_authtok_check — authentication and password management module, 265 pam_authtok_get — authentication and password management module, 267 pam_authtok_store — password management module, 269 pam_dial_auth — authentication management for dialups, 272 pam_krb5 — authentication, account, session and password management for Kerberos V5, 273 pam_krb5 — authentication, account, session, and password management PAM modules for Kerberos V5, 273 pam_passwd_auth — authentication module for password, 283 pam_projects — account management PAM module for projects, 285 pam_rhosts_auth - authentication management using ruserok(), 286 pam_sample — sample module for PAM, 289 pam_unix — authentication, account, session and password management for UNIX, 293 pam_unix_account — PAM account management module for UNIX, 296

348 man pages section 5: Standards, Environments, and Macros • December 2003

pam_unix_session — session management PAM module for UNIX, 299
password management module — pam_authtok_store, 269
POSIX — standards and specifications supported by Solaris, 328
POSIX.1 — standards and specifications supported by Solaris, 328
POSIX.2 — standards and specifications supported by Solaris, 328
POSIX.2 — standards and specifications supported by Solaris, 328
POSIX.2 — standards and specifications supported by Solaris, 328

R

rbac — role-based access control, 301
regex — internationalized basic and extended regular expression matching, 304
regular expression compile and match routines — advance, 313
— compile, 313
— regexp, 313
— step, 313
role-based access control — rbac, 301

S

ftp — authentication system, 320 session management PAM module for UNIX ---pam_unix_session, 299 sgml - Standard Generalized Markup Language, 322 sgml - Standard Generalized Markup language, RefEntry, 322 sgml — Standard Generalized Markup Language RefMeta, 322 RefNameDiv, 323 RefSect1, 324 RefSect2, 324 RefSynopsisDiv, 324 shell environment, conventional names for terminals - term, 334 solbook - Standard Generalized Markup Language, 322

special character definitions for eqn eqnchar, 42 stability - attributes of interfaces, 18 Standard Generalized Markup Language — sgml, 322 — solbook, 322 standards - standards and specifications supported by Solaris, 328 standards and specifications supported by Solaris — ANSI, 328 standards and specifications supported by Solaris — C++, 328 standards and specifications supported by Solaris — C, 328 standards and specifications supported by Solaris – ISO, 328 standards and specifications supported by Solaris — POSIX.1, 328 standards and specifications supported by Solaris — POSIX.2, 328 standards and specifications supported by Solaris — POSIX, 328 standards and specifications supported by Solaris — standards, 328 standards and specifications supported by Solaris — SUS, 328 standards and specifications supported by Solaris — SUSv2, 328 standards and specifications supported by Solaris — SVID3, 328 standards and specifications supported by Solaris — SVID, 328 standards and specifications supported by Solaris — XNS4, 328 standards and specifications supported by Solaris — XNS5, 328 standards and specifications supported by Solaris — XNS, 328 standards and specifications supported by Solaris — XPG3, 328 standards and specifications supported by Solaris — XPG4, 328 standards and specifications supported by Solaris — XPG4v2, 328 standards and specifications supported by Solaris — XPG, 328 sticky — mark files for special treatment, 333 subset of a user's environment that depends on language and cultural conventions locale, 212

- SUS standards and specifications supported by Solaris, 328
- SUSv2 standards and specifications supported by Solaris, 328
- SVID standards and specifications supported by Solaris, 328

SVID3 — standards and specifications supported by Solaris, 328

Т

term — conventional names for terminals, 334 terminals, conventional names — term, 334 traditional UNIX crypt algorithm crypt_unix, 32 transitional compilation environment lfcompile64, 206 transitional interfaces for 64-bit file offsets lf64, 196

U

unicode, code set conversion tables iconv_unicode, 185 user environment, — environ, 37

V

vgrindefs - vgrind language definitions, 338

W

wbem — Web-Based Enterprise Management, 341 Web-Based Enterprise Management wbem, 341

X

- XNS standards and specifications supported by Solaris, 328
- XNS4 standards and specifications supported by Solaris, 328
- XNS5 standards and specifications supported by Solaris, 328
- XPG standards and specifications supported by Solaris, 328
- XPG3 standards and specifications supported by Solaris, 328
- XPG4 standards and specifications supported by Solaris, 328
- XPG4v2 standards and specifications supported by Solaris, 328

350 man pages section 5: Standards, Environments, and Macros • December 2003